

Most human things go in pairs.

Alcmaeon, ~ 450 BC

true	false	good	bad
right	left	up	down
front	back	future	past
light	dark	hot	cold
matter	antimatter	boson	fermion

How can we formalize a *general* concept of duality? The Chinese tried yin-yang theory, which inspired Leibniz to develop binary notation, which in turn underlies digital computation!

But what's the state of the art *now*?

In category theory the fundamental duality is the act of reversing an arrow:

$$\bullet \rightarrow \bullet \quad \rightsquigarrow \quad \bullet \leftarrow \bullet$$

We use this to model switching past and future, false and true, small and big...

Every category \mathcal{C} has an *opposite* \mathcal{C}^{op} , where the arrows are reversed. This is a symmetry of the category of categories:

$$\text{op}: \text{Cat} \rightarrow \text{Cat}$$

and indeed the only nontrivial one:

$$\text{Aut}(\text{Cat}) = \mathbb{Z}/2$$

In logic, the simplest duality is negation. It's order-reversing:

if P implies Q then $\neg Q$ implies $\neg P$

and—ignoring intuitionism!—it's an involution:

$$\neg\neg P = P$$

Thus if \mathcal{C} is a category of propositions and proofs, we expect a functor:

$$\neg: \mathcal{C} \rightarrow \mathcal{C}^{\text{op}}$$

with a natural isomorphism:

$$\neg^2 \cong 1_{\mathcal{C}}$$

This has *two* analogues in quantum theory.

One shows up already in the category of finite-dimensional vector spaces, $\mathbf{FinVect}$. Every vector space V has a dual V^* . Taking the dual is contravariant:

$$\text{if } f: V \rightarrow W \text{ then } f^*: W^* \rightarrow V^*$$

and—ignoring infinite-dimensional spaces!—it's an involution:

$$V^{**} \cong V$$

This kind of duality is captured by the idea of a $*$ -autonomous category.

Recall that a *symmetric monoidal category* is roughly a category \mathcal{C} with a unit object $I \in \mathcal{C}$ and a tensor product

$$\otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$$

that is unital, associative and commutative up to coherent natural isomorphisms.

A symmetric monoidal category \mathcal{C} is *closed* if for every object $A \in \mathcal{C}$ the functor

$$X \mapsto A \otimes X$$

has a right adjoint, called the *internal hom*:

$$X \mapsto A \multimap X$$

In other words, there's a natural bijection

$$\text{hom}(A \otimes X, Y) \cong \text{hom}(X, A \multimap Y)$$

In FinVect , $A \otimes X$ is the usual tensor product, and $A \multimap Y$ is the vector space of linear maps from A to Y .

In classical logic, $A \otimes X$ is the proposition $A \& X$, and $A \multimap Y$ is the proposition $A \Rightarrow Y$.

A symmetric monoidal closed category \mathcal{C} is **-autonomous* if it has an object \perp such that if we set

$$X^* = X \multimap \perp$$

then the canonical morphism $X \rightarrow X^{**}$ is an isomorphism. We call \perp a *dualizing object*.

Any *-autonomous category comes with a functor:

$$*: \mathcal{C} \rightarrow \mathcal{C}^{\text{op}}$$

and a natural isomorphism:

$$*^2 \cong 1_{\mathcal{C}}$$

In $\mathbf{FinVect}$ we take $\perp = \mathbb{C}$, so

$$X^* = X \multimap \perp$$

is the usual dual of the vector space X . Here \perp is the unit for the tensor product. Whenever this happens we say \mathcal{C} is *compact*. In a compact category, $A \multimap X \cong A^* \otimes X$.

In classical logic we take $\perp = \text{FALSE}$, so

$$X^* = X \multimap \perp$$

is the usual negation $\neg X$. Here \perp is not the unit for the tensor product: the unit for $\&$ is TRUE .

In the category of finite-dimensional Hilbert spaces, we also have a second kind of duality. Besides the duals for objects with

$$\text{hom}(A \otimes X, Y) \cong \text{hom}(X, A^* \otimes Y)$$

we have duals for morphisms: given $f: X \rightarrow Y$ we have $f^\dagger: Y \rightarrow X$ with

$$\langle fx, y \rangle = \langle x, f^\dagger y \rangle$$

for all $x \in X$ and $y \in Y$.

