

5/14/02

We saw we can define a group object in any Cartesian category, i.e. any category w/ finite product (equiv. to a category w/ binary products and a terminal object).

We can also define a homomorphism $f: G \rightarrow G'$ between group objects:

• should preserve mult; inverses, identities

$$\begin{array}{ccc} \textcircled{1} & G \times G & \xrightarrow{m} & G \\ & \downarrow f \times f & \curvearrowright & \downarrow f \\ & G' \times G' & \xrightarrow{m'} & G' \end{array} \quad \text{preserves mult.}$$

$$\begin{array}{ccc} \textcircled{2} & 1 & \xrightarrow{i} & G \\ & \downarrow 1 & \curvearrowright & \downarrow f \\ & 1 & \xrightarrow{i'} & G' \end{array} \quad \text{preserve identity}$$

$$\begin{array}{ccc} \textcircled{3} & G & \xrightarrow{\text{inv}} & G \\ & \downarrow f & & \downarrow f \\ & G' & \xrightarrow{\text{inv}'} & G' \end{array} \quad \text{preserves inverses}$$

Secretly - these commutative diagrams involve a natural transformation.

Even more secretly - the group objects are functors.

Given a Cartesian category, we get a new category:

Defn: If C is a Cartesian category, let $C\text{-Grp}$ be the category whose objects are grp. objects in C and whose morphisms are homomorphisms.

A grp obj. in the cat. of sets is a group.

C	$C\text{-Grp}$
Set	Grp (groups)
Top	Top Grp (top groups mult is cont.)
Diff	Diff-Grp (Lie Groups)
(1) Grp	Ab-Grp (Grp-Grp) (abelian grps)
(2) Ab-Grp	(Grp-Grp-Grp) (Ab-Grp-Grp) Ab-Grp (abelian grps) (same as above)
(3) Vect products = \oplus term. obj = 0	Vect
Top-Grp	Ab-Top-Grp (abelian top. grps)
Diff-Grp	Ab-Diff-Grp (abelian Lie grp)
(4) Cat	Cat-Grp (categorical grps) (*)

(1) What's a group object in Grp ?

A group G w/ homomorphism: $m: G \times G \rightarrow G$
 (which is not the mult. in the group)
 (the original mult isn't a hano.)

unit map: $i: \mathbf{1} \rightarrow G$ "unit homomorphism"

↑ the 1-elt group (terminal grp)
 (think of it as picking out a particular elt of G)

Since i is a hano, must send identity elt of group
 maps

$1 \in \mathbf{1}$ to $1 \in G$. (no new data)

and "inv" hano: $\text{inv}: G \rightarrow G$

satisfying axioms of a group object.

m is a hano:

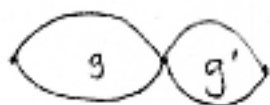
$$m(gg', hh') = m((g, g')(h, h')) = m(g, g')m(h, h'), \quad g, g', h, h' \in G$$

Let $m(g, g') = g \cdot g'$.

So we get:

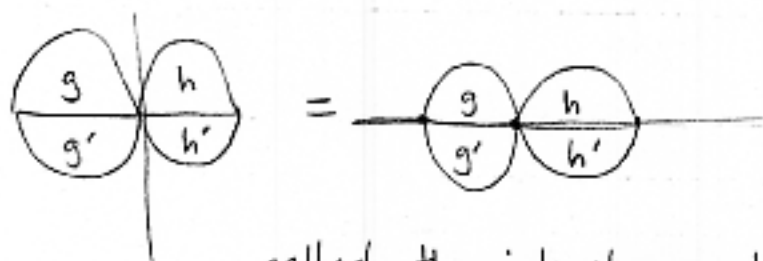
$$(gg') \cdot (hh') = (g \cdot g')(h \cdot h')$$

In a 2-category we have 2 forms of "composition".


 gg'

 $g \cdot g'$

so the law on the prev. pg looks like:

$$(gg') \cdot (hh') = (g \cdot g')(h \cdot h')$$



called the interchange law.

Guess: A group object in \mathbf{Grp} is a 2-cat
w/ only one object, and only one morphism
- i.e. some kind of commutative monoid.

These 2 multiplications are the same!

The identity 1 is the identity for the old \cdot , new
(horiz \cdot , vert) mult:

$$\begin{aligned}
 \text{g} \circ \text{h} &= \begin{array}{|c|c|} \hline \text{g} & 1 \\ \hline 1 & \text{h} \\ \hline \end{array} \\
 &= \begin{array}{|c|} \hline \text{g} \\ \hline \text{h} \\ \hline \end{array} \\
 &= \begin{array}{|c|c|} \hline 1 & \text{g} \\ \hline \text{h} & 1 \\ \hline \end{array} \\
 &= \text{h} \circ \text{g}
 \end{aligned}$$

$$\Rightarrow gh = g \cdot h = hg = h \cdot g$$

So - there is only one mult \cdot , it is commut., so we get a commut. grp.

(2) argument same as above; doesn't change if G above was abelian

(3) A v. space V w/ linear map $V \times V \rightarrow V$, etc...

same argument: new mult = add. in v. space

but we have no extra structure - nothing new,
and new mult = addition in v. space.

* A group object in Cat is a category object in Grp. 103

(4) grp objects in the category of categories

(*) categorical groups = strict 2-groups

= (strict 2-categories w/ one object and all morphisms a_i , 2-morphisms invertible)
(this defn is equiv to categorical grps)

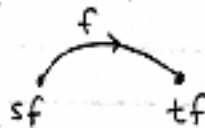
A grp is a special kind of category, so let's talk about:

Category Objects

what's a category object in some category C ?

