

Parallel Transport in Low Dimensions

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1. ordinary connections in terms of parallel transport functors

- (a) the usual definition of a connection in terms of a choice of horizontal subspaces
- (b) this allows to lift vectors and, in turn, paths, from base space to the total space
- (c) this lift amounts to a functor from paths to torsor isomorphisms
- (d) functors from paths to torsor isomorphisms arising this way have the special property of having smooth local trivializations
- (e) this smooth local trivialization of a functor reproduces the familiar differential cocycle relations
- (f) we can combine paths in patches with jumps between patches to a category that covers the original path category, and how our local data defines a functor on that cover
- (g) this is an example of an anafunctor

2. categorified parallel transport: 2-anafunctors

- (a) first categorify the domain: 2-paths
- (b) then categorify the codomain: 2-groups and 2-group torsors
- (c) finally categorify the notion of "smooth local trivialization"; draw the same diagram as before, but explain how now the triangle is filled by a 2-morphism that makes a tetrahedron 2-commute
- (d) this is an example of descent data that might be addressed as a 2-anafunctor
- (e) the main example: 2-anafunctor with values in strict 2-group
- (f) in particular: the issue of fake flatness

3. Chern-Simons transport

- (a) warmup: 2-transport with values in $\text{INN}(G)$ is the same as G 1-transport
- (b) there is a general principle behind this: Schreier theory
- (c) this makes us want to look at $\text{INN}(G_2)$
- (d) the curvature and Bianchi identities of $\text{INN}(G_2)$ -transport; these characterize the corresponding Lie 3-algebra
- (e) fact: there is a Lie 3-algebra, $\text{cs}(g)$, such that connections with values in it come from Chern-Simons 3-forms
- (f) fact: $\text{cs}(g)$ sits inside $\text{Lie}(\text{INN}(\text{String}(G)))$

1 Parallel 1-Transport: the motivating example

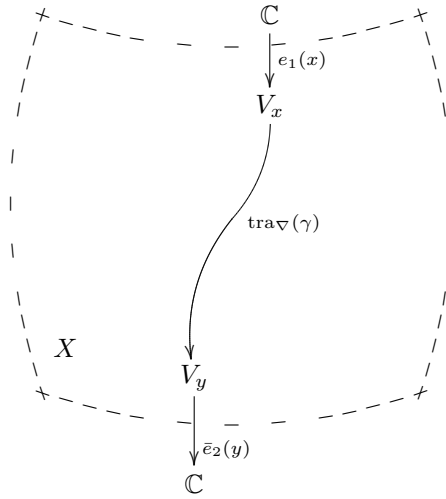
Models of the physics of charged particles are usually formulated in terms of vector bundles

$$V \rightarrow X$$

with connection

$$\nabla.$$

The part of this formalism most directly connected to what we actually observe in nature is the **parallel transport**.



1.1 Connections give rise to parallel transport functors.

One way to think of a connection of a principal bundle is to say that a connection is a prescription that tells us at each point of a principal bundle which tangent vectors are supposed to be **parallel** to the base space.

More precisely:

Definition 1 (connection in terms of horizontal subspaces) Let $p : B \rightarrow X$ be a smooth principal G -bundle. For each point $b \in B$ of the total space, let

$$V_b := \ker(p_b^*) \subset T_b B$$

be the **vertical subspace** of the tangent space at that point. V_b is the space of vectors at b that are tangent to the fiber.

Then a **connection** on the principal bundle is a smooth G -invariant choice of complements H_b of V_b

$$T_b B = V_b \oplus H_b$$

$$H_{gb} = g^* H_b$$

for all $b \in B$. H_b is called the **horizontal subspace** of the tangent space at b .

For our purposes, the point of this definition is the following: since p_b^* restricted to H_b is an isomorphism, it follows that a connection allows us to **lift** vectors in base space to **parallel** vectors on the total space.

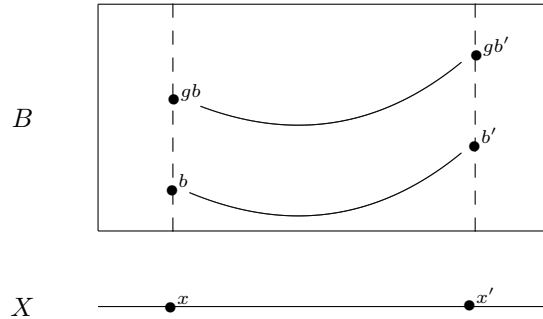
We can integrate this procedure and find for each path

$$x \xrightarrow{\gamma} x'$$

in base space a path

$$b \xrightarrow{\tilde{\gamma}} b'$$

in the total space, which is everywhere parallel to γ .



