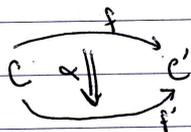


Learn the def. of

- category C
- functor $f: C \rightarrow C'$
- natural transformation



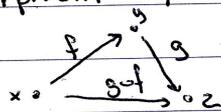
Duality

Every category C has an opposite category C^{op}

C^{op} has the same objects as C ,

but the morphisms are "turned around".

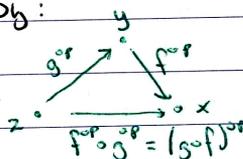
So there is a 1-1 correspondence between morphisms in C & morphisms in C^{op} , with $f: x \rightarrow y$ in C corresponding to a morphism $f^{op}: y \rightarrow x$ in C^{op} .



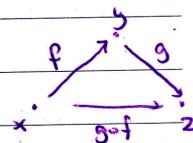
We compose morphisms in C^{op} by:

$$f^{op} \circ g^{op} = (g \circ f)^{op}$$

In C^{op} :



In C :



The study of how categories C relate to their partners C^{op} is called duality.

Note $(C^{op})^{op} = C$

'just like' for finite-dim vector spaces $(V^*)^* \cong V$

↳ natural isomorphism

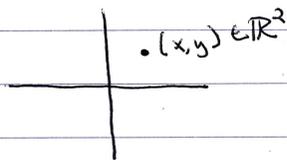
It turns out that the dual of geometry is algebra.

In geometry we study "points"; in algebra we study addition & multiplication.

Descartes realized we can reduce (a lot of) geometry to algebra:

this is called "analytic geometry"

We can associate to any finite-dimensional vector space V (over the real numbers) a commutative ring $\mathcal{O}(V)$ consisting of all polynomial functions on V , with usual addition & multiplication.



If $V = \mathbb{R}^n$, the algebra $\mathcal{O}(V)$ consists of polynomials in the coordinate functions x_1, \dots, x_n . $\mathcal{O}(V) = \mathbb{R}[x_1, \dots, x_n]$
polynomials in x_1, \dots, x_n

So: we go from a "space" V (a bunch of points) to an algebra $\mathcal{O}(V)$.
 Then we can describe certain subspaces of V :

$$X \xrightarrow{1:1} V$$

as quotient ~~algebras~~ ^{rings} of $\mathcal{O}(V)$:

$$\mathcal{O}(V) \xrightarrow{\text{onto}} \mathcal{O}(X) = \mathcal{O}(V) / \mathcal{I} \leftarrow \text{an ideal}$$

Example the unit circle is a subspace of the plane:

$$S' \longrightarrow \mathbb{R}^2$$

$$\text{where } S' = \{(x, y) : x^2 + y^2 - 1 = 0\}$$

Then there is an algebra $\mathcal{O}(S')$ of polynomial functions on the unit circle, with $\mathcal{O}(S') = \mathbb{R}[x, y] / \langle x^2 + y^2 - 1 \rangle$ ← the ideal generated by $x^2 + y^2 - 1$

So the 1-1 map $S' \longrightarrow \mathbb{R}^2$ gets turned around, giving $\mathcal{O}(\mathbb{R}^2) \longrightarrow \mathcal{O}(S')$ which is just restriction: $f \in \mathcal{O}(\mathbb{R}^2)$ gives $f|_{S'} \in \mathcal{O}(S')$.

Moreover $f, g \in \mathcal{O}(V)$ restricted to the same function on S' iff

$$f - g \in \langle x^2 + y^2 - 1 \rangle$$

$$\text{meaning } f - g = (x^2 + y^2 - 1)h \text{ for some } h \in \mathcal{O}(V).$$

Algebraic geometry is the study of geometry using commutative rings.

Our idea is: subspaces of V should correspond to quotient rings of $\mathcal{O}(V)$, or ideal of $\mathcal{O}(V)$

Problems:

1) What about $\langle x^2 + y^2 + 1 \rangle \subseteq \mathcal{O}(\mathbb{R}^2)$

$$\{ (x^2 + y^2 + 1)h : h \in \mathcal{O}(\mathbb{R}^2) \}$$

The function $x^2 + y^2 + 1$ doesn't vanish on \mathbb{R}^2 , so it seems the subspace of \mathbb{R}^2 corresponding to the ideal is \emptyset .

But there's another, simpler ideal that corresponds to $\emptyset \subseteq \mathbb{R}^2$.

Namely $\langle 1 \rangle$.

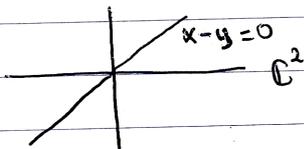
$$\emptyset = \{ (x, y) : 1 = 0 \}$$

We're getting 2 ideals corresponding to the same subspace.