In short, *a Hilbert space is like a miniature category*, with an amplitude to go from x to y instead of a set of ways to go from X to Y .

Duality for morphisms is captured by the concept of a *dagger category*: one where every morphism

$$f: X \rightarrow Y$$

gives a morphism

$$f^\dagger: Y \rightarrow X$$

with

$$(ST)^\dagger = T^\dagger S^\dagger \quad T^{\dagger\dagger} = T$$

Any dagger category comes with a functor

$$\dagger: \mathcal{C} \rightarrow \mathcal{C}^{\text{op}}$$

and a natural isomorphism

$$\dagger^2 \cong 1_{\mathcal{C}}$$

Duality for objects and morphisms fit together in the concept of a *dagger-compact category*.

This is a compact dagger category for which all relevant natural isomorphisms are *unitary* ($u^\dagger = u^{-1}$), and the dagger of the *unit*

$$X^* \otimes X \rightarrow I$$

is the *counit*

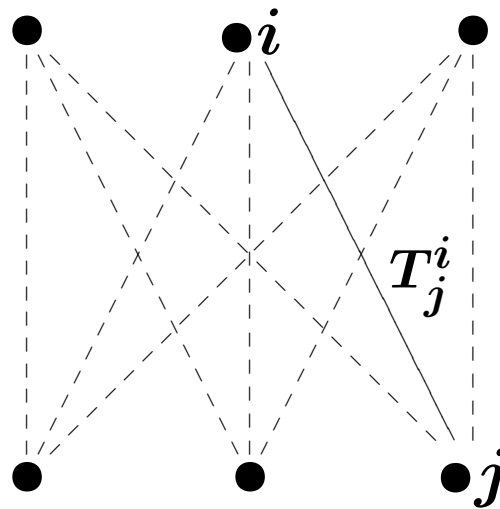
$$I \rightarrow X \otimes X^*$$

followed by the isomorphism

$$X \otimes X^* \xrightarrow{\sim} X^* \otimes X$$

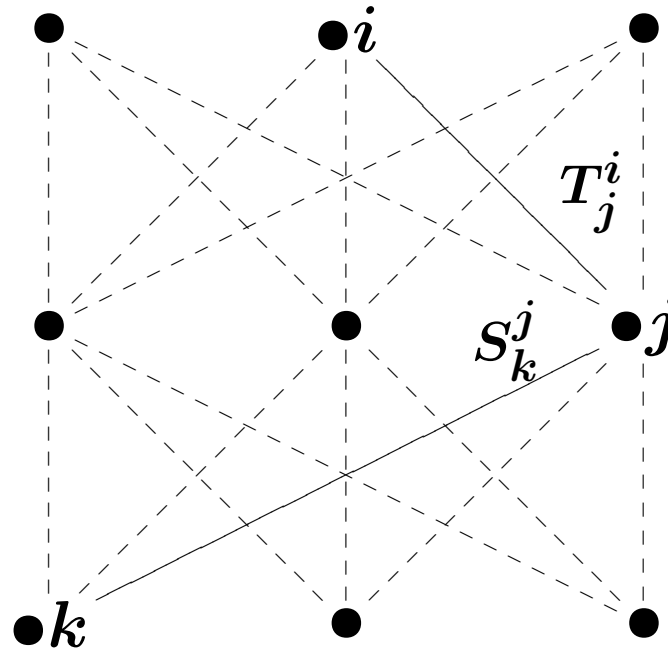
Dagger-compact categories are deeply related to ‘matrix mechanics’.

Heisenberg used matrices with complex entries to describe processes in quantum mechanics:



For each input state i and output state j , the process T gives a complex number $T_j^i \in \mathbb{C}$, the *amplitude* to go from i to j .

To compose processes, we sum over paths:



$$(ST)_k^i = \sum_j S_k^j \times T_j^i$$

In the continuum limit, such sums become *path integrals*.

Matrix mechanics also works with other rigs (= rings without negatives):

- $[0, \infty)$ with its usual $+$ and \times

Now T_j^i gives the *probability* to go from i to j .

- $\{\text{TRUE}, \text{FALSE}\}$ with \vee as $+$ and $\&$ as \times

Now T_j^i gives the *possibility* to go from i to j .