We say that b' is obtained from **parallel transporting** b along γ from the fiber B_x to the fiber $B_{x'}$.

This way a connection assigns, by parallel transport, to each path γ in base space a map

$$\text{tra}(\gamma) : B_x \rightarrow B_{x'}$$

between the fibers over the endpoints

This assignment of maps between fibers to paths in base space has some special properties:

- The G -invariance of the choice of horizontal subspaces implies that these maps between the fibers commute with the G -action on the fibers.
- In particular, this implies that these maps are invertible, since G acts freely and transitively on each fiber.
- The map $\text{tra}(\gamma)$ is independent of the parameterization of γ .
- If $\tilde{\gamma}$ is obtained from γ by reversing the direction, then $\text{tra}(\tilde{\gamma})$ is the inverse of $\text{tra}(\gamma)$.
- If γ is the composition of two paths γ_1 and γ_2 , then

$$\text{tra}(\gamma) = \text{tra}(\gamma_2) \circ \text{tra}(\gamma_1).$$

choice of horizontal subspaces $H_b \subset T_b B$	choice of functor $\mathcal{P}_1(X) \rightarrow G\text{Tor}$
allows to lift vectors $v \in T_{p(b)} X$ to parallel vectors $\tilde{v} \in T_b B$	allows to lift paths $\gamma \in \mathcal{P}_1(X)$ to fiber isomorphisms $\text{tra}(\gamma) : B_x \rightarrow B_y$
differential description of connection	integrated description of connection

Table 1: The ordinary definition of a **connection** on a principal bundle in terms of horizontal subspaces can be understood as the differential description of the concept of parallel transport.

Clearly, all this is trying to tell us that **parallel transport is a functor**

$$\text{tra} : \mathcal{P}_1(X) \rightarrow G\text{Tor}$$

that sends paths in base space to morphisms of G -torsors.

To make this precise, we need to specify what the groupoid of paths in base space that we are talking about is like.

Definition 2 *The objects of $\mathcal{P}_1(X)$ are points in X . The morphisms*

$$x \xrightarrow{\gamma} y$$

of $\mathcal{P}_1(X)$ are equivalence classes of smooth maps $\gamma : [0, 1] \rightarrow X$ with $\gamma(0) = x$ and $\gamma(1) = y$, which are constant in a neighbourhood of 0 and in a neighbourhood of 1, and where two maps are considered equivalent if they are related by an orientation-preserving diffeomorphism. Composition of morphisms is by the obvious concatenation of these maps, modulo the relation that paths related by an orientation reversing diffeomorphism are mutually inverse.

Given this definition of the groupoid of paths in X , our list of properties of parallel transport implies

Proposition 1 *Given a principal G -bundle with connection $B \rightarrow X$, parallel transport in that bundle is a functor*

$$\text{tra} : \mathcal{P}_1(X) \rightarrow G\text{Tor}.$$

The entire discussion generalizes directly to associated bundles.

Proposition 2 *Given a vector bundle with connection $V \rightarrow X$, parallel transport in that bundle is a functor*

$$\text{tra} : \mathcal{P}_1(X) \rightarrow \text{Vect}.$$

It would be nice if these statements had a converse. We cannot expect every functor from paths to G -torsors to define a smooth principal bundle, or from paths to vector spaces to define a smooth vector bundle.

But transport functors that are *smooth* and *locally trivializable* in some suitable sense should do.

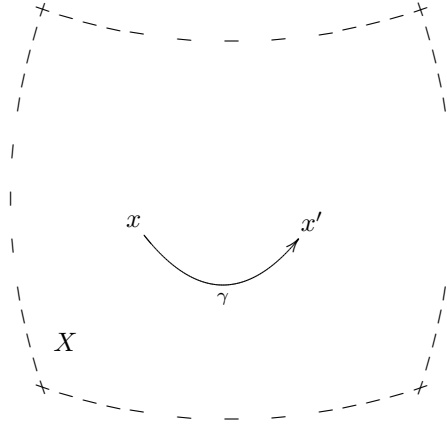


Figure 1: A **morphisms of the path groupoid** of X is an oriented path γ cobounding two points x and x' in X . Paths that differ by orientation-preserving diffeomorphism are identified. This ensures strict associativity of composition. Paths that differ by orientation-reversing diffeomorphism are taken to represent mutually inverse morphisms.

1.2 Locally trivial smooth transport functors

Let us locally trivialize our principal bundle and see what this does to the corresponding parallel transport functor.