One way out: use \mathbb{C} instead of \mathbb{R} .

2) Alas, using \mathbb{C} doesn't completely fix the problem:

2 different ideals can correspond to the same subspace



There is a (complex) line in \mathbb{C}^2 given by $x=y$, with ideal

$$\langle x-y \rangle \subseteq \mathbb{C}[x, y]$$

But $(x-y)^2$ also vanishes only on this line, so we're getting a different ideal defining the same subspace $\langle (x-y)^2 \rangle \subseteq \mathbb{C}[x, y]$

Algebraic geometers came up with a way around this...
but Grothendieck came along and found a better solution.

He cut the Gordian knot, and defined a new kind of space called an affine scheme such that the correspondence between algebra and geometry is perfect.

We're going to make up a category AffSch where objects are "affine schemes" and morphisms are maps between them, such that $\text{AffSch}^{\text{op}} = \text{CommRing}$

What is AffSch ?

Take op of both sides:

$$(\text{AffSch}^{\text{op}})^{\text{op}} = \text{CommRing}^{\text{op}}$$

$$\text{OR } \text{AffSch} = \text{CommRing}^{\text{op}}$$

Example the circle is an affine scheme, namely the comm. ring: $\mathbb{Z}[x,y]/\langle x^2+y^2-1 \rangle$

The plane is an affine scheme, $\mathbb{Z}[x,y]$

"The circle is included in the plane" means we have a homomorphism of comm. rings

$$\mathbb{Z}[x,y] \longrightarrow \mathbb{Z}[x,y]/\langle x^2+y^2-1 \rangle$$

namely the quotient map.

$$\text{We also have: } \mathbb{R}[x,y] \longrightarrow \mathbb{R}[x,y]/\langle x^2+y^2-1 \rangle$$

In "noncommutative geometry" we try to invent some new kind of "space" so that

$$\text{AffSch} = \text{CommRing}^{\text{op}}$$

sets generalized to something like

$$??? = \text{Ring}^{\text{op}}$$

Geometry

Algebraic geometry:

$$\mathbb{C} = [\text{affine schemes}]$$

Topology

$$\mathbb{C} = [\text{compact Hausdorff spaces}]$$

Set theory

$$\mathbb{C} = [\text{sets}]$$

Commutative Algebra

Ring Theory

$$\mathbb{C}^{\text{op}} = [\text{commutative rings}]$$

 \mathbb{C}^* -algebra Theory

$$\mathbb{C}^{\text{op}} = [\text{commutative } \mathbb{C}^* \text{-algebras}]$$

Logic

$$\mathbb{C}^{\text{op}} = [\text{atomic Boolean algebras}]$$

Look at $\mathbb{C}X_{\text{Haus}} = [\text{compact Hausdorff space, continuous maps}]$ From a compact Hausdorff space X on algebra

$$\mathbb{C}(X) = \{f: X \rightarrow \mathbb{C} : f \text{ is continuous}\}$$

This is an algebra with

$$(f+g)(x) = f(x) + g(x)$$

$$(fg)(x) = f(x)g(x)$$

$$(cf)(x) = cf(x) \quad c \in \mathbb{C}$$

It's a commutative algebra

It's a $*$ -algebra with $(f^*)(x) = \overline{f(x)}$,meaning an algebra A with $*$: $A \rightarrow A$ s.t.

$$(f+g)^* = f^* + g^*$$

$$(fg)^* = g^* f^*$$

$$(cf)^* = \overline{c} f^*$$

Also $\mathbb{C}(X)$ has a norm $\|f\| = \sup_{x \in X} |f(x)|$ This makes sense since X is compact.This makes $\mathbb{C}(X)$ into a \mathbb{C}^* -algebra, meaning that:

$$\|fg\| \leq \|f\| \|g\|$$

$$\|f^*\| = \|f\|$$

$$\|f^* f\| = \|f\|^2$$

" \mathbb{C}^* -axiom "

$$\|f^* f\| = \sup_{x \in X} |(f^* f)(x)|$$

$$= \sup_{x \in X} |f(x)|^2$$

$$= \left(\sup_{x \in X} |f(x)| \right)^2$$

$$= \|f\|^2$$

So $\mathbb{C}(X)$ is a commutative \mathbb{C}^* -alg.

next: can we take a morphism $\varphi: X \rightarrow Y$ in CHaus, that is a cont. map, and turn it into a (homo)morphism of comm. C^* -algebras.

