

Structure Types

We've defined a structure type to be a functor

$$F: \text{FinSet}_0 \rightarrow \text{Set}$$

We think of F as assigning to any finite set S the "set of all F -structures on S ", F_S .

Given a structure type F , we can define its generating function $|F| \in \mathbb{C}[[z]]$, by

$$|F|(z) = \sum_{n \geq 0} \frac{|F_n|}{n!} z^n$$

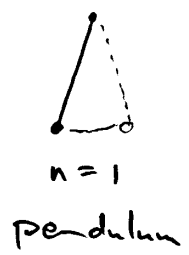
where F_n is the set of F -strs. on the n -elt set, " n ", & $|F_n|$ is its cardinality.

We saw that these give new insight into the quantum harmonic oscillator. This quarter, we'll

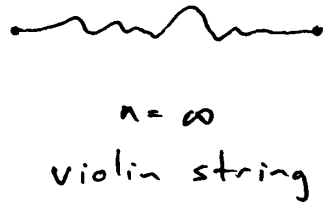
- 1) Generalize to formal power series with many variables (possibly infinitely many) which are generating fns of " n -sorted str. types" $F: \text{FinSet}_0^n \rightarrow \text{Set}$.

This lets us categorify the quantum harmonic oscillator with n degrees of freedom, especially $n = \infty$.

Quantum field theory is about oscillators w. infinitely many degrees of freedom:



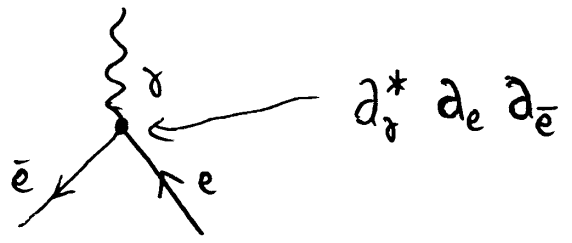
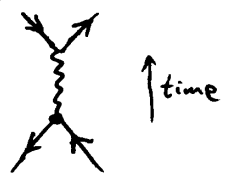
Config space is \mathbb{R}



config space is infinite-dimensional, " \mathbb{R}^∞ "

Here creation and annihilation operators create "quanta" of different sorts... whereas last quarter, in the $n=1$ case, we only had one creation operator & one annihilation operator. We'll be able to

categorify the theory of Feynman diagrams:



2) Understand more deeply how structure types are "categorified formal power series." We've seen $\text{hom}(\text{FinSet}_0, \text{Set})$ is a category analogous to the set $\mathbb{C}[[z]]$. We need to make this analogy precise.

To do this, we'll need to generalize structure types to "stuff types." We'll learn about: properties, structure, 'n' stuff! We need this to categorify Feynman diagrams

3) Have Fun: use structure types to do all sorts of combinatorics, number theory, ...

Example of Fun:

Recall:

1) If F is "being a totally ordered n -elt. set",

$$|F|(z) = z^n$$

since $|F_n| = n!$ & $|F_m| = 0$ if $m \neq n$.

So we call F " Z^n ":

$$|Z^n| = z^n.$$

2) If F & G are str. types, to put an $F+G$ -str on S we: choose either F or G & then put such a structure on S .

$$|F+G| = |F| + |G|$$

3) To put an FG-str. on S we chop S into 2 subsets & put an F str. on first part, a G-str. on second part.

$$|FG| = |F||G|.$$

Now: let an F-str. on S be a way of totally ordering it & then chopping it up into "blocks" of size 1 or 2:

$$7 \quad 1 \mid 2 \mid 4 \quad 3 \mid 5 \mid 6 \quad \text{is an F-str. on } \{1, \dots, 7\}$$

A subset $T \subseteq S$ is a block if $x, y \in T$, $x \leq z \leq y \Rightarrow z \in T$

$$|F_0| = 1$$

$$|F_1| = 1$$

$$|F_2| = 4$$

$$\begin{array}{cc} xy & yx \\ x|y & y|x \end{array}$$

Note:

$$F \cong \mathbb{Z}^2 F + \mathbb{Z} F + 1$$

i.e. to put an F str. on S , either chop it in two & put the structure of being a totally ordered 2-elt set ~~on the first part~~ or being a totally ordered 1-elt set on the first part, and then put an F str. on the second part. (or $F = \emptyset$)

