

Reeb Graph Smoothing Via Cosheaves

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Why category theory?

- It is a convenient language for describing **persistence modules**.
- It gives clues to finding the 'right' definitions and concepts.
- It gives immediate access to deeper theorems.
- We are free to drop it when it doesn't fit.

Category Theory for Applied Topology

Preordered sets

Let P be a set with a reflexive transitive relation \leq . Then

- objects = { elements of P }
- morphisms = { relations $x \leq y$ }

defines a category \mathbf{P} .

Directed graphs

A directed graph defines a category:



or



or



(Identities and composites are implicit.)

Sublevelset persistent homology

Let $f : X \rightarrow \mathbf{R}$. Consider the category \mathbf{n} defined by

$$0 \longrightarrow 1 \longrightarrow \cdots \longrightarrow n-1,$$

and select $a_0 \leq a_1 \leq \cdots \leq a_{n-1}$. From

$$X^{a_0} \xrightarrow{\subseteq} X^{a_1} \xrightarrow{\subseteq} \cdots \xrightarrow{\subseteq} X^{a_{n-1}},$$

construct

$$H(X^{a_0}) \longrightarrow H(X^{a_1}) \longrightarrow \cdots \longrightarrow H(X^{a_{n-1}}).$$

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Definitions

- $X^t := f^{-1}(-\infty, t]$ sublevelset
- $X_s := f^{-1}[s, +\infty)$ superlevelset
- $X_s^t := f^{-1}[s, t]$ interlevelset

Sublevelset persistent homology

The 'persistence module'

$$H(X^{a_0}) \longrightarrow H(X^{a_1}) \longrightarrow \cdots \longrightarrow H(X^{a_{n-1}})$$

can be thought of as a **functor**

$$\mathbf{n} \xrightarrow{F} \mathbf{Top} \xrightarrow{H} \mathbf{Vect.}$$

This means:

- For each object of \mathbf{n} we have a vector space.
- For each morphism of \mathbf{n} we have a linear map.

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Generalized persistence modules (Bubenik, Scott 2014)

A 'generalized persistence module' is simply a functor $\mathbb{V} : \mathbf{C} \rightarrow \mathbf{D}$.

- Usually \mathbf{C} is a pre-ordered set, such as $\mathbf{n}, \mathbf{N}, \mathbf{Z}, \mathbf{R}$.
- Usually \mathbf{D} is an abelian category, such as $\mathbf{Vect}, \mathbf{Ab}$.

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Categories of functors

The collection of functors $\mathbf{C} \rightarrow \mathbf{D}$ is itself a category, denoted $\mathbf{D}^{\mathbf{C}}$. The morphisms are **natural transformations** $\phi : \mathbb{V} \Rightarrow \mathbb{W}$, defined by the following data:

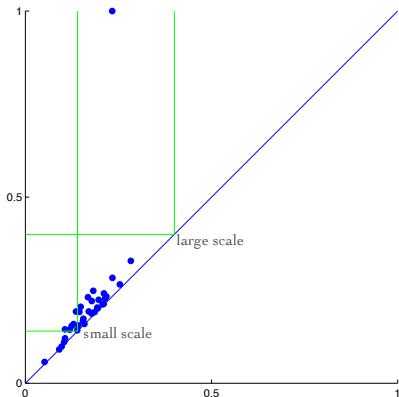
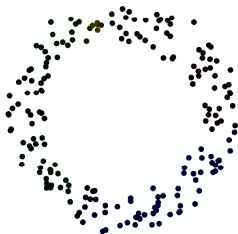
- For every $c \in \mathbf{C}$ there is a map $\phi_c : V_c \rightarrow W_c$.
- For every map $f : c \rightarrow c'$ in \mathbf{C} the diagram

$$\begin{array}{ccc} V_c & \xrightarrow{\phi_c} & W_c \\ V[f] \downarrow & & \downarrow W[f] \\ V_{c'} & \xrightarrow{\phi_{c'}} & W_{c'} \end{array}$$

is required to commute.

Story 1: Persistence diagrams

Persistent homology takes a filtered space $\mathbb{X} = \{X_t \mid t \in \mathbf{R}\}$ and returns a **barcode** of intervals $[p, q) \subset \mathbf{R}$ or a **persistence diagram** of points $(p, q) \in \mathbf{R}^2$.



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How is this defined?

Algorithmic approach (Edelsbrunner, Letscher, Zomorodian 2000).

- Discretize the t -variable.
- Present \mathbb{X} as a finite list of cells, attached in sequence.
- Some cells σ generate new homology cycles.
- Other cells τ destroy cycles created by an earlier σ .
- There is an interval $[t_\sigma, t_\tau)$ for each such pair (σ, τ) .
- There is an interval $[t_\sigma, +\infty)$ for each σ whose cycle is never destroyed.

Story 1: Persistence diagrams

Using commutative algebra (Zomorodian, Carlsson 2003).

- Discretize the t -variable to integers: $t = 0, 1, 2, \dots$
- Present \mathbb{X} as an increasing sequence:

$$\mathbb{X}: \quad X_0 \subset X_1 \subset X_2 \subset \dots$$

- Apply a homology functor $H = H(-; \mathbf{k})$ to the sequence:

$$H(\mathbb{X}): \quad H(X_0) \rightarrow H(X_1) \rightarrow H(X_2) \rightarrow \dots$$

- Observe that $H(\mathbb{X})$ is a graded module over the polynomial ring $\mathbf{k}[z]$, where z acts by shifting to the right.
- Decompose this graded module as a direct sum of cyclic submodules.
- Summands $z^s \mathbf{k}[z]/(z^{t-s})$ are recorded as intervals $[s, t)$.
- Summands $z^s \mathbf{k}[z]$ are recorded as intervals $[s, +\infty)$.

Story 1: Persistence diagrams

Using quiver theory (Carlsson, dS 2010).