A homomorphism between C^* -algs, say $F: A \rightarrow B$, is a map s.t.

$$F(a+b) = F(a) + F(b) \quad a, b \in A$$

$$F(ab) = F(a)F(b)$$

$$F(ca) = cF(a) \quad c \in \mathbb{C}$$

$$F(a^*) = F(a)^*$$

$$\exists K > 0 \text{ s.t. } \|F(a)\| \leq K \|a\| \quad \forall a \in A$$

All these imply $\|F(a)\| = \|a\|$

So we get a category

Comm C^* Alg = [comm. C^* -algebras, C^* -algebra homomorphisms]

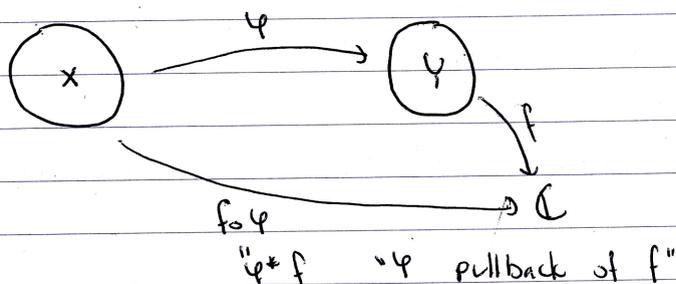
How does a cont. map $\varphi: X \rightarrow Y$ between compact Hausdorff spaces give a C^* -alg. homo. between $C(X)$ and $C(Y)$?

We'll get one,

$$\varphi^*: C(Y) \longrightarrow C(X)$$

$$\text{by: } \varphi^*(f)(x) = f(\varphi(x)) \quad f \in C(Y), \quad x \in X$$

$$\text{or: } \varphi^*(f) = f \circ \varphi$$



This is why algebra is the "dual" of geometry - it goes backwards:

$$\varphi: X \longrightarrow Y \quad \text{gives} \quad \varphi^*: C(Y) \longrightarrow C(X)$$

$$\text{Also } (\varphi \circ \psi)^* = \varphi^* \circ \psi^* \quad (\text{check this})$$

So we're getting a functor:

$$C: \text{Chtaus} \longrightarrow \text{Comm } C^* \text{Alg}^{\text{op}}$$

$$X \longmapsto C(X)$$

$$\varphi: X \rightarrow Y \longmapsto \varphi^*: C(Y) \rightarrow C(X)$$

Gelfand-Naimark Thm: This functor is an equivalence of categories.

I.e. there's a functor going back:

$$\text{Spec}: \text{Comm } C^* \text{Alg}^{\text{op}} \longrightarrow \text{Chtaus}$$

s.t. $\text{Spec} \circ C \cong \text{Id}_{\text{Chtaus}}$ and $C \circ \text{Spec} \cong \text{Id}_{\text{Comm } C^* \text{ Alg}^{\text{op}}}$
↑ natural isomorphism

What's Spec?

Given a comm. C^* -alg. A , how do we get a space $\text{Spec}(A)$?

Let's do $A = C(X)$.

Then $\text{Spec}(C(X))$ should give back X .

How do we recover the points of X starting from $C(X)$?

What's a point in X ?

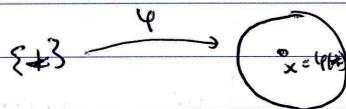
In terms of Chtaus, what's a point of X ?

It's a map $\varphi: \{*\} \rightarrow X$ where $\{*\}$ is the one-point space.

i.e. given $x \in X$ there's a map

$$\varphi: \{*\} \longrightarrow X$$

$$* \longmapsto x$$



↳ conversely any map $\varphi: \{*\} \rightarrow X$ determines a point in X .

Our functor $C: \text{Chtaus} \rightarrow \text{Comm } C^* \text{ Alg}^{\text{op}}$ will turn $\varphi: \{*\} \rightarrow X$ into a homomorphism

$$\varphi^*: C(X) \longrightarrow C(\{*\})$$

$$f \longmapsto f \circ \varphi$$

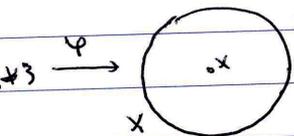
In fact $C(\{*\}) \cong \mathbb{C}$ where $g \in C(\{*\})$ gives $g(*) \in \mathbb{C}$

So we get

$$\varphi^*: C(X) \longrightarrow C(\{*\}) \xrightarrow{\sim} \mathbb{C}$$

$$f \longmapsto f \circ \varphi \longmapsto \underbrace{f \circ \varphi(*)}_{f(x)}$$

A point $x \in X$



gives a homomorphism $C(X) \rightarrow \mathbb{C}$
 $f \mapsto f(x)$

In short: any point $x \in X$ gives a homomorphism from $C(X)$ to \mathbb{C} called evaluation at x .

! pt \forall hom.