So we choose a good cover

$$p : U \rightarrow X$$

of base space by open contractible sets.

This allows us to pull back structures over X to U , where they may be trivializable.

For the bundle $B \rightarrow X$ this means that we can choose a bundle isomorphism

$$t : \pi^* B \xrightarrow{\sim} U \times G .$$

This amounts to choosing, in a smooth way, for each point (x, i) in U an isomorphism of G -torsors

$$t(x, i) : B_x \xrightarrow{\sim} G .$$

We can translate this into a similar local trivialization of the corresponding

local connection 1-form $A \in \Omega^1(U, \mathfrak{g})$	smooth transport functor $\text{tra}_U : \mathcal{P}_1(U) \rightarrow \Sigma(G)$
transition function $g \in \Omega^0(U^{[2]}, G)$	natural isomorphism $g : p_1^* \text{tra}_U \rightarrow p_2^* \text{tra}_U$
$A_i = g_{ij} A_j g_{ij}^{-1} + g_{ij} dg_{ij}^{-1}$	
$g_{ij} g_{jk} = g_{ik}$	

Table 2: The **differential cocycle** data describing the local trivialization of a principal bundle with connection is the **descent data** of a local smooth transport functor.

transport functor. Its pullback to U is

$$\begin{array}{ccc}
 \mathcal{P}_1(U) & \xrightarrow{p} & \mathcal{P}_1(X) \\
 & & \downarrow \text{tra} \\
 & & G\text{Tor}
 \end{array}$$

Which structure on U could this be isomorphic to? Notice that, using the identifications of fibers with copies of G above, we can form a functor

$$\text{tra}_U : \mathcal{P}_1(U) \rightarrow \text{Aut}_{G\text{Tor}}(G)$$

such that

$$\begin{array}{ccc}
 B_x & \xrightarrow{\text{tra}(\gamma)} & B_y \\
 \downarrow t(x,i) & & \downarrow t(y,i) \\
 G & \xrightarrow{\text{tra}_U(\gamma)} & G
 \end{array}$$

for any path γ in U_i .

An automorphism of G regarded as a G -torsor over itself is nothing but an element of G . In other words, we have a canonical injection

$$i : \Sigma(G) \xrightarrow{\subset} G\text{Tor}$$

of the category $\Sigma(G)$ with a single object and G worth of morphisms into the category of G -torsors.

Using this injection, we can think of tra_U as a functor with values in $G\text{Tor}$ that factors through i :

$$\begin{array}{ccc} \mathcal{P}_1(U) & & \\ \text{tra}_U \downarrow & & \\ \Sigma(G) & \xrightarrow{\rho} & G\text{Tor} \end{array}$$

Taken together, we find that the local trivialization $t : p^*B \rightarrow U \times G$ of the principal bundle corresponds to a morphism

$$\begin{array}{ccc} \mathcal{P}_1(U) & \xrightarrow{p} & \mathcal{P}_1(X) \\ \text{tra}_U \downarrow & \swarrow \tilde{\sim}_t & \downarrow \text{tra} \\ \Sigma(G) & \xrightarrow{i} & G\text{Tor} \end{array}$$

that relates the corresponding parallel transport functor to a functor on $\mathcal{P}_1(U)$ with values in $\Sigma(G)$.

In contrast to the category $G\text{Tor}$, the category $\Sigma(G)$ is naturally a **smooth category**, namely a category internal to smooth spaces. The same is true for the path groupoid. Since the bundle with connection that we started with was smooth, the functor tra_U is a smooth functor between smooth categories.

The smoothness of a smooth functor implies that the functor is specified by its derivatives. Functoriality then implies that already the derivatives at all identity morphisms suffice:

Proposition 3 *Smooth functors*

$$\text{tra}_U : \mathcal{P}_1(U) \rightarrow \Sigma(G)$$

are in bijection with \mathfrak{g} -valued 1-forms A on U :

$$\text{tra}_U(\gamma) = \text{P exp} \left(\int_0^1 \gamma^* A \right).$$

On double intersections of the cover, the local trivialization of our bundle yields a smooth natural isomorphism

$$g := p_2^* t \circ p_1^* t : p_1^* \text{tra}_U \rightarrow p_2^* \text{tra}_U.$$

Proposition 4 *Such smooth natural isomorphisms between smooth functors coming from 1-forms A and A' , respectively, are in bijection with smooth functions g with values in G such that*

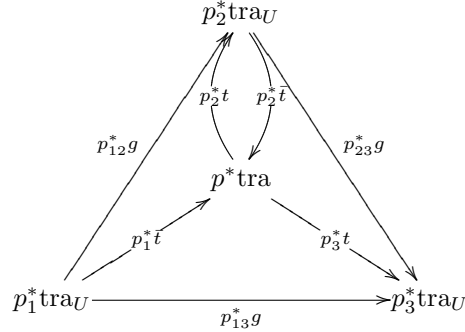
$$A = gA'g^{-1} + gdg^{-1} .$$

These G -valued functions are nothing but the **transition function** describing the local trivialization of our bundle B .