It will have:

- an object $O \in C$ of objects
- an object $M \in C$ of morphisms
- source s , target morphisms $s, t: M \rightarrow O$



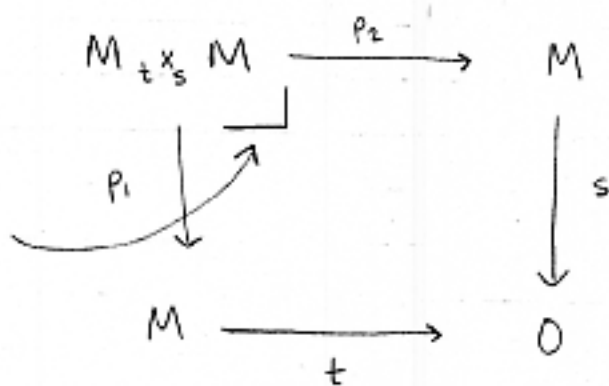
- identity $i: O \rightarrow M$
 $x \mapsto 1_x$ (identity-morphism) morphism

• composition morphism :

we can only compose composable morphisms

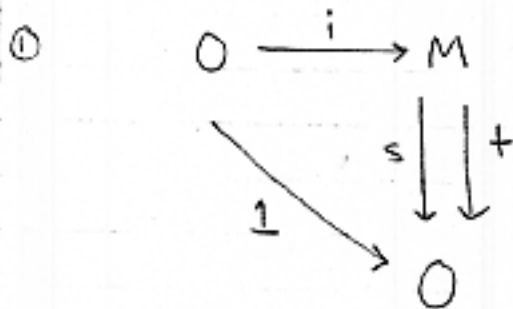
(pullback - subset of a product)
(pushout - quotient)

$\circ : M \times_t M \rightarrow M$ defined only on the pullback

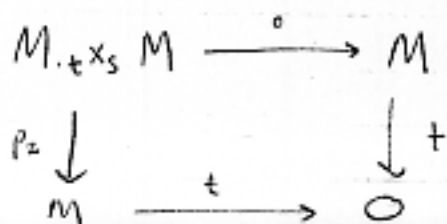


means
this is a
pullback
square.

such that the following commute: (defn of category)



② source / target of compositions must work out



$$\textcircled{2} \quad \begin{array}{ccc} 0_{1 \times s} M & \xrightarrow{i \times 1} & M_{t \times s} M \\ & \searrow P_2 & \downarrow \circ \\ & & M \end{array}$$

• composing on right by identity
 $f \cdot 1 = f$

$$\begin{array}{ccc} M_{t \times 1} 0 & \xrightarrow{1 \times i} & M_{t \times s} M \\ & \searrow P_1 & \downarrow \circ \\ & & M \end{array}$$

• composing on left by identity $1 \cdot f = f$

③ associativity

$$\begin{array}{ccc} M_{t \times s} M_{t \times s} M & \xrightarrow{1 \times 0} & M_{t \times s} M \\ \downarrow 0 \times 1 & & \downarrow \circ \\ M_{t \times s} M & \xrightarrow{\circ} & M \end{array}$$

* Moral: we can define "category object in C " if C has pullbacks.

Prop: A category has finite limits iff it has pullbacks and terminal objects.
(These give finite products.)

We can also define functors (functor-morphism) between category objects:

$F: (O, M) \longrightarrow (O', M')$ consists of morphism

$$F_o: O \longrightarrow O'$$

$$F_m: M \longrightarrow M'$$

st the usual properties of a functor hold.

Defn: Given a category w/ pullbacks C , we let $C\text{-Cat}$ be the category where objects are category objects in C and morphisms are functors between those.

Note: often people use "topological category" when \mathcal{O} is discrete.

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<u>C</u>	<u>C-Cat</u>
Set	Cat
(1) Top	Top-Cat \approx "topological category" (top. space of objects, space of morphisms, comp. is cart.)
(2) Diff	Diff-Cat \approx "smooth category" (Ehresman: "differentiable category")

(1) an example of a topological category is the fundamental groupoid.



Here $\mathcal{O} = X$

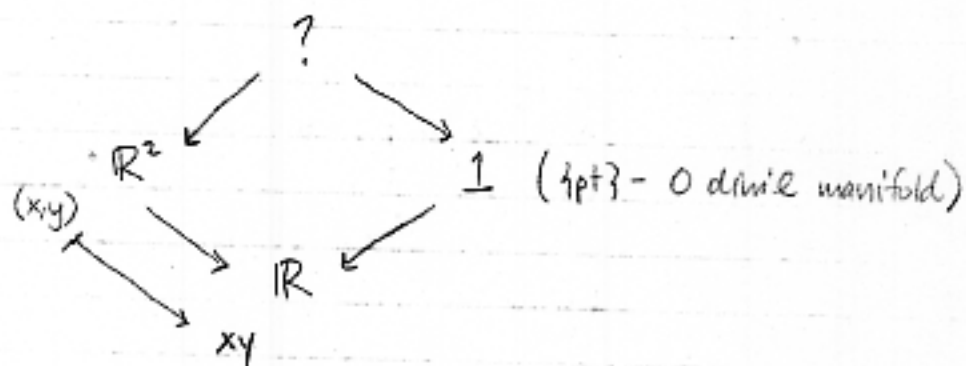
or - $\pi_1^{\text{thin}}(M)$

(2) But Diff doesn't always have pullbacks.

$$\begin{array}{ccc} & ? & A \\ & & \downarrow \\ B & \longrightarrow & X \end{array}$$

But we just need the pullbacks to exist that we're going to use.

Ex)

Hope $? \subseteq \mathbb{R}^2 \times 1$

$$? = \{ (x,y) \mid xy=0 \} \subseteq \mathbb{R}^2 \times 1$$

$$\{xy=0\} \in \mathbb{R}^2$$



not a smooth submanifold

o

(...)

...