Lemma Distinct points of X give distinct homomorphisms $C(X) \rightarrow \mathbb{C}$.

("There are enough continuous functions to separate points" for a compact Hausdorff space)
 \rightarrow Stone-Weierstrass Theorem

! pt \forall hom.

Lemma Any C^* -alg. homomorphism $\Psi: C(X) \rightarrow \mathbb{C}$ comes from a point $x \in X$ via: $\Psi(f) = f(x) \quad \forall f \in C(X)$.

So we get a 1-1 correspondence between points $x \in X$ and homomorphisms $\Psi: C(X) \rightarrow \mathbb{C}$.

So given any comm. C^* -algebra A we define a set of points
 $\text{Spec}(A) = \{ \Psi: A \rightarrow \mathbb{C} : \Psi \text{ is a } C^* \text{-alg. homomorphism} \}$

There's a topology making $\text{Spec}(A)$ into a compact Hausdorff space.
In this topology Ψ_i converges to Ψ iff $\Psi_i(a) \rightarrow \Psi(a)$ for all $a \in A$

Finally, given a C^* -alg. homo. $F: A \rightarrow B$, how do we get a map of spaces $\text{Spec}(F): \text{Spec}(B) \rightarrow \text{Spec}(A)$?

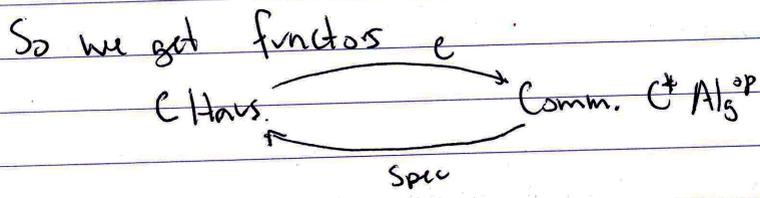
$$\text{Spec}(F)(\Psi)(a) = \Psi(F(a))$$

$$\Psi: B \rightarrow \mathbb{C} \quad C^* \text{-alg. homo.}$$

$$a \in A$$

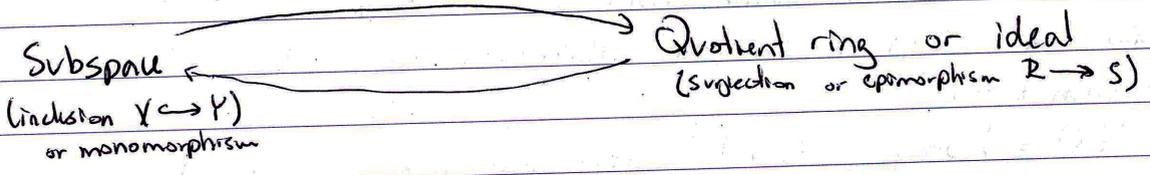
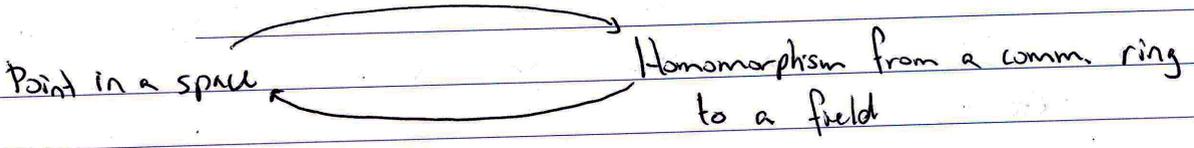
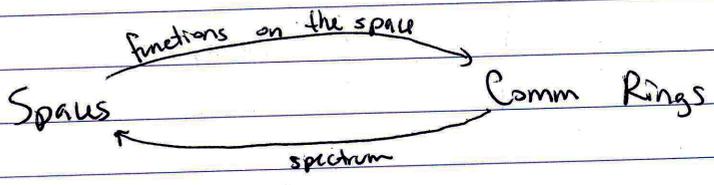
$$F(a) \in B$$

$$\Psi(F(a)) \in \mathbb{C}$$



which are inverses (up to nat. iso.)

Note



The duality between set theory & logic.

The basic idea:

operations on subsets of X	logical operations
\cup	\vee (or)
\cap	\wedge (and)
\emptyset	$F = 0$
X	$T = 1$
complement	\neg (not)

e.g. $x \in S \cup T \iff x \in S \text{ or } x \in T$
etc.

Let's fit this into the same mold as last time.