- Discretize the t -variable to integers: $t = 0, 1, \dots, n - 1$.
- Present \mathbb{X} as a sequence of spaces with maps:

$$\mathbb{X}: X_0 \rightarrow X_1 \rightarrow \dots \rightarrow X_{n-1}$$

- Apply a homology functor $H = H(-; \mathbf{k})$ to the sequence:

$$H(\mathbb{X}): H(X_0) \rightarrow H(X_1) \rightarrow \dots \rightarrow H(X_{n-1})$$

- Observe that $H(\mathbb{X})$ is a representation of the quiver $\bullet \rightarrow \bullet \rightarrow \dots \rightarrow \bullet$.
- Decompose $H(\mathbb{X})$ as a direct sum of indecomposable representations.
- According to Gabriel (1970), the indecomposables are precisely the **intervals**:

$$0 \rightarrow \dots \rightarrow 0 \rightarrow \mathbf{k} \rightarrow \dots \rightarrow \mathbf{k} \rightarrow 0 \rightarrow \dots \rightarrow 0$$

The list of summands of $H(\mathbb{X})$ gives the persistence intervals.

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When the arrows have mixed orientations \leftarrow, \rightarrow , we get **zigzag persistence**.

Story 1: Persistence diagrams

What if we wish to work with a continuous parameter?

Interval decomposition

- Let \mathbb{V} be a persistence module defined over the real numbers \mathbf{R} .
- Suppose

$$\mathbb{V} = \bigoplus_{k \in K} \mathbb{I}_{[a_k, b_k]}$$

where $\mathbb{I} = \mathbb{I}_{[a, b]}$ denotes the persistence module with

$$I_t = \begin{cases} \mathbf{k} & \text{if } t \in [a, b] \\ 0 & \text{otherwise} \end{cases}$$

and all maps i_s^t having full rank. (Open, half-open intervals allowed too.)

- Then we can define the persistence diagram to be

$$\text{Dgm}(\mathbb{V}) = \{(a_k, b_k) \mid k \in K\},$$

a multiset of points in the half-plane above the diagonal.

Story 1: Persistence diagrams

Problem

Not every \mathbb{V} decomposes into intervals.

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Theorem (Gabriel, Auslander, Ringel–Tachikawa, Webb, Crawley-Boevey)

Let \mathbb{V} be a persistence module over $\mathbf{T} \subseteq \mathbf{R}$. In either of the following situations, \mathbb{V} decomposes into interval modules:

- \mathbf{T} is a finite set; or
- Every V_t is finite-dimensional.

On the other hand, there exists a persistence module over \mathbf{Z} which does not admit an interval decomposition.

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Solution (Chazal, dS, Glisse, Oudot 2016)

Define a measure which counts the number of persistence points in an arbitrary rectangle. Infer the existence of the persistence diagram. This works if the maps $V_s \rightarrow V_t$ are finite-rank whenever $s < t$.

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Definition 1 (non-functorial)

Let

$$\mu([a, b] \times [c, d]) = r_b^c - r_a^c - r_b^d + r_a^d$$

where $r_s^t = \text{rank}(V_s \rightarrow V_t)$.

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$$\mu([a, b] \times [c, d]) = \dim(M_{a,b,c,d}\mathbb{V})$$

where

$$M_{a,b,c,d}\mathbb{V} = \frac{[\text{Im}(v_b^c) \cap \text{Ker}(v_c^d)]}{[\text{Im}(v_a^c) \cap \text{Ker}(v_c^d)]}.$$

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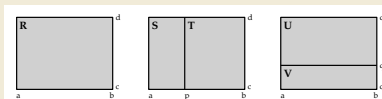
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Note. Each $M_{a,b,c,d}$ extends to a functor $\mathbf{Vect}^{\mathbb{R}} \rightarrow \mathbf{Vect}$.

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Solution step

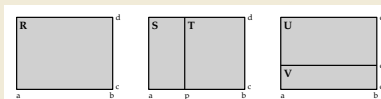
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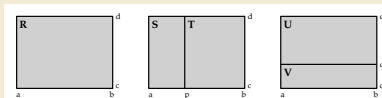
Proof 1 (for horizontal split)

$$r_b^c - r_a^c - r_b^d + r_a^d = (r_p^c - r_a^c - r_p^d + r_a^d) + (r_b^c - r_p^c - r_b^d + r_p^d)$$

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Proof 2 (for horizontal split)

There is a short exact sequence

$$0 \rightarrow \left[\frac{\text{Im}(v_p^c) \cap \text{Ker}(v_c^d)}{\text{Im}(v_a^c) \cap \text{Ker}(v_c^d)} \right] \rightarrow \left[\frac{\text{Im}(v_b^c) \cap \text{Ker}(v_c^d)}{\text{Im}(v_a^c) \cap \text{Ker}(v_c^d)} \right] \rightarrow \left[\frac{\text{Im}(v_b^c) \cap \text{Ker}(v_c^d)}{\text{Im}(v_p^c) \cap \text{Ker}(v_c^d)} \right] \rightarrow 0$$

or, in other words, a short exact sequence of functors

$$0 \longrightarrow M_{a,p,c,d} \longrightarrow M_{a,b,c,d} \longrightarrow M_{p,b,c,d} \longrightarrow 0$$

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Question (of Morozov)

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Answer 1: Constructing a functorial persistence diagram

Let $\mathbb{V} : \mathbf{R} \rightarrow \mathbf{Vect}$ be a persistence module. Select

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The **functorial persistence diagram** with respect to (a_n) is the function

$$(m, n) \mapsto M_{a_m, a_{m+1}, a_n, a_{n+1}} \mathbb{V}$$

for integers $m < n$. At each point there is a vector space.

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Pros and cons

- A map $\mathbb{V} \rightarrow \mathbb{W}$ between persistence modules induces a map between f.p.d.
- This method defines a persistence diagram in **any abelian category**.
- It is not so easy to change the discretization.

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- It is not so easy to change the discretization.
- **What is the right metric between these diagrams?**

Story 2: Interleaving

Stability theorem (Cohen-Steiner, Edelsbrunner, Harer 2007)

The map $\{\text{persistence modules}\} \rightarrow \{\text{diagrams}\}$ is 1-Lipschitz.

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Relators

The metrics on the two spaces are defined in terms of 'relators'.

- Two persistence modules may be related by an **interleaving**.
- Two diagrams may be related by a **matching**.