- $\mathbb{R}^{\text{MIN}} = \mathbb{R} \cup \{+\infty\}$ with MIN as $+$ and $+$ as \times

Now T_j^i gives the *action* to go from i to j .

In the continuum limit, matrix mechanics using \mathbb{R}^{MIN} gives the *principle of least action* in classical mechanics — see work on ‘idempotent analysis’ and ‘tropical algebra’.

For any commutative rig R , there is a compact category $\text{Mat}(R)$ where:

- objects are natural numbers;
- morphisms $T: m \rightarrow n$ are $n \times m$ matrices with entries in R .

If R is also a $*$ -rig:

$$(a + b)^* = a^* + b^*, \quad (ab)^* = b^* a^*, \quad a^{**} = a$$

then $\text{Mat}(R)$ is dagger-compact, with

$$(T^\dagger)_j^i = (T_i^j)^*$$

Every rig becomes a $*$ -rig with $a^* = a$. So, $\text{Mat}(\mathbb{R}^{\text{MIN}})$ is dagger-compact and *we can apply dagger-compact categories to classical mechanics!*

$\text{Mat}(\mathbb{R})$ is equivalent to the category of finite-dimensional real Hilbert spaces.

$\text{Mat}(\mathbb{C})$ is equivalent to the category of finite-dimensional complex Hilbert spaces.

$\text{Mat}(\mathbb{H})$ is equivalent to the category of finite-dimensional quaternionic Hilbert spaces.

The first two are dagger-compact. The third is not monoidal but still a dagger-category. *Why does Nature love complex Hilbert spaces best?*

My answer:

In fact Nature loves all three, each in its own way!

To understand this, it helps to equip our Hilbert spaces with more structure.

Start with a compact topological group G . Let $\text{Rep}(G)$ be the category of continuous unitary representations of G on finite-dimensional *complex* Hilbert spaces.

This category is dagger-compact. It's crucial in quantum physics, where G describes symmetries!

But Freeman Dyson's 'Threefold Way' reveals that continuous unitary representations of G on finite-dimensional *real* and *quaternionic* Hilbert spaces are *also* contained in $\text{Rep}(G)$!

We can ask if an object $X \in \text{Rep}(G)$ is its own dual. If X is *irreducible* — not a direct sum of other objects in a non-trivial way — there are three mutually exclusive choices:

- $X \not\cong X^*$.
- $X \cong X^*$ and X is *real*: it comes from a representation of G on a real Hilbert space H :

$$X = \mathbb{C} \otimes_{\mathbb{R}} H$$

- $X \cong X^*$ and X is *quaternionic*: it comes from a representation of G on a quaternionic Hilbert space H :

$$X = \text{the underlying complex space of } H$$

Why? If X is irreducible there's a 1-dimensional space of morphisms $f: X \rightarrow X$. So if $X \cong X^*$ there's a 1d space of morphisms $f: X \rightarrow X^*$, and thus a 1d space of morphisms

$$g: X \otimes X \rightarrow \mathbb{C}$$

But

$$X \otimes X \cong S^2 X \oplus \Lambda^2 X$$

so either there exists a nonzero g that is *symmetric*:

$$g(x, y) = g(y, x)$$

or one that is *skew-symmetric*:

$$g(x, y) = -g(y, x)$$

One or the other, not both!

Either way, we can write

$$g(x, y) = \langle jx, y \rangle$$

for some real-linear operator $j: X \rightarrow X$ with

$$ji = -ij$$

If g is symmetric, we can rescale j to achieve

$$j^2 = 1$$

Then j acts like complex conjugation, so

$$H = \{x \in X : jx = x\}$$

is a real representation of G with:

$$X = \mathbb{C} \otimes_{\mathbb{R}} H$$

If g is skew-symmetric, we can rescale j to achieve

$$j^2 = -1$$

Then i , j and $k = ij$ act like the quaternions:

$$i^2 = j^2 = k^2 = ijk = -1$$

so we obtain a representation of G on a quaternionic Hilbert space H with:

$X =$ the underlying complex space of H

Why can we rescale j to achieve either $j^2 = \pm 1$ but not both?