The cocycle condition

$$g_{ij}g_{jk} = g_{ik}$$

which they satisfy is an expression of the existence of this triangle:



Using the familiar fact that principal G -bundles with connection are equivalent to differential 1-cocycles, we find that principal G -bundles with connection are equivalent to descent data

$$\text{Trans}_{i,p}$$

for smooth transport functors taking values in $\Sigma(G)$.

Again, all these considerations go through completely analogously for vector bundles. All we need to do is to replace the injection

$$i : \Sigma(G) \longrightarrow G\text{Tor}$$

by a *representation*

$$\rho : \Sigma(G) \longrightarrow \text{Vect} .$$

1.3 Anafunctors

There is an equivalent way to talk about functors on paths of a cover that are related by isomorphisms on double intersections such that a triangle commutes on triple intersections.

As Toby Bartels and John Baez emphasized, we can think of this situation as characterizing an **anafunctor** - a functor not directly acting on its domain, but on a *cover* of that domain.

Namely, if we let

$$\mathcal{P}_1(U^\bullet)$$

be the category whose morphisms are combinations of paths in U with jumps from one patch into the other, then our locally trivial transport functor tra_U with transitions g is encoded in the span

$$\begin{array}{ccc} \mathcal{P}_1(U^\bullet) & \xrightarrow{(\text{tra}_U, g)} & \Sigma(G) . \\ \downarrow p & & \\ \mathcal{P}_1(X) & & \end{array}$$

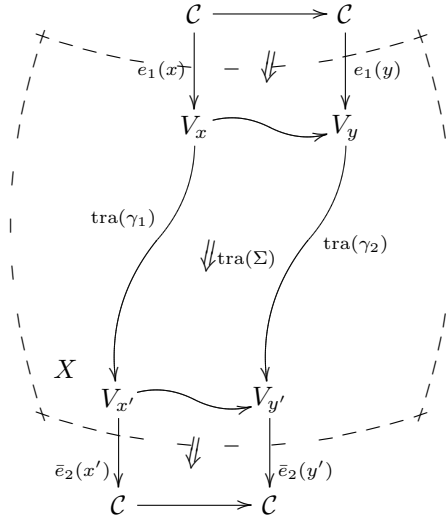
2 Parallel 2-Transport: 2-Bundles with Connection

We wish to categorify the description of bundles with connection in terms of descent data of smooth parallel transport functors.

This requires that we

- a) find suitable categorifications of the domain $\mathcal{P}_1(X)$ and codomain, $G\text{Tor}$ or Vect of our parallel transport functors
- b) find a suitable categorification of the descent data, i.e. find a suitable notion of **2-anafunctor**.

When such a categorification is available, we can study parallel transport of strings across surfaces:



parallel 1-transport	parallel 2-transport
domain: path groupoid $\mathcal{P}_1(X)$	domain: 2-path 2-groupoid $\mathcal{P}_2(X)$
codomain: vector spaces Vect	codomain: 2-vector spaces 2Vect
structure group: G	structure 2-group: G_2
representation $\rho : \Sigma(G) \rightarrow \text{Vect}$	2-representation $\rho : \Sigma(G_2) \rightarrow 2\text{Vect}$
trivial vector bundles with connection	trivial 2-vector bundle with connection
smooth functors: $\mathcal{P}_1(X) \longrightarrow \Sigma(G) \xrightarrow{\rho} \text{Vect}$	smooth 2-functor: $\mathcal{P}_2(X) \longrightarrow \Sigma(G_2) \xrightarrow{\rho} 2\text{Vect}$
vector bundles with connection	2-vector bundle with connection
smooth anafunctors: $\mathcal{P}_1(U^\bullet) \longrightarrow \Sigma(G) \xrightarrow{\rho} \text{Vect}$	smooth 2-anafunctor: $\mathcal{P}_2(U^\bullet) \longrightarrow \Sigma(G_2) \xrightarrow{\rho} 2\text{Vect}$
$p \downarrow$ $\mathcal{P}_1(X)$	$p \downarrow$ $\mathcal{P}_2(X)$

Table 3: On the left, our description of bundles with connection in terms of parallel transport functors. On the right our categorification of this situation.