Every relator, of each type, has a size associated with it. The metrics are defined by finding the infimum of the size of relators between a given pair of objects.

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Stability theorem (Cohen-Steiner, Edelsbrunner, Harer 2007)

If two persistence modules admit an ϵ -interleaving, then their persistence diagrams admit an ϵ -matching.

Story 2: Interleaving

Definition

Let \mathbb{V}, \mathbb{W} be persistence modules. An ϵ -interleaving between \mathbb{V}, \mathbb{W} is a pair (Φ, Ψ) where $\Phi = (\phi_t)$ and $\Psi = (\psi_t)$ are collections of maps

$$\phi_t : V_t \rightarrow W_{t+\epsilon}$$

$$\psi_t : W_t \rightarrow V_{t+\epsilon}$$

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Glisse's Lemma (Chazal, Cohen-Steiner, Glisse, Guibas, Oudot 2009)

The proof of the stability theorem relies on the following fact. If \mathbb{V}, \mathbb{W} are ϵ -interleaved, then there is a 1-parameter family

$$(\mathbb{V}_s \mid s \in [0, \epsilon])$$

with $\mathbb{V}_0 = \mathbb{V}$ and $\mathbb{V}_\epsilon = \mathbb{W}$, and where $\mathbb{V}_r, \mathbb{V}_s$ are $|r - s|$ -interleaved for all r, s .

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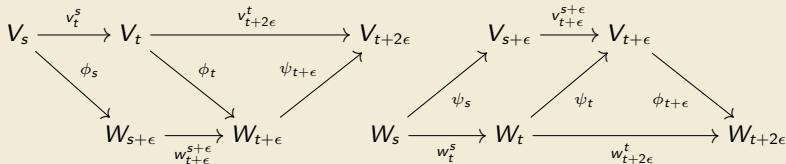
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The [various conditions] amount to the assertion that there is a unique way to get from any of the V_t, W_t to any other. All compositions of the $v_s^t, w_s^t, \phi_t, \psi_t$ with the same start and end point must agree.

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Interleavor categories (Chazal, dS, Glisse, Oudot 2016)

An ϵ -interleaved pair of modules $(\mathbb{V}, \mathbb{W}, \Phi, \Psi)$ is 'the same thing' as a persistence module defined over the set $\mathbf{I} = \mathbf{R} \times \{0, \epsilon\}$ (two copies of the real line) with the partial order

$$(s, a) \leq (t, b) \Leftrightarrow \begin{cases} s \leq t & \text{if } a = b \\ s + \epsilon \leq t & \text{if } a \neq b \end{cases}$$

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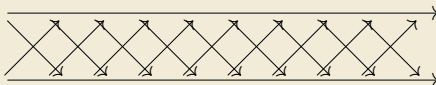
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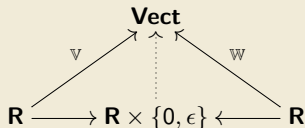
$\mathbf{R} \times \{0, \epsilon\}$:



Story 2: Interleaving

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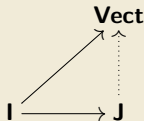
Proof of Glisse's Lemma

Consider the set $\mathbf{J} = \mathbf{R} \times [0, \epsilon]$ with the partial order

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This contains the interleavor category \mathbf{I} as a sub-poset. An ϵ -interleaving between two persistence modules corresponds to a functor $\mathbf{I} \rightarrow \mathbf{Vect}$ which restricts to \mathbb{V}, \mathbb{W} on the two respective copies of the real line.

An interpolation (\mathbb{V}_t) is found constructing an extension of the functor to \mathbf{J} :



Since \mathbf{I} is a full subcategory of \mathbf{J} , and \mathbf{Vect} contains all limits and colimits, the problem is solved by taking a left or right Kan extension.

Story 2: Interleaving

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Story 2: Interleaving

Question (of Morozov)

Is the persistence diagram functorial?

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Is the persistence diagram functorial?

Answer 2

The persistence diagram is a map

$$\{\text{persistence modules}\} \rightarrow \{\text{diagrams in the upper half-plane}\}$$

What are the morphisms that make these into categories?

- A morphism $\mathbb{V}_1 \rightarrow \mathbb{V}_2$ could be an interleaving pair (ϕ, ψ) .
- A morphism $\text{Dgm}_1 \rightarrow \text{Dgm}_2$ could be a matching between points.

For both notions there is an associative composition law with identities.

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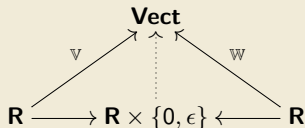
Answer 2⁺ (Bauer, Lesnick 2015)

Almost. See recent work of Ulrich Bauer and Michael Lesnick.

Story 3: Interleaving Metrics

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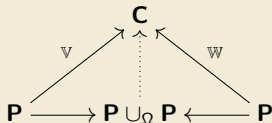
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Story 3: Interleaving Metrics

Interleavings for generalized persistence modules over a poset

Two persistence modules $\mathbb{V}, \mathbb{W} : \mathbf{P} \rightarrow \mathbf{C}$ are Ω -interleaved iff the following functor extension problem has a solution:



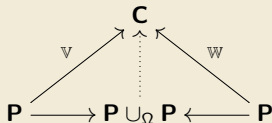
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where $\Omega : \mathbf{P} \rightarrow \mathbf{P}$ is a **translation**.

Story 3: Interleaving Metrics

Translations (Bubenik, dS, Scott 2015)

Trans_P is the poset of functions $\Omega : \mathbf{P} \rightarrow \mathbf{P}$ that are order-preserving and satisfy $x \leq \Omega x$ for all $x \in \mathbf{P}$.

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A **superlinear family** is a 1-parameter family of translations of \mathbf{P}

$$(\Omega_\epsilon \mid \epsilon \in [0, \infty))$$

such that

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Examples of superlinear families

- $\mathbf{P} = \mathbf{R}$,
 $\Omega_\epsilon(t) = t + \epsilon$.
- $\mathbf{P} = \{\text{compact intervals in the real line}\}$,
 $\Omega_\epsilon([a, b]) = [a - \epsilon, b + \epsilon]$.
- $\mathbf{P} = \{\text{closed subsets of a metric space } X\}$,
 $\Omega_\epsilon(V) = V^\epsilon = \{x \in X \text{ such that } d(x, V) \leq \epsilon\}$.