Suppose g is symmetric, so that

$$\langle jx, y \rangle = g(x, y) = g(y, x) = \langle jy, x \rangle$$

Then

$$\langle j^2x, x \rangle = \langle jx, jx \rangle \geq 0$$

so $j^2 > 0$. Letting $J = \alpha j$ with $\alpha \in \mathbb{R}$ we have

$$J^2 = \alpha^2 j^2$$

so by correctly choosing α we can achieve

$$J^2 = 1$$

Similarly, if g is skew-symmetric $j^2 < 0$ and we can rescale j to achieve $J^2 = -1$.

An example: $G = \text{SU}(2)$. Every object in $\text{Rep}(\text{SU}(2))$ is self-dual. There is one irreducible representation for each ‘spin’ $j = 0, \frac{1}{2}, 1, \dots$

If $j \in \mathbb{Z}$, the spin- j rep is real.

If $j \in \mathbb{Z} + \frac{1}{2}$, the spin- j rep is quaternionic.

For example: the spin-1 or *vector* rep on \mathbb{C}^3 is the complexification of a real rep on \mathbb{R}^3 .

But the spin- $\frac{1}{2}$ or *spinor* rep on \mathbb{C}^2 is the underlying complex rep of a quaternionic rep on \mathbb{H} .

In short: spin- $\frac{1}{2}$ particles are quaternions!

But what is the physical meaning of j ? For any $iJ \in \mathfrak{su}(2)$:

$$j \exp(i\theta J) = \exp(i\theta J) j$$

and thus

$$j iJ = iJ j$$

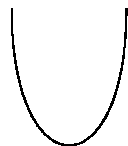
but $ij = -ji$ so

$$jJj^{-1} = -J$$

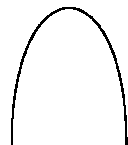
So: j reverses the angular momentum J .

This calculation works for all reps of $SU(2)$. Thus the meaning of j is *time reversal*.

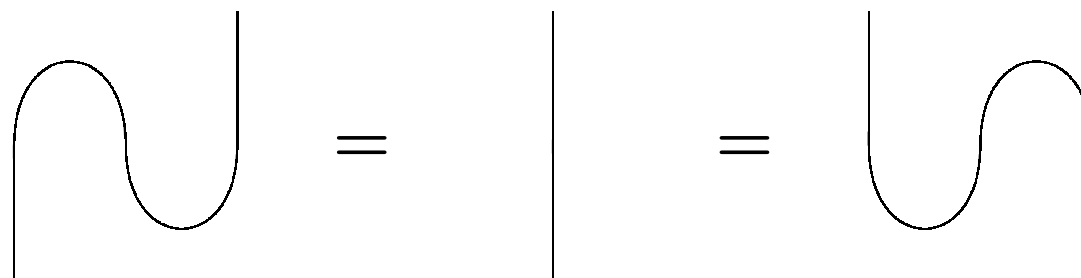
For any self-dual $X \in \text{Rep}(G)$ we can write $g: X \otimes X \rightarrow \mathbb{C}$ as a ‘cup’:



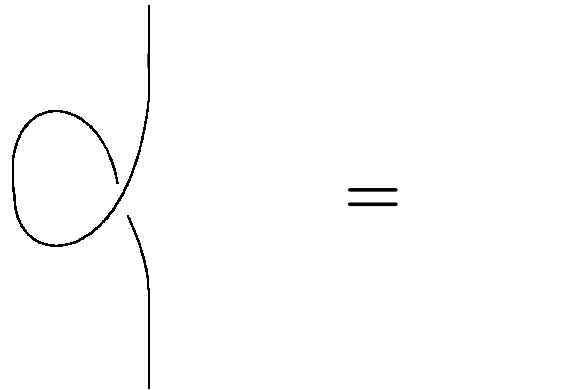
Since g defines an isomorphism $X \cong X^*$ we have a corresponding ‘cap’:



satisfying the *zig-zag identities*:

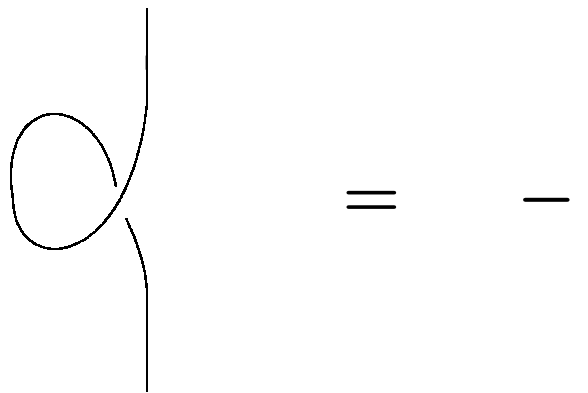


For a real self-dual object ($j^2 = 1$) we have



A diagrammatic equation showing a loop on a vertical line equal to a single vertical line. The loop is formed by a vertical line that curves to the left, forms a loop, and then continues as a vertical line. This is set equal to a single vertical line.

but for a quaternionic one ($j^2 = -1$) we have



A diagrammatic equation showing a loop on a vertical line equal to the negative of a single vertical line. The loop is formed by a vertical line that curves to the left, forms a loop, and then continues as a vertical line. This is set equal to a vertical line with a minus sign in front of it.

This is the behavior we expect from *bosons* and *fermions*, respectively.

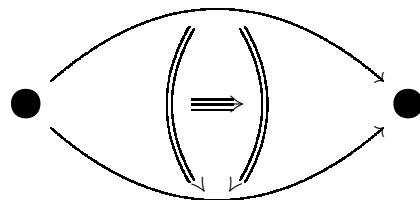
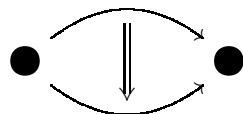
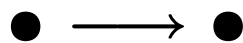
THE THREEFOLD WAY

complex	$X \not\cong X^*$	unitary	$U(n)$
real	$X \cong X^*$ $j^2 = 1$	orthogonal	$O(n)$
quaternionic	$X \cong X^*$ $j^2 = -1$	symplectic	$Sp(n)$

This accounts for the ‘classical groups’, leaving 5 exceptional groups related to the octonions and ‘triality’. These are important in superstring theory...

...but that’s a story for another day.

When it comes to duality, categories are just the beginning. We also have n -categories:



and so on...

An $(n + k)$ -category with only one j -morphism for $j < k$ is called a k -monoidal n -category. We have guesses about these...

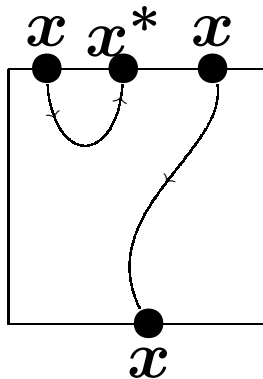
THE PERIODIC TABLE

	$n = 0$	$n = 1$	$n = 2$
$k = 0$	sets	categories	2-categories
$k = 1$	monoids	monoidal categories	monoidal 2-categories
$k = 2$	commutative monoids	braided monoidal categories	braided monoidal 2-categories
$k = 3$	“	symmetric monoidal categories	syllactic monoidal 2-categories
$k = 4$	“	“	symmetric monoidal 2-categories
$k = 5$	“	“	“
$k = 6$	“	“	“

Back in 1995, James Dolan and I proposed studying n -categories with *duals at all levels*, in order to make precise and prove:

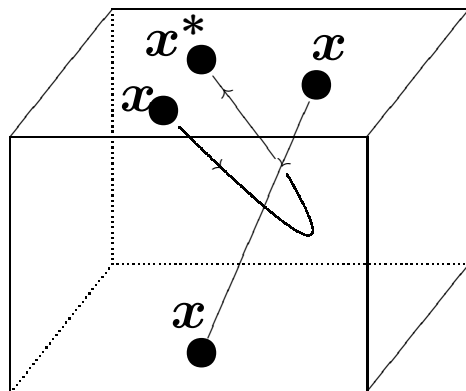
THE TANGLE HYPOTHESIS: The n -category of framed n -dimensional tangles in a $(n + k)$ -dimensional cube is the free k -monoidal n -category with duals on one object: the positively oriented point.

When $n = 1$, $k = 1$ we get a monoidal category with duals:



Morphisms here model worldlines of particles in 2d space-time.