2.1 2-Paths

We want 2-morphisms in $\mathcal{P}_2(X)$ to look like little surface elements. There are various choices one could make concerning the degree of invertibility and strictness of composition of the 1-morphisms involved. For our purposes, it is useful to make

Definition 3 *The 2-path 2-groupoid $\mathcal{P}_2(X)$ has as objects the points of X , has as morphisms classes of oriented paths in X modulo orientation preserving diffeomorphism, and 2-morphisms thin homotopy classes of oriented surfaces cobounding such paths.*

The 2-path 2-groupoid is a strict 2-category. Composition is strictly associative. However, it is not a strict 2-groupoid, since a path is not strictly but weakly inverse to its orientation-reversed path.

2.2 2-Groups

From the point of view of parallel transport, structure groups G arise as the local trivializations of the transport groupoid. Hence the important characterizing property is that a group is a groupoid with a single object.

This immediately suggests the kind of categorification we need

Definition 4 *A 2-group is a 2-groupoid with a single object.*

Here we want to work within the 3-category of *strict* 2-categories, strict 2-functors between them, pseudonatural transformations between those and mod-

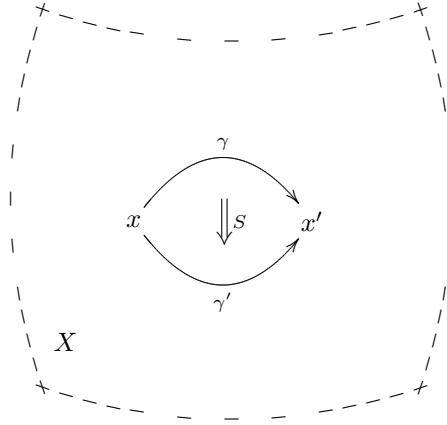


Figure 2: A **2-morphisms of the 2-path 2-groupoid** of X is a thin-homotopy class of a surface S cobounding two diffeomorphism classes γ and γ' of paths which in turn cobound two points in x and x' .

ificatiopn between the latter. For that reason we restrict attention to *strict 2-groups*.

Definition 5 A **strict 2-group** is a *strict 2-groupoid with a single object*.

Strict 2-groups turn out to have a useful description in terms of crossed modules.

Definition 6 A **crossed module** of groups is a pair (G_0, G_1) of groups, together with homomorphisms

$$G_1 \xrightarrow{t} G_0 \xrightarrow{\alpha} \text{Aut}(G_1)$$

such that t is equivariant with respect to the action induced by α , i.e. such that

$$\begin{array}{ccc} G_1 & \xrightarrow{\text{Ad}} & \text{Aut}(G_1) \\ & \searrow t & \nearrow \alpha \\ & G_0 & \end{array}$$

$$\Leftrightarrow \alpha(t(h))(h') = h h' h^{-1}$$

and such that

$$t(\alpha(g)(h)) = g t(h) g^{-1}.$$

Namely we have

Theorem 1 (classic) *The 2-category of 2-groups is equivalent to the 2-category of crossed modules.*

This equivalence is induced by identifying G_0 with the set of morphisms

$$\text{Mor}_1 = \left\{ \bullet \xrightarrow{g} \bullet \mid g \in G_0 \right\}$$

of the 2-groupoid; G_1 with the kernel of the source map, i.e. with those 2-morphisms starting at the identity

$$\left\{ \begin{array}{c} \bullet \begin{array}{c} \xrightarrow{\text{Id}} \\ \Downarrow h \\ \xrightarrow{t(h)} \end{array} \bullet \\ \left| \right. \\ h \in G_1 \end{array} \right\};$$

and the set of all morphisms with the semidirect product $G_1 \times G_0$ as

$$\text{Mor}_2 = \left\{ \bullet \begin{array}{c} \xrightarrow{\text{Id}} \\ \Downarrow h \\ \xrightarrow{t(h)} \end{array} \bullet \xrightarrow{g} \bullet \left| \right. h \in G_1, g \in G_0 \right\}.$$