We saw that given a compact Hausdorff space X , we get a commutative C^* -alg.

$$C(X) = \text{hom}_{\text{Top}}(X, \mathbb{C}) \quad (\text{continuous maps from } X \text{ to } \mathbb{C})$$

Also, given a comm. C^* -algebra A , we get a compact Hausdorff space

$$\text{Spec}(A) = \text{hom}_{\text{Comm } C^* \text{ Algs}}(A, \mathbb{C}) \quad (\text{homomorphisms of } C^* \text{-algebras})$$

We call \mathbb{C} here a dualizing object:

"the same object x in 2 different categories \mathcal{C} & \mathcal{D} " such that

$$\mathcal{C} \longrightarrow \mathcal{D}^{\text{op}}$$

$$c \longmapsto \text{hom}_{\mathcal{C}}(c, x) \in \mathcal{D}$$

for some mysterious reason...

&

$$\mathcal{D}^{\text{op}} \longrightarrow \mathcal{C}$$

$$d \longmapsto \text{hom}_{\mathcal{D}}(d, x) \in \mathcal{C}$$

for some mysterious reason...

are inverses,

so \mathcal{C} & \mathcal{D}^{op} are equivalent.

Similarly we'll relate sets & Boolean algebras using the dualizing object

$$2 = \{0, 1\} \cong \{F, T\}$$

Boolean algebras are a little bit like C^* -algebras, but:

Comm. C^* -algebras	Boolean algebras
\mathbb{C}	$2 = \{F, T\}$
$+$	\vee
\cdot	\wedge
0	F
1	T

So, given a set X , we'll make a Boolean algebra
 $2^X = \text{hom}_{\text{set}}(X, 2)$

An element here is a fn. $f: X \rightarrow \{0, 1\}$.

Any subset $S \subseteq X$ gives such a fn:

$$\chi_S(x) = \begin{cases} 0 & x \notin S \\ 1 & x \in S \end{cases}$$

& conversely any such function gives a subset of X , so 2^X is just another way to think of the power set of X .

The operations $\cup, \cap, \overset{c}{}$ on subsets of X corr. to operations \vee, \wedge, \neg on functions $f: X \rightarrow \{0, 1\}$

$$\chi_{S \cup T} = \chi_S \vee \chi_T$$

$$\text{where } \chi_S \vee \chi_T(x) = \chi_S(x) \vee \chi_T(x)$$

& so on.

$(2^S, \vee, \wedge, 0, 1)$ will be a Boolean algebra:

$$\text{Note } S \subseteq T \iff \chi_S \leq \chi_T$$

i.e. if $\chi_S(x) = 1$ then $\chi_T(x) = 1$

Note " \leq " is not a separate concept:

$$\chi_S \leq \chi_T \iff \chi_S \wedge \chi_T = \chi_S$$

$$\iff \chi_S \vee \chi_T = \chi_T$$

Def A partially ordered set (A, \leq) is called a lattice if every pair $a, b \in A$ has a least upper bound $a \vee b$ & greatest lower bound $a \wedge b$, and also a least element $0 = F$ and a greatest element $1 = T$.

Def A distributive lattice is one where \wedge & \vee distribute over each other.

Def A Boolean algebra is a distributive lattice A where every element $x \in A$ has a complement $\neg x$ such that

$$x \wedge \neg x = F \quad x \vee \neg x = T$$

(If a complement exists, it's unique.)

Ex For any set S , 2^S is a Boolean algebra with pointwise defined \leq :
given $f, g \in 2^S$ we say $f \leq g$ if $f(x) \leq g(x) \quad \forall x \in S$
It thus has pointwise defined $\vee, \wedge, 0, 1, \& \neg$,
e.g. $(\neg f)(x) = \neg f(x)$

Alas, not every Boolean algebra is isomorphic to one of this form!

The Boolean algebras of the form 2^S are "complete atomic Boolean algebras".

Def A complete Boolean alg. A is one where every subset $S \subseteq A$ has a l.u.b. $\bigvee_{x \in S} x$ and g.l.b. $\bigwedge_{x \in S} x$, and these distribute over each other.

Def An atom in a Boolean algebra A is an element $x \in A$ st. $x \neq 0$ and if $y < x$ then $y = 0$.

Ex In 2^S the atoms correspond to the elements of S , or singletons $\{s\} \subseteq S$.

Def A Boolean algebra A is atomic if $\forall x \in A, x = \bigvee_{y \in A} y$ where $y_i \in A$ are atoms.

There's a category CABA of complete atomic Boolean algebras & homomorphisms of complete Boolean algebras:

$$\varphi: A \rightarrow B \text{ preserving } \vee, \wedge, 0, 1, \top, \perp$$

There's a category Set of sets and functions.