So:

$$|F| = z^2 |F| + z |F| + 1$$

$$(1 - z - z^2) |F| = 1$$

$$|F|(z) = \frac{1}{1 - z - z^2}$$

Note:

$$|F|(z) = 1 + (z + z^2) + (z + z^2)^2 + (z + z^2)^3 + \dots$$

$$= 1 +$$

$$z + z^2 +$$

$$z^2 + 2z^3 + z^4 +$$

$$z^3 + 3z^4 + 3z^5 + z^6 +$$

$$z^4 + 4z^5 + 6z^6 + 4z^7 + z^8 + \dots$$

$$= 1 + 1z + 2z^2 + 3z^3 + 5z^4 + \dots$$

$$= 1 + z + \frac{4z^2}{2!} + \frac{18z^3}{3!} + \dots$$

Note: $|F_n| = f_n \cdot n!$, where $n!$ comes from the choice of orderings & f_n comes from choice of block structure:

$$f_4 = 5 \quad \begin{array}{l} 1 \cdot | \cdot | \cdot | \cdot | \cdot | \\ 1 \cdot | \cdot \cdot | \cdot | \cdot | \\ 1 \cdot | \cdot \cdot \cdot | \cdot | \cdot | \end{array} \quad \begin{array}{l} | \cdot \cdot | \cdot | \cdot | \\ | \cdot | \cdot | \cdot \cdot | \\ | \cdot \cdot \cdot | \cdot | \end{array}$$

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Since $|F| = \sum_{n \geq 0} f_n z^n$,

$$|F| = z^2 |F| + z |F| + 1$$

says $\sum_{n \geq 0} f_n z^n = \sum_{n \geq 0} f_n z^{n+1} + \sum_{n \geq 0} f_n z^{n+2} + 1$

$$\sum_{n \geq 0} f_n z^n = \sum_{n \geq 1} f_{n-1} z^n + \sum_{n \geq 2} f_{n-2} z^n + 1$$

so

$$f_n = f_{n-1} + f_{n-2} \quad \text{for } n \geq 2$$

& we know $f_0 = f_1 = 1$, so f_n is the n th Fibonacci number!

8 Jan 2004

(Lots of people gone to $\frac{\text{MAA}}{\text{AMS}}$ meeting today...)

We saw that if F is the structure type "total ordering & partition into blocks of length 1 or 2, then

$$F \cong \mathbb{Z}^2 F + \mathbb{Z} F + 1$$

$$\text{so } |F| = \frac{1}{1-z-z^2}$$

We can study asymptotics of $|F_n| = n! f_n$ by studying radius of convergence, poles, etc. of $|F|(z)$. This is quite general, so let Φ

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be any str. type & $|\Phi|(z) = \sum_{n \geq 0} \frac{|\Phi_n|}{n!} z^n = \sum_{n \geq 0} a_n z^n$

(where $a_n = f_n$ in our example). Hadamard's

Theorem says that if R is the radius of

convergence of $\sum_{n \geq 0} a_n z^n$ then $\frac{1}{R} = \limsup \sqrt[n]{|a_n|}$

This implies that for any $\varepsilon > 0$, eventually

$$\sqrt[n]{|a_n|} \leq \frac{1}{R} + \varepsilon$$

or

$$|a_n| \leq \left(\frac{1}{R} + \varepsilon\right)^n$$

We don't get any good lower bounds unless $\lim \sqrt[n]{|a_n|}$ exists, in which case $\forall \varepsilon > 0$ we eventually have

$$|a_n| \geq \left(\frac{1}{R} - \varepsilon\right)^n$$

In our example:

$$\frac{1}{1-z-z^2}$$

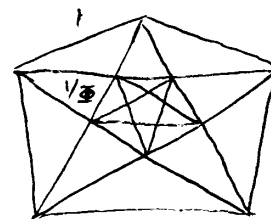
has simple poles at roots of $z^2 + z - 1 = 0$

$$z = \frac{-1 \pm \sqrt{1+4}}{2} = \frac{-1 \pm \sqrt{5}}{2}$$

Golden ratio $\frac{\sqrt{5}+1}{2} = \Phi = 1.6180339\dots$

$$\frac{\sqrt{5}-1}{2} = \Phi^{-1} = 0.6180339\dots$$

$$\Phi = 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \dots}}}}}$$



$$= \sqrt{1 + \sqrt{1 + \sqrt{1 + \sqrt{1 + \sqrt{1 + \dots}}}}}$$

So: the radius of convergence of $\sum_{n \geq 0} f_n z^n$ is

$$\left(\text{distance of closest pole from origin} \right) \text{ of } \frac{1}{1-z-z^2} = \frac{1}{\Phi}$$

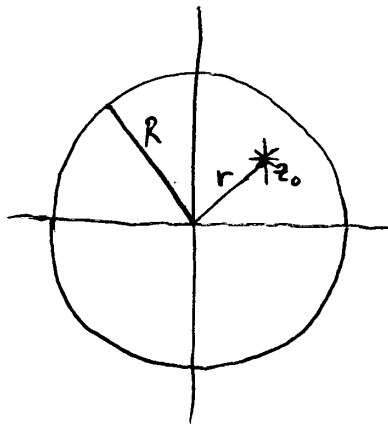
So, Hadamard's Theorem says

$$|f_n| \leq (\Phi + \epsilon)^n$$

for large enough n .