The main fact to keep in mind, especially for the discussion in section ??, is the following rule for horizontal and vertical composition of 2-group elements (their precise form depends on some conventions that we chose to fix):

$$\begin{array}{c} \bullet \begin{array}{c} \xrightarrow{g_1} \\ \Downarrow h_1 \\ \xrightarrow{g'_1} \end{array} \bullet \begin{array}{c} \xrightarrow{g_2} \\ \Downarrow h_2 \\ \xrightarrow{g'_2} \end{array} \bullet = \bullet \begin{array}{c} \xrightarrow{g_1 \cdot g_2} \\ \Downarrow h_1 \cdot \alpha(g_1)(h_2) \\ \xrightarrow{g'_1 \cdot g'_2} \end{array} \bullet \end{array}$$

and

$$\bullet \begin{array}{c} \xrightarrow{g_1} \\ \Downarrow h_1 \\ \xrightarrow{g_2} \\ \Downarrow h_2 \\ \xrightarrow{g_3} \end{array} \bullet = \bullet \begin{array}{c} \xrightarrow{g_1} \\ \Downarrow h_2 \cdot h_1 \\ \xrightarrow{g_3} \end{array} \bullet,$$

where the dot on the right hand side indicates the ordinary product in the respective group.

Example 1

The two standard classes of examples for strict 2-groups and crossed modules are the following:

- Let G be any group, regarded as a groupoid with a single object. Then the automorphism functor 2-category $\text{Aut}_{\text{Cat}}(G)$ is a 2-group. It corresponds to the crossed module

$$G \xrightarrow{\text{Ad}} \text{Aut}(G) \xrightarrow{\text{Id}} \text{Aut}(G) .$$

- Every central extension

$$1 \longrightarrow K \longrightarrow H \longrightarrow G \longrightarrow 1$$

with the usual action of G on H defines a crossed module.

2.3 Transition of 2-functors and 2-anafunctors

We say a 2-functor $\text{tra} : \mathcal{P}_2(X) \rightarrow T$ is p -locally i -trivializable if there is an equivalence

$$\begin{array}{ccc} \mathcal{P}_2(U) & \xrightarrow{p} & \mathcal{P}_2(X) . \\ \text{tra}_U \downarrow & \swarrow \sim_t & \downarrow \text{tra} \\ T' & \xrightarrow{i} & T \end{array}$$

Proposition 5 *The resulting transitions*

$$\begin{array}{ccc} & p_2^* \text{tra}_U & \\ p_{12}^* g \nearrow & & \searrow p_{23}^* g \\ p_1^* \text{tra}_U & \xrightarrow{p_{13}^* g} & p_3^* \text{tra}_U \\ & f \Downarrow & \end{array} \quad \equiv \quad \begin{array}{ccc} & p_2^* \text{tra}_U & \\ p_{12}^* g \nearrow & \begin{array}{c} p_2^* t \Downarrow p_2^* \tilde{t} \\ p^* \text{tra} \end{array} & \searrow p_{23}^* g \\ p_1^* \text{tra}_U & \xrightarrow{p_{13}^* g} & p_3^* \text{tra}_U \\ & \begin{array}{c} p_1^* \tilde{t} \nearrow \\ p_3^* \tilde{t} \searrow \\ \sim \Downarrow \end{array} & \end{array}$$

1-anafunctors	2-anafunctors
$\begin{array}{ccc} \mathcal{P}_1(U) & \xrightarrow{p} & \mathcal{P}_1(X) \\ \text{tra}_U \downarrow & \searrow \sim_t & \downarrow \text{tra} \\ T' & \xrightarrow{i} & T \end{array}$	$\begin{array}{ccc} \mathcal{P}_2(U) & \xrightarrow{p} & \mathcal{P}_2(X) \\ \text{tra}_U \downarrow & \searrow \sim_t & \downarrow \text{tra} \\ T' & \xrightarrow{i} & T \end{array}$
$\begin{array}{ccc} & p_2^* \text{tra}_U & \\ p_{12}^* g \nearrow & & \searrow p_{23}^* g \\ p_1^* \text{tra}_U & \xrightarrow{p_{13}^* g} & p_3^* \text{tra}_U \end{array}$	$\begin{array}{ccc} p_2^* \text{tra}_U & \xrightarrow{p_{23}^* g} & p_3^* \text{tra}_U \\ p_{12}^* g \nearrow & p_{123}^* f \searrow & \nearrow p_{13}^* g \\ p_1^* \text{tra}_U & \xrightarrow{p_{14}^* g} & p_4^* \text{tra}_U \\ & \Downarrow p_{134}^* f & \downarrow p_{34}^* g \\ & & p_3^* \text{tra}_U \\ & & \searrow p_{234}^* f \\ & & p_2^* \text{tra}_U \\ & & \downarrow p_{24}^* g \\ & & p_1^* \text{tra}_U \\ & & \xrightarrow{p_{14}^* g} & p_4^* \text{tra}_U \end{array}$
$\begin{array}{ccc} \mathcal{P}_1(U^\bullet) & \xrightarrow{(\text{tra}_U, g)} & T' \\ p \downarrow & & \\ \mathcal{P}_1(X) & & \end{array}$	$\begin{array}{ccc} \mathcal{P}_2(U^\bullet) & \xrightarrow{(\text{tra}_U, g, f)} & T' \\ p \downarrow & & \\ \mathcal{P}_2(X) & & \end{array}$