Thm Set is equivalent to CABA^{op} via these functors

$$\text{Set} \longrightarrow \text{CABA}^{\text{op}}$$

$$S \longmapsto 2^S = \text{hom}_{\text{Set}}(S, 2)$$

and

$$\text{CABA}^{\text{op}} \longrightarrow \text{Set}$$

$$A \longmapsto \text{hom}_{\text{CABA}}(A, 2)$$

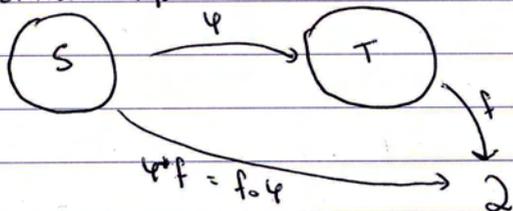
where $2 = \{F, T\}$ is a CABA in the obvious way.

Given a function $\varphi: S \rightarrow T$ we get a complete Boolean alg. homo.

$$\varphi^*(f)(s) = f(\varphi(s))$$

$$f \in 2^T \text{ i.e. } f: T \rightarrow 2 \quad s \in S$$

This is called a pullback:



Ex $\{f \in L^\infty[0,1] : f(x) = \overset{F}{0} \text{ or } \overset{T}{1} \text{ for all } x \in X\}$

This is a complete but not atomic Boolean algebra under pointwise operations.

(There are no atoms)

Geometry

Algebraic Geometry (affine schemes)

Topology (compact Hausdorff spaces)

Set Theory

?

?

?

Algebra

Commutative Rings

Commutative C^* -algebras

Complete atomic Boolean algebras (Proposition Logic)

Boolean algebras

Linear algebra (finite-diml vector spaces)

Finite abelian groups

The opposite of the category of all Boolean algebras is the category of Stone spaces: compact Hausdorff spaces that are totally disconnected: every open set is closed (the vice versa)

The Boolean algebra of a Stone space X consists of its open subsets,

with $A \cup B$ as " \vee "

$A \cap B$ as " \wedge "

A^c as " \neg "

Let FinVect be the category of finite-dimensional vector spaces over favorite field (e.g. \mathbb{R}) & linear maps.

What's $\text{FinVect}^{\text{op}}$?

A typical morphism in FinVect is $T: \mathbb{R}^n \rightarrow \mathbb{R}^m$.

A morphism in $\text{FinVect}^{\text{op}}$ is thus $T^{\text{op}}: \mathbb{R}^m \rightarrow \mathbb{R}^n$

suspiciously similar to the transpose $T^t: \mathbb{R}^m \rightarrow \mathbb{R}^n$ in FinVect .

In fact, we have an equivalence $\text{FinVect} \xrightarrow{\sim} \text{FinVect}^{\text{op}}$, with

$T: \mathbb{R}^n \rightarrow \mathbb{R}^m$ in FinVect \longmapsto $T^t: \mathbb{R}^m \rightarrow \mathbb{R}^n$ in FinVect

or $(T^t)^{\text{op}}: \mathbb{R}^n \rightarrow \mathbb{R}^m$ in $\text{FinVect}^{\text{op}}$

We can also get the equivalence $\text{FinVect} \cong \text{FinVect}^{\text{op}}$ using $\mathbb{R} \in \text{FinVect}$ as our dualizing object:

$\text{FinVect} \xrightarrow{\quad} \text{FinVect}^{\text{op}}$

$V \longmapsto \text{hom}(V, \mathbb{R}) = V^*$

$T: V \rightarrow W \longmapsto T^t: W^* \rightarrow V^*$ in FinVect

$(T^t)^{\text{op}}: V^* \rightarrow W^*$ in $\text{FinVect}^{\text{op}}$

So, FinVect straddles the worlds of geometry & algebra, being its own opposite.

Also, the category [finite abelian groups, group homomorphism] is its own "op".

Galois Theory

Galois theory is secretly about dualities between posets.

Def A poset is a partially ordered set (S, \leq) where \leq is reflexive, transitive, and antisymmetric: $x \leq y$ & $y \leq x \Rightarrow x = y$.

If (S, \leq) is a poset, we get a category with elements of S as objects & there exists a unique morphism $f: x \rightarrow y$ iff $x \leq y$ ($x, y \in S$), and no morphisms $f: x \rightarrow y$ otherwise.

In fact, the categories we get this way are precisely those with:

- 1) at most 1 morphism from any object x to any object y
- 2) if there are morphisms $f: x \rightarrow y$ & $g: y \rightarrow x$, then $x = y$.

So to a category theorist, a poset is a category with these 2 properties.

Given categories of this kind, a functor is really just an order preserving map $f: (S, \leq) \rightarrow (T, \leq)$, i.e. a function s.t. $x \leq y$ in $S \Rightarrow f(x) \leq f(y)$ in T .

