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The duality between set theory & logic

Idea:

Subset ^(of X) operations	Logic operations
\cup	\vee (or)
\cap	\wedge (and)
\emptyset	F (false)
X	T (true)
c (complement)	\neg (not)

e.g. $x \in S \cup T \iff x \in S \vee x \in T$

Last time we saw that given a compact Hausdorff space X , we get a commutative C^* -algebra

$$C(X) = \text{hom}_{\text{Top}}(X, \mathbb{C}) \quad (\text{continuous maps from } X \text{ to } \mathbb{C})$$

and, given a commutative C^* -algebra A , we get a compact Hausdorff space

$$\text{Spec}(A) = \text{hom}_{\text{Comm } C^* \text{ Alg}}(A, \mathbb{C}) \quad (\text{homomorphisms of } C^* \text{-algebras})$$

We call \mathbb{C} here a dualizing object: "the same object x in two different categories \mathcal{C} & \mathcal{D} " such that

$$\begin{aligned} \mathcal{C} &\longrightarrow \mathcal{D}^{\text{op}} \\ c &\longmapsto \text{hom}_{\mathcal{D}}(c, x) \in \mathcal{D} \end{aligned}$$

for some mysterious reason...

and

$$\begin{aligned} \mathcal{D}^{\text{op}} &\longrightarrow \mathcal{C} \\ d &\longmapsto \text{hom}_{\mathcal{C}}(d, x) \in \mathcal{C} \end{aligned}$$

are inverse functors, so \mathcal{C} & \mathcal{D}^{op} are equivalent.

Let's make the mystery work again!

We'll relate sets & Boolean algebras using the dualizing object
 $2 = \{0, 1\} \cong \{F, T\}$.

Boolean algebras are a little bit like C^* -algebras, but:

<u>Commutative C^*-algebras</u>	<u>Boolean algebras</u>
$+$	\vee
\cdot	\wedge
0	F
1	T
\mathbb{C}	$2 = \{F, T\}$

Given a set X , we'll make a Boolean algebra

$$2^X = \text{hom}_{\text{Set}}(X, 2).$$

An element here is a function $f: X \rightarrow \{0, 1\}$. Any subset $S \subseteq X$ gives a characteristic function:

$$\chi_S(x) = \begin{cases} 0 & x \notin S \\ 1 & x \in S \end{cases}$$

and conversely any such function gives a subset of X , so 2^X is just another way to think of the power set of X .

The operations \cup, \cap, \complement on subsets of X correspond to operations \vee, \wedge, \neg on functions $f: X \rightarrow \{0, 1\}$

$$\chi_{S \cup T} = \chi_S \vee \chi_T \quad \text{where} \quad \chi_S \vee \chi_T = \chi_S(x) \vee \chi_T(x) \quad \text{etc.}$$

$(2^S, \vee, \wedge, 0, 1)$ will be a Boolean algebra.

De Note: $S \subseteq T \Leftrightarrow \chi_S \leq \chi_T$
i.e. if $\chi_S(x) = 1$ then $\chi_T(x) = 1$.

Note: " \leq " is not a separate concept:

$$\begin{aligned}\chi_S \leq \chi_T &\Leftrightarrow \chi_S \wedge \chi_T = \chi_S \\ &\Leftrightarrow \chi_S \vee \chi_T = \chi_T\end{aligned}$$

Def: A partially ordered set (A, \leq) is called a lattice if every pair $a, b \in A$ has a least upper bound $a \vee b$ and greatest lower bound $a \wedge b$, and also a least element $0 = F$ and a greatest element $1 = T$.

Def: A distributive lattice is a lattice where \wedge & \vee distribute over each other.

Def: A Boolean algebra is a distributive lattice where every element $x \in A$ has a complement $\neg x$ such that

$$x \wedge \neg x = F$$

$$x \vee \neg x = T.$$

(If a complement exists, it is unique.)

Example: For any set S , 2^S is a Boolean algebra with pointwise defined \leq ; given $f, g \in 2^S$ we say $f \leq g$ if $f(x) \leq g(x) \forall x \in S$.

It thus has pointwise defined $\vee, \wedge, 0, 1, \neg$.

E.g. $(\neg f)(x) = \neg f(x)$

Not every Boolean algebra is isomorphic to one of this form!
 The Boolean algebras of the form 2^S are "complete atomic Boolean Algebras".

Def: A complete Boolean algebra A is one where every subset $S \subseteq A$ has a least upper bound $\bigvee_{x \in S} x$ and greatest lower bound $\bigwedge_{x \in S} x$, and these distribute over each other.

Def: An atom in a Boolean algebra A is an element $x \in A$ such that $x \neq 0$ and if $y < x$ then $y = 0$.

Example: in 2^S the atoms correspond to the elements of S , or singletons $\{s\} \subseteq S$.

Def: A Boolean algebra is atomic if $\forall x \in A, x = \bigvee_{x_2 \in A} y_2$ where $y_2 \in A$ are atoms.

There is a category CABA of complete atomic Boolean algebras and homomorphisms of complete Boolean algebras:

$\phi: A \rightarrow B$ preserving $\vee, \wedge, 0, 1, \neg$, \bigvee, \bigwedge
 enough to preserve these two

There is a category Set of sets & functions.

Thm Set is equivalent to $CABA^{op}$ via functors

$$\text{Set} \longrightarrow CABA^{op}$$

$$S \longmapsto 2^S = \text{hom}_{\text{Set}}(S, 2) \quad \text{and}$$

$$CABA^{op} \longrightarrow \text{Set}$$

$$A \longmapsto \text{hom}_{CABA}(A, 2), \quad \text{where } 2 = \{0, 1\} \text{ is a CABA}$$

in the obvious way.

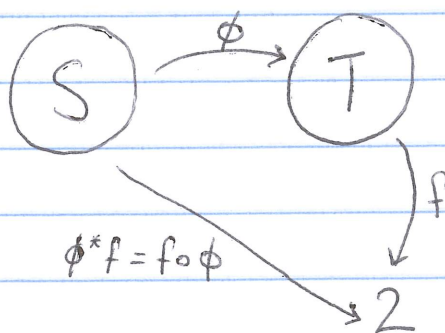
Given a function $\phi: S \rightarrow T$ we get a complete Boolean algebra homomorphism

$$\phi^*: 2^T \rightarrow 2^S \quad \text{by:}$$

$$\underbrace{\phi^*(f)}_{\in 2^S}(s) = f(\phi(s))$$

$$f \in 2^T \\ s \in S$$

This is called a pullback:



measurable subsets of \mathbb{R}

atomic but not complete Boolean algebra under pointwise operations

$$\{f \in L^\infty[0,1] : f(x) = 0 \text{ or } 1 \quad \forall x \in X\}$$

Complete but not atomic Boolean algebra under pointwise operations.