Table 4: We generalize 1-anafunctors to **2-anafunctors** by regarding an anafunctor as an instance of **descent data** or **transition data**.

make a tetrahedron 2-commute

$$\begin{array}{ccc} p_2^* \text{tra}_U & \xrightarrow{p_{23}^* g} & p_3^* \text{tra}_U \\ p_{12}^* g \nearrow & p_{123}^* f \searrow & \nearrow p_{13}^* g \\ p_1^* \text{tra}_U & \xrightarrow{p_{14}^* g} & p_4^* \text{tra}_U \\ & \Downarrow p_{134}^* f & \downarrow p_{34}^* g \\ & & p_3^* \text{tra}_U \end{array} = \begin{array}{ccc} p_2^* \text{tra}_U & \xrightarrow{p_{23}^* g} & p_3^* \text{tra}_U \\ p_{12}^* g \nearrow & p_{234}^* f \searrow & \nearrow p_{24}^* g \\ p_1^* \text{tra}_U & \xrightarrow{p_{14}^* g} & p_4^* \text{tra}_U \\ & \Downarrow p_{124}^* f & \downarrow p_{34}^* g \\ & & p_3^* \text{tra}_U \end{array}$$

We say tra is equipped with a smooth structure, if it is equipped with a fixed p -local i -trivialization such that all the transition data is smooth.

We may address these transitions as **differential 2-cocycles**, or, taking them as a categorification of anafunctors, as **2-anafunctors**.

For different choices of i , we find various structures invented by various authors:

Proposition 6 *Principal and line bundle gerbes with connection, as well as the differential 2-cocycles characterizing them, are p -local i -transition data for 2-*

morphism of parallel transport codomain			name of corresponding 2-anafunctor
$i : T' \rightarrow T$			p -local i -transition
$\Sigma(\Sigma(U(1)))$	$\xrightarrow{\text{Id}}$	$\Sigma(\Sigma(U(1)))$	Deligne 2-cocycle
$\Sigma(\Sigma(U(1)))$	\rightarrow	$\Sigma(1d\text{Vect}_{\mathbb{C}})$	line bundle gerbe with connection (Murray)
$\Sigma(\text{AUT}(H))$	$\xrightarrow{\text{Id}}$	$\Sigma(\text{AUT}(H))$	fake-flat Breen-Messing 2-cocycle
$\Sigma(\text{AUT}(H))$	\rightarrow	$\Sigma(H\text{BiTor})$	fake-flat Aschieri-Jurčo bibundle gerbe with connection

Table 5: **Whatever the entries of this table mean** – an explanation of which does not fit in here – the message is that **various higher structures** that people have studied **are secretly examples of 2-anafunctors**.

functors, with p a given surjective submersion and i as indicated in the following table.

The above table says in particular that smooth 2-functors

$$\text{tra} : \mathcal{P}_2(X) \rightarrow \Sigma(\text{AUT}(H))$$

are in bijection with pairs consisting a $\text{Lie}(\text{AUT}(H))$ -valued 1-form A and a $\text{Lie}(H)$ -valued 2-form B

$$\text{tra} : \begin{array}{ccc} 0 & \xrightarrow{\gamma_1} & x \\ \gamma_3 \downarrow & \swarrow & \downarrow \gamma_2 \\ y & \xrightarrow{\gamma_4} & x + y \end{array} \mapsto \begin{array}{ccc} \bullet & \xrightarrow{1+A(\gamma_1)+\dots} & \bullet \\ 1+A(\gamma_3)+\dots \downarrow & \swarrow 1+B(\gamma_1, \gamma_3)+\dots & \downarrow 1+A(\gamma_2)+\dots \\ \bullet & \xrightarrow{1+A(\gamma_4)+\dots} & \bullet \end{array},$$

such that

$$\beta = F_A + \text{ad}(B)$$

vanishes. This β is the **2-form curvature** or **fake curvature**. The true curvature