Given a category of this sort coming from the poset (S, \leq) , its opposite comes from the poset (S, \leq^{op}) where $x \leq^{op} y$ iff $y \leq x$. We'll write $x \geq y$ for $x \leq^{op} y$.

What are adjoint functors between categories of this sort?

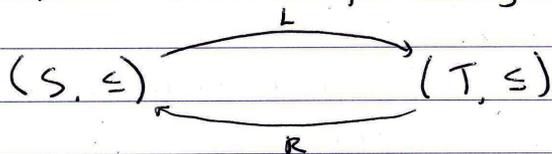
Def Given categories C, D , we say a functor $L: C \rightarrow D$ is the left adjoint of a functor $R: D \rightarrow C$, or R is the right adjoint of L , if there is a natural 1-1 correspondence

$$\text{hom}_D(Lx, y) \cong \text{hom}_C(x, Ry) \quad \forall x \in C, y \in D$$

Ex Let $L: \text{Set} \rightarrow \text{Grp}$ send any set S to the free group on S
 and $R: \text{Grp} \rightarrow \text{Set}$ send any group G to its underlying set.
 Here $\text{hom}_{\text{Grp}}(LS, G) \cong \text{hom}_{\text{Set}}(S, R(G))$

L = liberty!
 = freedom

Ex What are adjoint functors between posets (S, \leq) & (T, \leq) ?
 It's a pair of order-preserving functions



such that: $Lx \leq y \iff x \leq Ry$

This comes from $\text{hom}_b(Lx, y) \cong \text{hom}_c(x, Ry)$

Def A pair of adjoint functors between posets is called a Galois correspondence.

Thm Suppose $(S, \leq) \overset{L}{\underset{R}{\rightleftarrows}} (T, \leq)$ is a Galois correspondence. Then we get an order-preserving map $RL: (S, \leq) \rightarrow (S, \leq)$.

Let's write \bar{x} for RLx .

Then $x \leq \bar{x} \quad \forall x \in S$

$(Lx \leq Lx \Rightarrow x \leq RLx)$

and $\bar{\bar{x}} = \bar{x} \quad \forall x \in S$

So we say $\bar{}$ is a closure operator on the poset (S, \leq) .

Similarly write y° for LRy .

Then $y^\circ \leq y \quad \forall y \in T$

and $(y^\circ)^\circ = y^\circ \quad \forall y \in T$

So $^\circ$ behaves like the "interior" operation on subsets of a top. space — it's a closure operator on $(T, \leq)^\text{op}$.

Finally, L & R give a bijection between closed elements of S

(meaning $x \in S$ w/ $\bar{x} = x$) & open elements of T (meaning $y \in T$ s.t. $y^\circ = y$).

Galois Theory

Suppose you have any kind of algebraic gadget - a set with some operations obeying some axioms.

e.g. monoids, ^{+, -, 0} groups, ^{+, x, 0, 1} rings, fields

Then we can define a "subgadget" of a gadget K to be a subset $k \subseteq K$ closed under all the operations.

The gadgets F with
 $k \subseteq F \subseteq K$

form a poset with \subseteq as the partial ordering. Let's call this poset D .

Galois theory uses groups to study D .

Any gadget K has a group $\text{Aut}(K)$ of "automorphisms", i.e. 1-1 & onto functions $g: K \rightarrow K$ that preserve all the operations,

e.g.

$$\begin{aligned} g(x+y) &= gx+gy \\ g(xy) &= (gx)(gy) \\ g(0) &= 0 \\ g(1) &= 1 \end{aligned}$$

We say an element $x \in K$ is fixed by $g \in \text{Aut}(K)$ if $gx = x$.

We say a subgadget $F \subseteq K$ is fixed by $g \in \text{Aut}(K)$ if $gx = x$ for each $x \in F$.

Note the subset $\{g \in \text{Aut}(K) : g \text{ fixes } F\}$ is a subgroup of $\text{Aut}(K)$.

The subgroup of $\text{Aut}(K)$ fixing the subgadget $k \subseteq K$ is called the Galois group $G(K|k)$.

Let C be the poset of subgroups of $G(K|k)$, where the partial ordering is \subseteq . The idea is to use C to study D .

We'll do this by constructing a Galois correspondence $C \begin{matrix} \xrightarrow{L} \\ \xleftarrow{R} \end{matrix} D^{\text{op}}$
 i.e. order-preserving maps obeying
 $L_G \subseteq F \iff G \supseteq R_F$

What's R ?

It maps gadgets $k \subseteq F \subseteq K$ to subgroups of the Galois group $G(K|k)$

It works as follows:

$$R_F = \{g \in \text{Aut}(K) : g \text{ fixes } F\}$$

To show $R: D^{op} \rightarrow C$ is order preserving (i.e. a functor) we need:

$$k \subseteq F \subseteq F' \subseteq K \Rightarrow R(F) \supseteq R(F')$$

This is true: that if g fixes F' & $F \subseteq F'$ then g fixes F .