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Interleaving distance (Bubenik, dS, Scott 2015)

Given a superlinear family (Ω_ϵ) of translations of \mathbf{P} , we define the interleaving distance

$$d_i(\mathbb{V}, \mathbb{W}) = \inf (\epsilon \mid \mathbb{V}, \mathbb{W} \text{ are } \Omega_\epsilon\text{-interleaved})$$

between generalized persistence modules $\mathbb{V}, \mathbb{W} : \mathbf{P} \rightarrow \mathbf{C}$.

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Given a sublinear projection family $\pi : \mathbf{Trans}_{\mathbf{P}} \rightarrow [0, \infty]$, we define the interleaving distance

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Story 3: Interleaving Metrics

Functoriality

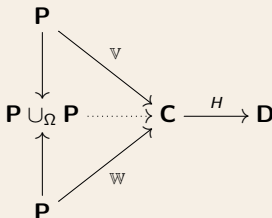
Suppose $\mathbb{V}, \mathbb{W} : \mathbf{P} \rightarrow \mathbf{C}$ and $H : \mathbf{C} \rightarrow \mathbf{D}$ are functors. Then

$$d_i(H\mathbb{V}, H\mathbb{W}) \leq d_i(\mathbb{V}, \mathbb{W})$$

for any superlinear family or sublinear projection.

Proof.

An Ω -interleaving of \mathbb{V}, \mathbb{W} gives an Ω -interleaving of $H\mathbb{V}, H\mathbb{W}$:



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Example: sublevelset persistence

Let X be a topological space and $f, g : X \rightarrow \mathbf{R}$ be functions with $\|f - g\|_\infty \leq \epsilon$.

- The persistence modules $\mathbb{V}, \mathbb{W} : \mathbf{R} \rightarrow \mathbf{Top}$ defined

$$\mathbb{V}(t) = f^{-1}(-\infty, t], \quad \mathbb{W}(t) = g^{-1}(-\infty, t],$$

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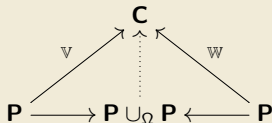
- For any homology functor $H : \mathbf{Top} \rightarrow \mathbf{Vect}$, the persistence modules $H\mathbb{V}, H\mathbb{W} : \mathbf{R} \rightarrow \mathbf{Vect}$ are ϵ -interleaved.

Story 3: Interleaving Metrics

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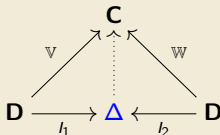
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Story 3: Interleaving Metrics

Interleavings for generalized persistence modules over an arbitrary category

Two persistence modules $\mathbb{V}, \mathbb{W} : \mathbf{D} \rightarrow \mathbf{C}$ are **Δ -interleaved** iff the following functor extension problem has a solution:



Here Δ is a **cospan**. The two functors l_1, l_2 are full-and-faithful. Every object of Δ is of the form $l_1(d)$ or $l_2(d)$.

Story 3: Interleaving Metrics

Example: dynamical system interleavings

Let \mathbf{D} be the category defined by the directed graph



Thus \mathbf{D} has one object and morphisms $\{0, 1, 2, 3, \dots\}$.

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with addition as composition.

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- Δ_n -interleavings are **shift-equivalences**.

Story 4: Set-Valued Persistence Modules

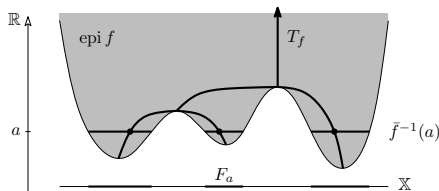
Merge trees (Cagliari, Ferri, Pozzi 2001, & Morozov, Beketayev, Weber 2013)

- A functor $\mathbf{T} : \mathbf{R} \rightarrow \mathbf{Set}$ can be thought of as a **merge tree**.
- Let X be a topological space and $f : X \rightarrow \mathbf{R}$ a function. Then

$$\mathbf{T}(t) = \pi_0 f^{-1}(-\infty, t]$$

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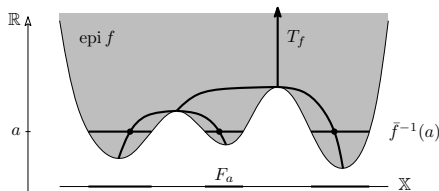
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Story 4: Set-Valued Persistence Modules

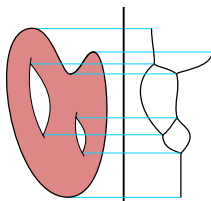
Reeb graphs (dS, Munch, Patel 2016)

- A functor $\mathbf{F} : \mathbf{Int} \rightarrow \mathbf{Set}$ can be thought of as a **graph over the real line**. (Technically we require \mathbf{F} to satisfy a **cosheaf condition**.)
- Let X be a topological space and $f : X \rightarrow \mathbf{R}$ a function. Then

$$\mathbf{F}_f(I) = \pi_0 f^{-1}(I)$$

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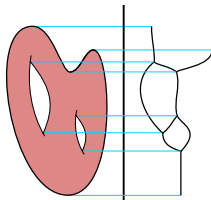
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Story 5: Reeb Graphs & Reeb Cosheaves

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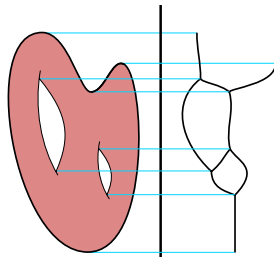
Reeb graphs