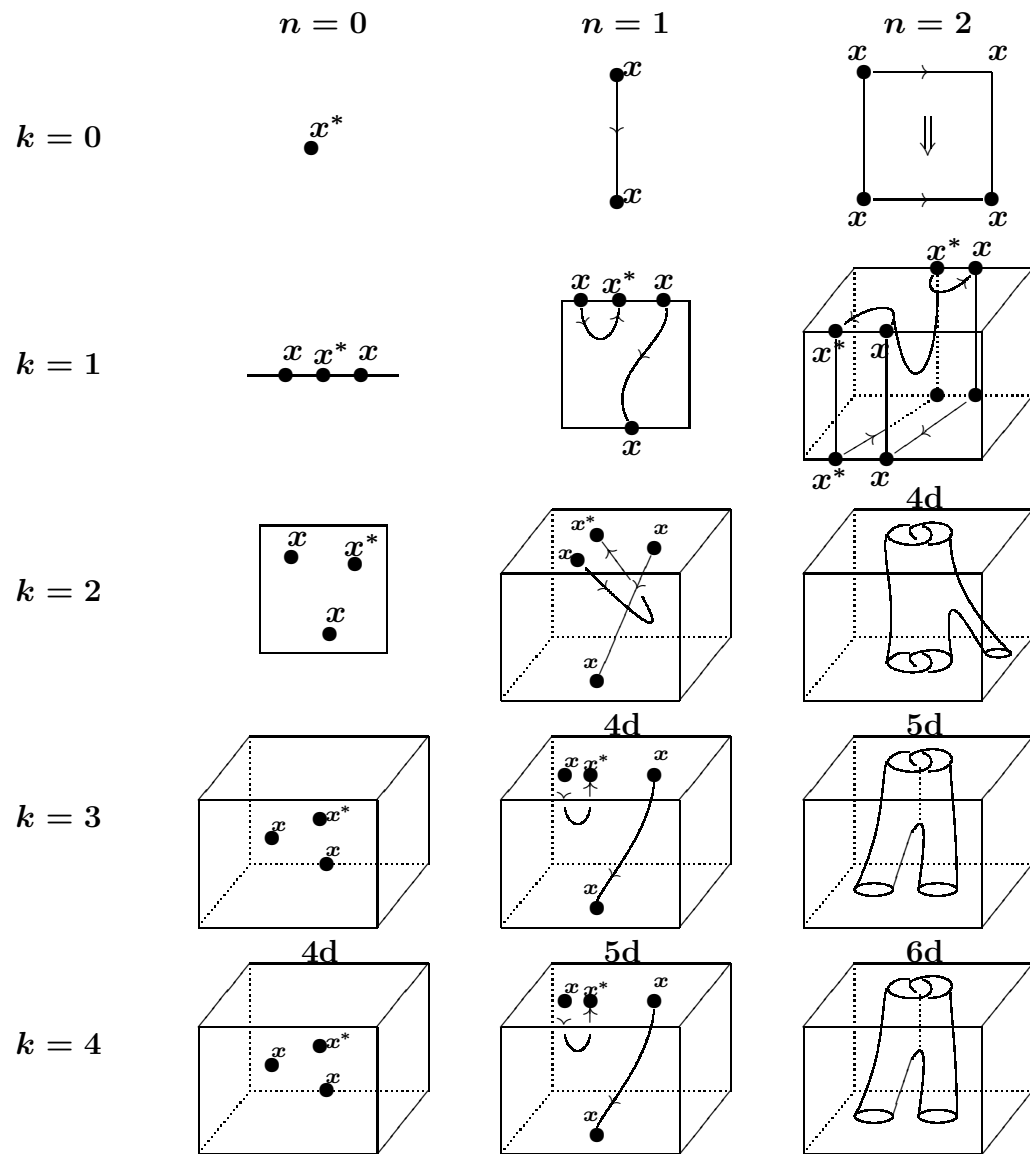
When $n = 1$, $k = 2$ we get a braided monoidal category with duals:



Morphisms here model worldlines of particles in 3d spacetime.

When $n = 1$, $k = 3$ we get a symmetric monoidal category with duals—i.e. a dagger-compact category. Morphisms here model models worldlines of particles in 4d spacetime. *Dagger-compact categories seem natural in logic and physics because we're used to particles in 4d spacetime!*

PERIODIC TABLE OF TANGLES



‘Categorified matrix mechanics’ gives examples of symmetric monoidal 2-categories with duals.