Supped-up Hadamard's Theorem:

Suppose $f(z) = \sum a_n z^n$ is analytic in a disc of radius R about the origin except for a simple pole at distance $r < R$ from the origin:



Then

$$|a_n| \sim c \left(\frac{1}{r}\right)^n \quad c > 0$$

↖ "asymptotic"

i.e.

$$\lim_{n \rightarrow \infty} \frac{|a_n|}{c \left(\frac{1}{r}\right)^n} = 1$$

Proof: Can find $k \neq 0$ s.t.

$$f(z) = \tilde{f}(z) + \frac{k}{z-z_0}$$

where f is analytic on the disc of radius R .

Let

$$\tilde{f}(z) = \sum \tilde{a}_n z^n$$

so

$$\sum a_n z^n = \sum \tilde{a}_n z^n - \frac{k}{z_0} \sum \left(\frac{z}{z_0}\right)^n$$

$$\begin{aligned} \frac{k}{z-z_0} &= \frac{-k/z_0}{1-z/z_0} \\ &= -\frac{k}{z_0} \sum \left(\frac{z}{z_0}\right)^n \end{aligned}$$

By Hadamard's Thm,

$$|\tilde{a}_n| \leq \left(\frac{1}{R} + \varepsilon\right)^n, \text{ eventually,}$$

but the other term has:

$$\left| -\frac{k}{z_0} \left(\frac{1}{z_0}\right)^n \right| = c \left(\frac{1}{r}\right)^n$$

Choose ε s.t. $\frac{1}{r} > \frac{1}{R} + \varepsilon$ and conclude $|\tilde{a}_n|$ is dwarfed by $c \left(\frac{1}{r}\right)^n$ for large n , so $|a_n| \sim c \left(\frac{1}{r}\right)^n$. \blacksquare

So in our example we actually find

$$f_n \sim c \Phi^n.$$

We can do much better, though.

$$\begin{aligned} \frac{1}{1-z-z^2} &= \frac{A}{z+\Phi} + \frac{B}{z-\Phi^{-1}} \\ &= \frac{A/\Phi}{1+z/\Phi} - \frac{B\Phi}{1-z\Phi} \\ &= \frac{A}{\Phi} \sum \left(\frac{-z}{\Phi}\right)^n - B\Phi \sum (z\Phi)^n = \sum f_n z^n \end{aligned}$$

~~In fact:~~

So:

$$f_n = \frac{A}{\Phi} \left(-\frac{1}{\Phi}\right)^n - B\Phi (\Phi)^n$$

In fact

$$f_n = \frac{\Phi^{n+1} - (-\Phi)^{-(n+1)}}{\sqrt{5}}$$

Since $\Phi^{-(n+1)} \rightarrow 0$ as $n \rightarrow \infty$, this means that

as $n \rightarrow +\infty$, $\left|f_n - \frac{\Phi^{n+1}}{\sqrt{5}}\right| \rightarrow 0$. I.e.

$\frac{\Phi^n}{\sqrt{5}}$ gets closer and closer to an integer! By

the way, $\exists a > 0$ s.t. $a \notin \mathbb{Q}$ but a^n gets closer & closer to integers as $n \rightarrow \infty$.

In fact:

$$\frac{\Phi^1}{\sqrt{5}} = .73\dots \text{ has } 1 = f_0 \text{ as its closest integer}$$

$$\frac{\Phi^2}{\sqrt{5}} = 1.13\dots \text{ has } 1 = f_1 \text{ " "}$$

$$\frac{\Phi^3}{\sqrt{5}} = 1.89\dots \text{ has } 2 = f_2 \text{ " "}$$

Therefore

$$f_n = \left\lfloor \frac{\Phi^n}{\sqrt{5}} \right\rfloor \leftarrow \text{closest integer}$$

Moral: The growth rate of $|\Phi_n|$ for some str. type Φ is controlled by singularities (poles, branch cuts, ...) of $|\Phi|(z)$. If $|\Phi|(z)$ is a rational function, we can use partial fractions to completely solve for $|\Phi_n|$ in "closed form."

So: there's a classification of str. types according to their generating fns. - e.g. "rational structure types" have rational generating functions: this gives a ^{full} subcategory of the category of str. types, $\text{hom}(\text{finSet}_0, \text{Set})$, and this subcategory is closed under $+$ & \cdot .