- An **R-space** (X, f) is a topological space X with a map $f : X \rightarrow \mathbf{R}$.
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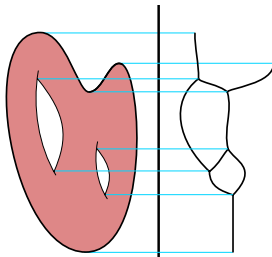
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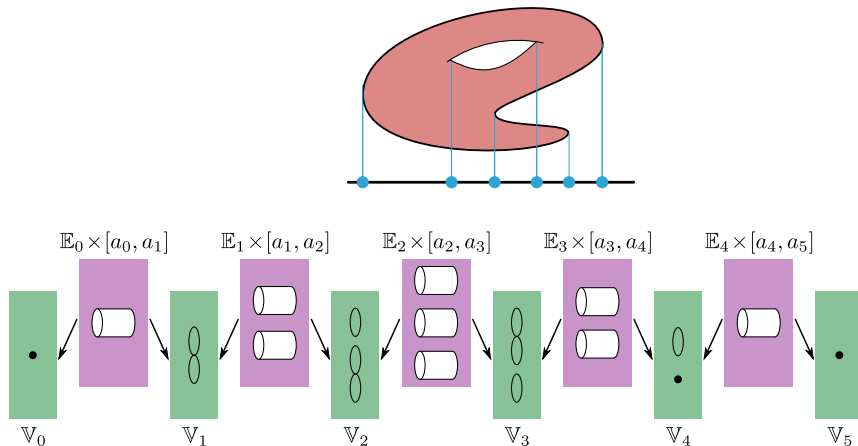
Reeb functor

- The **Reeb functor** converts a (constructible) **R-space** into a Reeb graph:

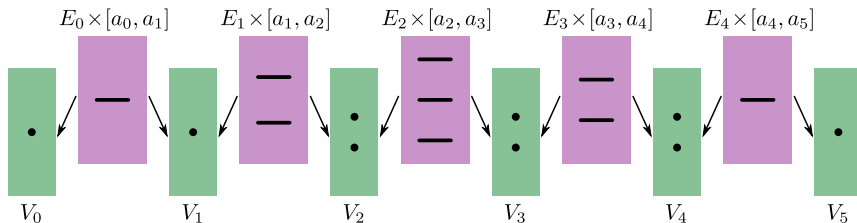
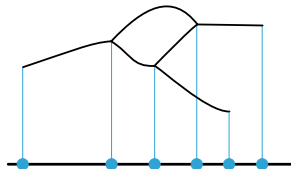
$$(X, f) \mapsto ((X/\sim), \bar{f})$$

where $x \sim y$ iff x, y are in the same component of the same levelset of f .

Story 5: Reeb Graphs & Reeb Cosheaves



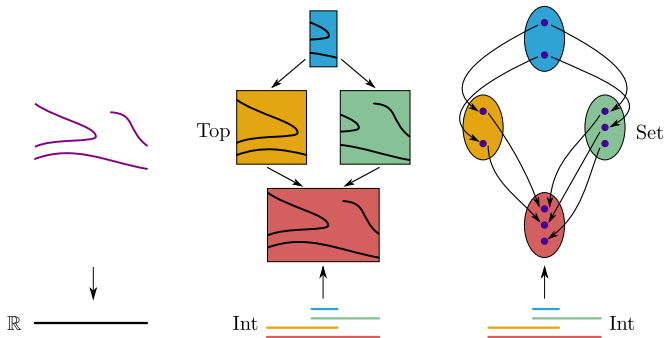
Story 5: Reeb Graphs & Reeb Cosheaves



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Reeb cosheaves (dS, Munch, Patel 2016)

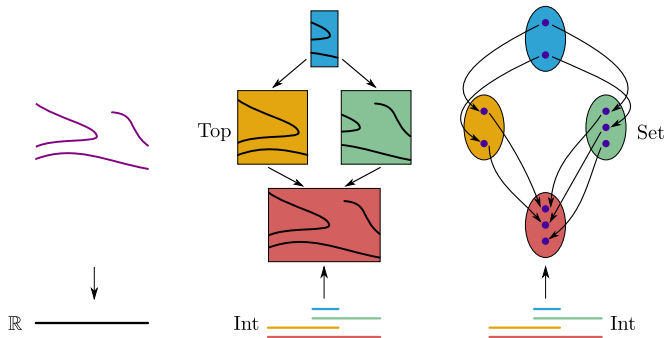
- Let **Int** denote the poset of open intervals, with respect to inclusion.
- A Reeb graph gives rise to a functor $\mathbf{F} : \mathbf{Int} \rightarrow \mathbf{Set}$ that is **constructible** and satisfies the **cosheaf condition** for unions of intervals.



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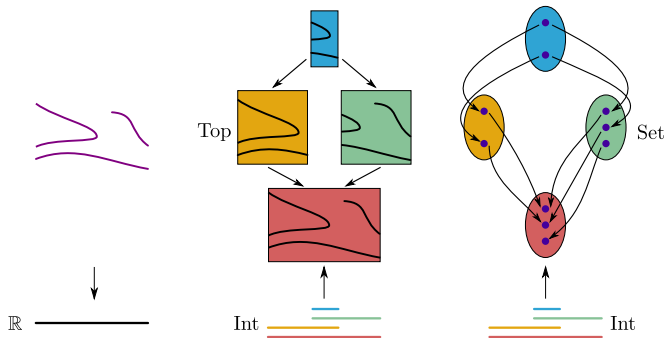


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Reeb cosheaves (dS, Munch, Patel 2016)

- Let **Int** denote the poset of open intervals, with respect to inclusion.
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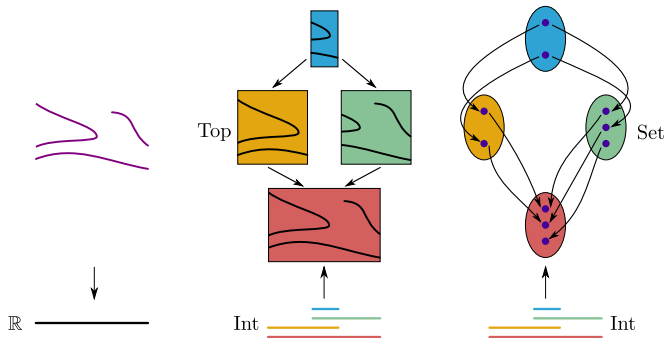