$$\begin{array}{ccc} \bullet & \xrightarrow{\text{tra}_A(\gamma_1)} & \bullet \\ \text{tra}_A(\gamma_3) \swarrow & & \swarrow \text{tra}_A(\gamma_2) \\ \bullet & \xrightarrow{\text{tra}_A(S_1)} & \bullet \\ \text{tra}_A(\gamma_3) \downarrow & \swarrow \text{tra}_A(S_2) & \downarrow \text{tra}_A(S_3) \\ \bullet & \xrightarrow{\text{tra}_A(S_4)} & \bullet \end{array} \rightarrow \begin{array}{ccc} \bullet & \xrightarrow{\text{tra}_A(\gamma_1)} & \bullet \\ \text{tra}_A(\gamma_3) \swarrow & & \swarrow \text{tra}_A(\gamma_2) \\ \bullet & \xrightarrow{\text{tra}_A(S_3)} & \bullet \\ \text{tra}_A(\gamma_3) \downarrow & \swarrow \text{tra}_A(S_4) & \downarrow \text{tra}_A(S_5) \\ \bullet & \xrightarrow{\text{tra}_A(S_4)} & \bullet \end{array}$$

is

$$H = d_A B.$$

People did consider connection data on gerbes that is not fake flat. By the above, this does not integrate to a parallel transport 2-functor with values in a 2-group.

But it does integrate to a pseudo 2-functor with values in a 3-group.

3 Parallel 3-Transport: Chern-Simons

As Danny Stevenson explains, we should expect general parallel transport with respect to a n -group G_n to involve the $(n+1)$ -group of **inner automorphisms**

$$\text{Ad}_q \quad : \quad G_2 \quad \rightarrow \quad G_2$$

of G_n .

For G a 1-group, G -1-transport is the same as $(\text{INN}(G) = (G \rightarrow G))$ -2-transport:

$$\text{curv}_A \quad : \quad \begin{array}{ccc} x_s & \xrightarrow{\gamma_1} & x_1 \\ \downarrow \gamma_3 & \searrow_S & \downarrow \gamma_2 \\ x_2 & \xrightarrow{\gamma_4} & x_t \end{array} \quad \mapsto \quad \begin{array}{ccc} \bullet & \xrightarrow{1+A(\gamma_1)+\dots} & \bullet \\ \downarrow 1+A(\gamma_3)+\dots & \searrow_{1+F_A(\gamma_3, \gamma_1)+\dots} & \downarrow 1+A(\gamma_2)+\dots \\ \bullet & \xrightarrow{1+A(\gamma_4)+\dots} & \bullet \end{array} .$$

Surfaces are sent to the integrated curvature of the parallel 1-transport.

But for G_2 a 2-group, $\text{INN}(G_2)$ -3-transport is inherently richer than G_2 -2-transport.

$\text{INN}(G_2)$ is no longer strict, but all nontrivial structure morphisms are unique. We can consider pseudo-2-transport

$$\mathcal{P}_2(X) \rightarrow \Sigma(\text{INN}(G_2))$$

that strictly respects horizontal and vertical composition by itself.

Proposition 7 *p -local $\text{Id}_{\Sigma(\text{INN}(G_2))}$ -trnsitions yield the full Breen-Messing co-cycle data, not restricted to vanishing fake curvature.*

We would like to understand $\text{INN}(G_2)$ for the case where $G_2 = \text{String}_G = (\hat{\Omega}_k G \rightarrow PG)$, the strict version of the string 2-group.

This is easiest at the level of Lie 3-algebras.

Proposition 8 *For any semisimple Lie algebra \mathfrak{g} , and any level $k \in \mathbb{N}$, there is a Lie-3-algebra*

$$\text{cs}(\mathfrak{g})$$

such that a 3-connection with values in that Lie 3-algebra is, locally, a \mathfrak{g} -valued 1-form A , a 2-form B and the Chern-Simons 3-form

$$H = k\text{CS}(A) + dB .$$

Proposition 9 *The Chern-Simons Lie 3-algebra sits inside the Lie 3-algebra of the inner automorphisms of the String 2-group*

$$\text{cs}(\mathfrak{g}) \xrightarrow{\subset} \text{Lie}(\text{INN}(\text{String}_G)) .$$

This indicates that $\text{INN}(\text{String}_G)$ -3-transport is in fact Chern-Simons 3-transport.

I expect that the above inclusion is in fact an equivalence, but this I could not prove yet.

If we consider Chern-Simons 2-gerbes *without* connection, the situation simplifies, since it is known that Chern-Simons 2-gerbes are characterized by having WZW transition 1-gerbes:

noticing that the maximal strict sub-3-group in $\text{INN}(\text{String}_G)$ is $(U(1) \circlearrowleft \hat{\Omega}_k G \rightarrow PG)$ we have

Proposition 10 *Transition gerbes for $(U(1) \rightarrow \hat{\Omega}_k G \rightarrow PG)$ -2-gerbes (without connection) are WZW gerbes.*