For example, there’s a symmetric monoidal 2-category \mathbf{Prof} where:

- an object is a small category \mathcal{C} ;
- a morphism T from \mathcal{C} to \mathcal{D} is a functor

$$T: \mathcal{C}^{\text{op}} \times \mathcal{D} \rightarrow \mathbf{Set}$$

also known as a *profunctor*;

- a 2-morphism is a natural transformation between profunctors.

\mathbf{Prof} has duals for objects: the dual of \mathcal{C} is \mathcal{C}^{op} . It doesn’t have duals (adjoints) for all morphisms—but a functor gives a profunctor with a right adjoint. It doesn’t have daggers for 2-morphisms.

We can do better by replacing Set here by a dagger-compact category.

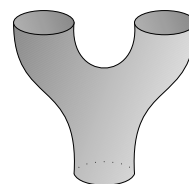
Furthermore, the symmetric monoidal 2-category of ‘finite-dimensional 2-Hilbert spaces’ has duals at all levels!

But let’s look at Prof....

A 2-dimensional topological quantum field theory gives a *Frobenius monoid* in $\mathbf{FinVect}$:



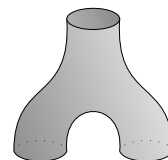
unit



multiplication

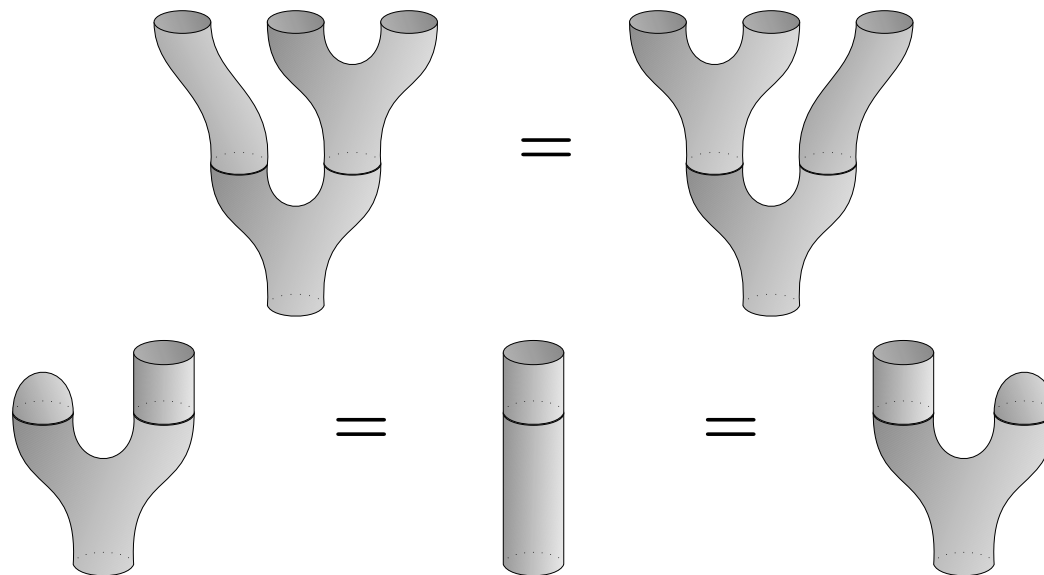


counit

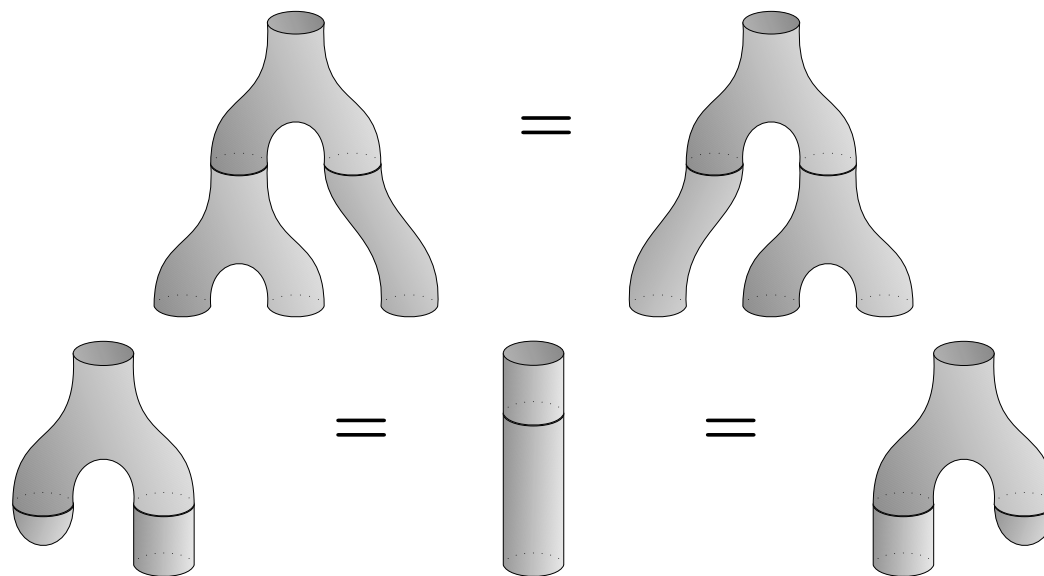


comultiplication

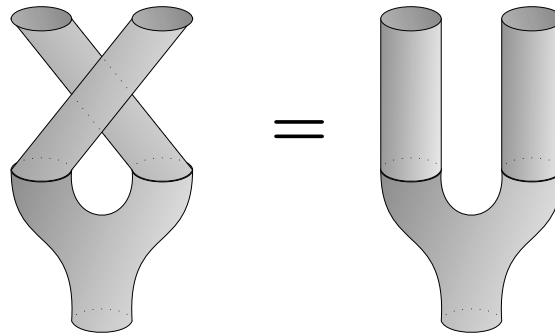