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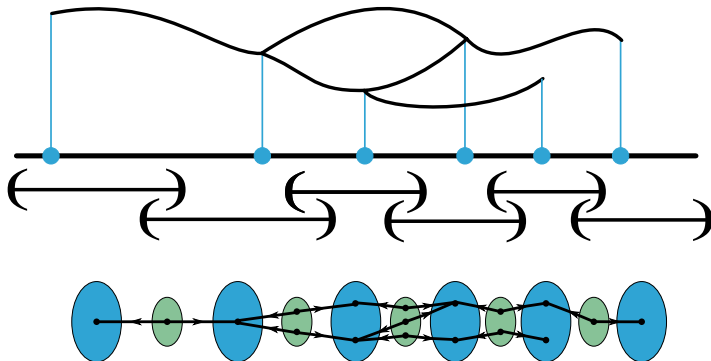


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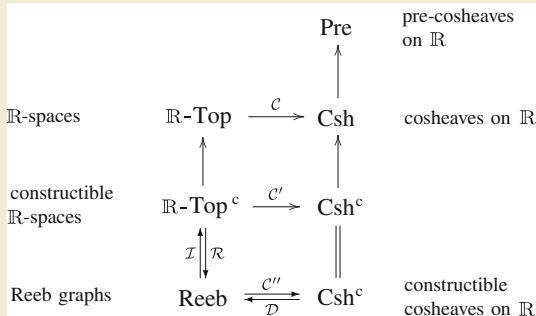
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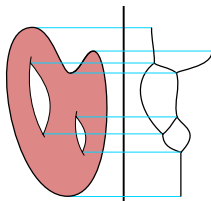
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Story 5: Reeb Graphs & Reeb Cosheaves



Reeb functor (two versions)

- The **Reeb functor** converts a (constructible) **R-space** into a Reeb graph:

$$(X, f) \longmapsto ((X/\sim), \bar{f})$$

where $x \sim y$ iff x, y are in the same component of the same levelset of f .

or

- The **Reeb functor** converts a constructible **R-space** into a Reeb cosheaf:

$$\mathbf{F}(I) = \pi_0 f^{-1}(I)$$

$$\mathbf{G}[I \subseteq J] = \pi_0 \left[f^{-1}(I) \subseteq f^{-1}(J) \right]$$

Story 5: Reeb Graphs & Reeb Cosheaves

Translation operators on \mathbf{Int}

We define a 1-parameter semigroup (Ω_ϵ) of functors $\mathbf{Int} \rightarrow \mathbf{Int}$ by setting

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Reeb interleaving

An ϵ -**interleaving** between \mathbf{F}, \mathbf{G} is given by two families of maps

$$\phi_I : \mathbf{F}(I) \rightarrow \mathbf{G}(I^\epsilon), \quad \psi_I : \mathbf{G}(I) \rightarrow \mathbf{F}(I^\epsilon)$$

which are natural with respect to inclusions $I \subseteq J$ and such that

$$\psi_{I^\epsilon} \circ \phi_I = \mathbf{F}[I \subseteq I^{2\epsilon}], \quad \phi_{I^\epsilon} \circ \psi_I = \mathbf{G}[I \subseteq I^{2\epsilon}]$$

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Stability Theorem

If $f, g : X \rightarrow \mathbf{R}$ with $\|f - g\|_\infty \leq \epsilon$ then $\mathrm{di}(\mathbf{F}, \mathbf{G}) \leq \epsilon$.

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If $\mathbf{F} : \mathbf{Int} \rightarrow \mathbf{Set}$ is a (constructible) cosheaf, then so is $\mathbf{F}\Omega_\epsilon : \mathbf{Int} \rightarrow \mathbf{Set}$.

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There is a 1-parameter semigroup of 'smoothing' operations on Reeb graphs.

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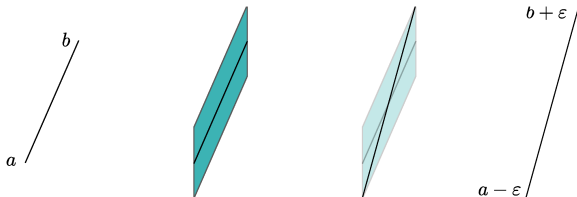
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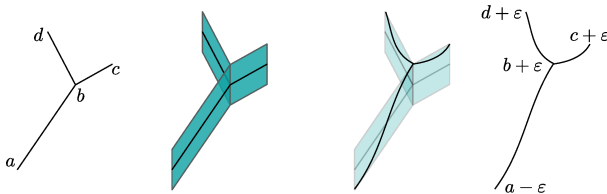
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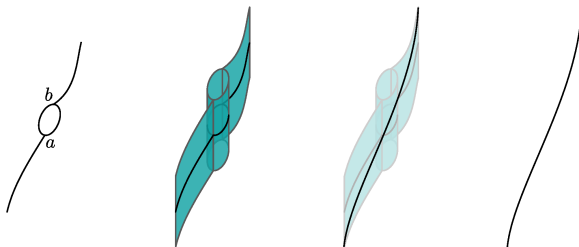
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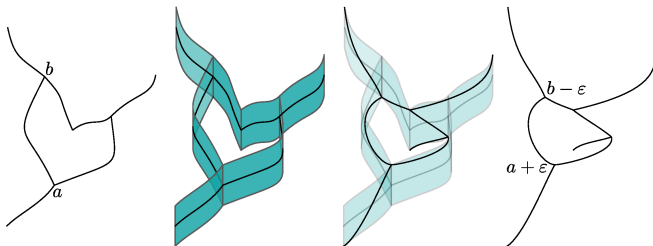
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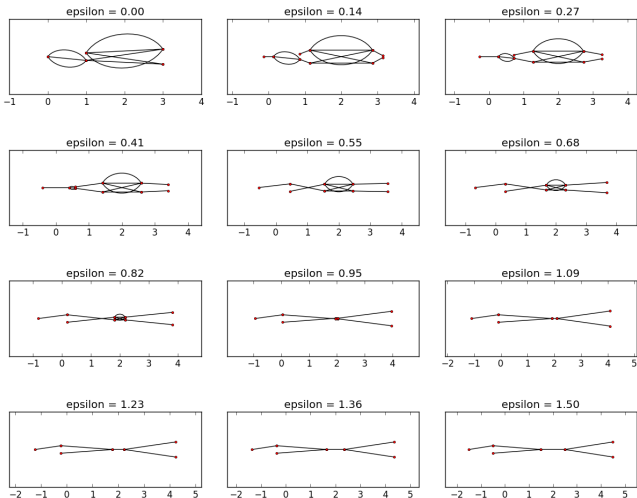
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Progressive smoothing algorithm by Dmitriy Smirnov & Song Yu:



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Discretized Reeb Graphs

- A discrete Reeb graph is a diagram

$$E \begin{array}{c} \xrightarrow{\ell} \\ \xrightarrow{r} \end{array} V \xrightarrow{\phi} \mathbf{R}$$

where E, V are finite sets and $\phi\ell(e) < \phi r(e)$ for each $e \in E$.

- Each $v \in V$ has a **left-** and **right-degree**:

$$\deg_l(v) = \#r^{-1}(v), \quad \deg_r(v) = \#\ell^{-1}(v), \quad \deg(v) = (\deg_l(v), \deg_r(v)).$$

- The discrete Reeb graph is **reduced** if $\deg(v) \neq (1, 1)$ for all v .

The **critical radius** of a reduced graph is

$$\epsilon_{crit} = \frac{1}{2} \min \{ \phi r(e) - \phi \ell(v) \mid e \in E, \deg_r(\ell(e)) > 1, \deg_l(r(e)) > 1 \}$$

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Progressive smoothing algorithm by Dmitriy Smirnov & Song Yu:

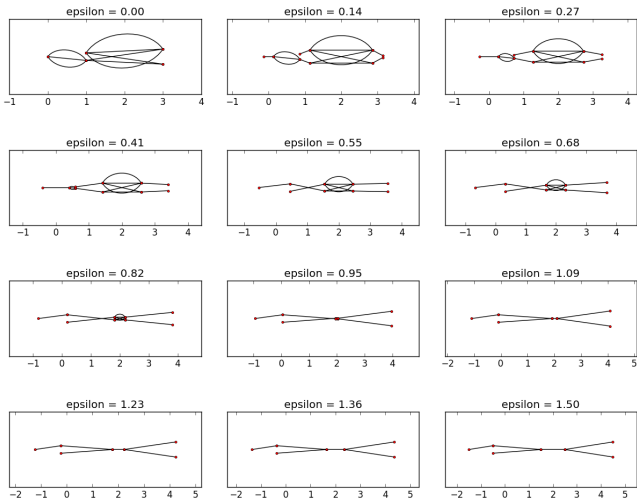
Algorithm: smooth by ϵ

- If $\deg(v) = (1, ?)$ then v moves by $+\epsilon$.
- If $\deg(v) = (?, 1)$ then v moves by $-\epsilon$.
- If $\deg(v) = (?, ?)$ then split v into two and move by $\pm\epsilon$.

Valid up to the critical radius. Recompute at critical radius and recurse.

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Story 6: Generalised Factors

Image persistence (Cohen-Steiner, Edelsbrunner, Harer, Morozov 2009)

Let $\mathbb{V}, \mathbb{W} : \mathbf{P} \rightarrow \mathbf{Vect}$ be persistence modules and let $\Phi : \mathbb{V} \Rightarrow \mathbb{W}$. Then we can define a persistence module $\text{Im}(\Phi)$ with

- $[\text{Im}(\Phi)](t) = \text{Im}(V_t \xrightarrow{\phi_t} W_t)$ for all t .
- $[\text{Im}(\Phi)](s \leq t)$ = the map induced by the horizontal maps in:

$$\begin{array}{ccc} V_s & \longrightarrow & V_t \\ \phi_s \downarrow & & \downarrow \phi_t \\ W_s & \longrightarrow & W_t \end{array}$$

We can similarly define $\text{Ker}(\Phi)$ and $\text{Coker}(\Phi)$.

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Example

Suppose $p : X \rightarrow Y$ is a map of spaces, $f : X \rightarrow \mathbf{R}$, and $g : Y \rightarrow \mathbf{R}$. If $f \leq gp$, then p carries the t -sublevelset of f into the t -sublevelset of g , for all t , and the persistence module $\text{Im}(H(p))$ is defined.

Story 6: Generalised Factors

Three ways of thinking of a map between persistence modules (over \mathbf{N} , say)

A functor $\mathbf{2} \rightarrow \mathbf{Vect}^{\mathbf{N}}$:

$$\begin{array}{ccccccc} F_0 & \longrightarrow & F_1 & \longrightarrow & F_2 & \longrightarrow & \cdots \\ & & & & \Downarrow & & \\ G_0 & \longrightarrow & G_1 & \longrightarrow & G_2 & \longrightarrow & \cdots \end{array}$$

A functor $\mathbf{N} \times \mathbf{2} \rightarrow \mathbf{Vect}$:

$$\begin{array}{ccccccc} F_0 & \longrightarrow & F_1 & \longrightarrow & F_2 & \longrightarrow & \cdots \\ \downarrow & & \downarrow & & \downarrow & & \\ G_0 & \longrightarrow & G_1 & \longrightarrow & G_2 & \longrightarrow & \cdots \end{array}$$

A functor $\mathbf{N} \rightarrow \mathbf{Vect}^{\mathbf{2}}$:

$$\begin{array}{ccccccc} F_0 & & F_1 & & F_2 & & \cdots \\ \downarrow & \Rightarrow & \downarrow & \Rightarrow & \downarrow & \Rightarrow & \\ G_0 & & G_1 & & G_2 & & \cdots \end{array}$$

Story 6: Generalised Factors

The exponential law

The following categories of functors

$$(\mathbf{D}^{\mathbf{P}})^{\mathbf{W}} = \mathbf{D}^{\mathbf{P} \times \mathbf{W}} = (\mathbf{D}^{\mathbf{W}})^{\mathbf{P}}$$

are equal for any three categories $\mathbf{D}, \mathbf{P}, \mathbf{W}$.