Associativity and unit laws:



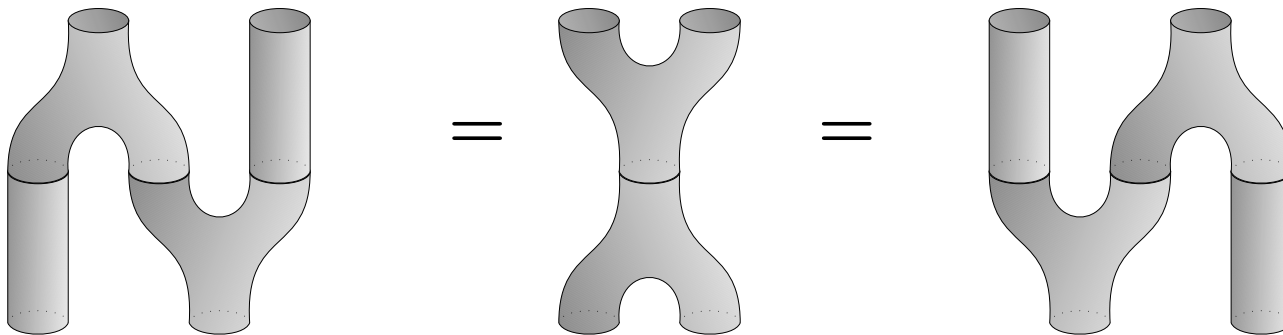
Coassociativity and counit laws:



Commutativity:



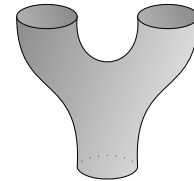
The Frobenius law:



Similarly, Ross Street showed that a $*$ -autonomous category \mathcal{C} gives a ‘weak Frobenius monoid’ in Prof!



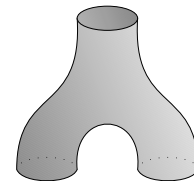
$\text{hom}(\text{TRUE}, A)$



$\text{hom}(A \& B, C)$



$\text{hom}(A, \text{FALSE})$



$\text{hom}(A, B \vee C)$

The Frobenius law is the ‘cut rule’ in logic!

SUMMARY

‘Matrix mechanics’ over commutative $*$ -rigs describes quantum *and classical* physics. See Aaron Fenyes’ no-cloning theorem for classical mechanics.

The study of duality unifies real, complex and quaternionic quantum mechanics into a single theory *which is already implicit* in standard physics.

Dagger-compact categories are the $n = 1$, $k = 3$ example of ‘ k -monoidal n -categories with duals’—the case most relevant to particles in 4d spacetime, but just one of many.

Treating Prof as a categorified version of FinVect relates *propositional logic*—i.e. $*$ -autonomous categories—to *categorified 2d topological quantum field theories!*