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Image, Kernel, Cokernel functors

The operations Im , Ker and Coker can be thought of as functors $\mathbf{Vect}^2 \rightarrow \mathbf{Vect}$.

- Each operation converts any $(V \xrightarrow{\alpha} W)$ into a vector space.
- Given a commutative square, there are induced maps between images, kernels, cokernels.

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Proposition (Bubenik, dS, Scott)

The image persistence of $\Phi : \mathbb{V} \Rightarrow \mathbb{W}$ is equal to the composite

$$\mathbf{P} \xrightarrow{\hat{\Phi}} \mathbf{Vect}^2 \xrightarrow{\text{Im}} \mathbf{Vect}$$

where $\hat{\Phi}$ is the interpretation of Φ as a functor $\mathbf{P} \rightarrow \mathbf{Vect}^2$.

Story 6: Generalised Factors

Generalized factor persistence (Bubenik, dS, Scott)

Given

- a category of persistence modules $\mathbf{D}^{\mathbf{P}}$;
- a category \mathbf{W} , which we call the *auxiliary category*;
- a functor $\mathbf{D}^{\mathbf{W}} \xrightarrow{N} \mathbf{E}$, which we call the *generalized factor*.

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- a functor $\mathbf{D}^W \xrightarrow{N} \mathbf{E}$, which we call the *generalized factor*.

Then any functor $F : \mathbf{W} \rightarrow \mathbf{D}^P$ determines a persistence module in \mathbf{E}^P , by

$$\begin{array}{ccc} (\mathbf{D}^P)^W = \mathbf{D}^{W \times P} = (\mathbf{D}^W)^P & \longrightarrow & \mathbf{E}^P \\ F \longmapsto & \hat{F} \longmapsto & N\hat{F} \end{array}$$

Story 6: Generalised Factors

Reductions of 2-dimensional persistence

Let $\mathbb{V} = (V(s, t)) \in \mathbf{Vect}^{\mathbb{R} \times \mathbb{R}}$ be a two-dimensional persistence module. Think of this as a family (\mathbb{W}_t) of 1-dimensional persistence modules. We will define various generalized factors $N : \mathbf{Vect}^{\mathbb{R}} \rightarrow \mathbf{Vect}$.

- Fix a and define $N(\mathbb{W}) = \mathbb{W}(a)$.
- Fix $a < b$ and define $N(\mathbb{W}) = \text{Im}(\mathbb{W}(a) \rightarrow \mathbb{W}(b))$.
- Fix $a < b \leq c < d$ and define

$$N(\mathbb{W}) = \left[\frac{\text{Im}(\mathbb{W}(b) \rightarrow \mathbb{W}(c)) \cap \text{Ker}(\mathbb{W}(c) \rightarrow \mathbb{W}(d))}{\text{Im}(\mathbb{W}(a) \rightarrow \mathbb{W}(c)) \cap \text{Ker}(\mathbb{W}(c) \rightarrow \mathbb{W}(d))} \right]$$

Then there is a 1-parameter persistence module associated to each of these functors.

Story 6: Generalised Factors

Zigzag factors

Suppose \mathbf{Z} is the category defined by:



An element of $\mathbf{Vect}^{\mathbf{Z}}$ is a diagram

$$\mathbb{W} : \quad W_1 \xrightarrow{f} W_2 \xleftarrow{g} W_3 \xrightarrow{h} W_4$$

Then, for example, the functor $\mathbf{Vect}^{\mathbf{Z}} \rightarrow \mathbf{Vect}$ defined by

$$N(\mathbb{W}) = \left[\frac{g(h^{-1}(0))}{f(W_1)} \right]$$

picks out the part of \mathbb{W} supported over W_2, W_3 .

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Therefore, given a zigzag of persistence modules

$$\mathbb{V}_1 \xrightarrow{f} \mathbb{V}_2 \xleftarrow{g} \mathbb{V}_3 \xrightarrow{h} \mathbb{V}_4$$

we can construct a single persistence module which extracts the $[2, 3]$ part.

Story 7: The observable category

Tame persistence modules

Let $\mathbb{V} : \mathbf{R} \rightarrow \mathbf{Vect}$ be a persistence module. If the maps $V_s \rightarrow V_t$ have finite rank whenever $s < t$, then \mathbb{V} has a persistence diagram. If \mathbb{V} has an interval decomposition, then the summands are identified exactly by the points in the diagram. However, it is not guaranteed that \mathbb{V} has an interval decomposition.

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Ephemeral modules (Chazal, Crawley-Boevey, dS 2016)

A persistence module \mathbb{V} is *ephemeral* if $v_s^t = 0$ whenever $s < t$.

Then:

- The ephemeral modules comprise a **Serre subcategory** of the category of persistence modules.
- We can form the Serre quotient category by formally inverting all maps whose kernels and cokernels are ephemeral.
- In this category, every q-tame persistence module admits an interval decomposition.

Perhaps this is the ‘correct’ category for real-parameter persistence?

Definition

A **Serre subcategory** is a full subcategory \mathbf{C} of an Abelian category such that for any short exact sequence

$$0 \longrightarrow U \longrightarrow V \longrightarrow W \longrightarrow 0$$

we have

$$V \in \mathbf{C} \Leftrightarrow U \in \mathbf{C} \text{ and } W \in \mathbf{C}.$$

Equivalently, the subcategory \mathbf{C} is closed under subobjects, quotient objects, and extensions.

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Noise systems (Scolamiero et al., 2016)

Noise in topological data analysis can be studied by considering a nested family $(\mathbf{C}_\epsilon \mid \epsilon \in [0, \infty))$ satisfying an enriched version of the Serre conditions:

$$\begin{aligned} V \in \mathbf{C}_\epsilon &\Rightarrow U \in \mathbf{C}_\epsilon \text{ and } W \in \mathbf{C}_\epsilon \\ V \in \mathbf{C}_{\epsilon_1 + \epsilon_2} &\Leftarrow U \in \mathbf{C}_{\epsilon_1} \text{ and } W \in \mathbf{C}_{\epsilon_2}. \end{aligned}$$

for any short exact sequence.

Acknowledgements

Collaborators

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