What's L ?

It maps subgroups $G \in \mathcal{G}(K|k)$ to gadgets between k & K .

It works as follows

$$LG = \{x \in K : G \text{ fixes } x\}$$

$$= \{x \in K : \forall g \in G \ g \text{ fixes } x\} \leftarrow \text{Note: this is a subobject of } K!$$

To show $L: C \rightarrow D^{op}$ is order-preserving we need:

$$G \subseteq G' \subseteq \mathcal{G}(K|k) \Rightarrow LG \supseteq LG'$$

This is true: it says that if $x \in F$ is fixed by all $g \in G'$ then it's fixed by all $g \in G$ (some $G \subseteq G'$)

Next: why is $L: C \rightleftharpoons D^{op}: R$ a Galois connection?

i.e., why is $LG \subseteq F \iff G \supseteq RF$

$LG \subseteq F$ means ^{every element of F} everything fixed by G is in F

$G \supseteq RF$ means everything fixing F is in G .

These are just two ways of saying the same thing.

Now we can relate nice subgadgets $k \subseteq F \subseteq K$ & nice subgroups $G \in \mathcal{G}(K|k)$ using the theorem we saw last time...

but now let's stick in an "op".

Thm Suppose $L: C \rightleftharpoons D^{op}: R$ is a Galois connection. Define

$$\bar{c} = RLc$$

$$c \in C$$

$$\bar{d} = LRd$$

$$d \in D^{op}$$

These are closure operators:

$$c \subseteq \bar{c} \quad \& \quad \bar{c} = \overline{\bar{c}}$$

$$d \subseteq \bar{d} \quad \& \quad \bar{d} = \overline{\bar{d}} \quad (\text{where } \subseteq \text{ is ordering on } D)$$

We say $c \in C$ is closed if $c = \bar{c}$, and similarly for $d \in D$. L & R give a 1-1 correspondence between closed elements of C & closed elements of D .

[If we would have done C^{op} instead, we would get open operators.]

In our application, what's a "closed" subgadget $k \subseteq F \subseteq K$?

It's one with $F = LRF$

$$= L \{g \in \text{Aut}(K) : g \text{ fixes } F\}$$

$$= \{x \in K : x \text{ is fixed by all } g \text{ that fixes } F\}$$

So a subgadget F is closed if it contains all $x \in K$ that are fixed by all $g \in G(K/k)$ that fix F .

What's a "closed" subgroup $G \subseteq G(K/k)$?

$$G = RL G$$

$$= R \{x \in K : x \text{ is fixed by } G\}$$

$$= \{g \in G(K/k) : gx = x \text{ for all } x \text{ fixed by } G\}$$

So: a subgroup G is closed if it's the group of all $g \in G(K/k)$ that fix all x fixed by G .

So: the hard part of Galois theory includes:

1) finding a more concrete characterization of the "closed" subfields $k \subseteq F \subseteq K$

2) Similarly for the closed subgroups

3) Understanding the poset \mathcal{C} — the poset of the Galois groups.

Groupoids

Def A morphism $f: x \rightarrow y$ in a category has an inverse $g: y \rightarrow x$ if $fg = 1_y$ & $gf = 1_x$

If f has an inverse, it's unique so we write it as f^{-1} .

A morphism with an inverse is called an isomorphism.

If there's an isomorphism $f: x \rightarrow y$ we say x & y are isomorphic.

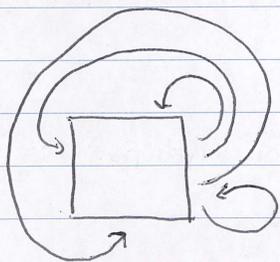
Def A groupoid is a category where all morphisms are isomorphisms.

Ex Any group G gives a groupoid with one object, $*$, and morphisms $g: * \rightarrow *$ corresponding to elements $g \in G$, with composition coming from multiplication in G .

Conversely any 1-object groupoid gives a group.

So "a group is a 1-object groupoid."

More generally, if C is any category & $x \in C$, the isomorphisms $f: x \rightarrow x$ form a group under composition, called the automorphism group $\text{Aut}(x)$.



$$\text{Aut}(\square) \cong \mathbb{Z}_4$$

(rotational symmetries)

Ex Given any category C there's a groupoid, the core of C , C_0 , whose objects are those of C & whose morphisms are the isomorphisms of C , composed as before.

Ex If $\text{Fin Set} = [\text{finite sets, functions}]$

then $\text{Fin Set}_0 = [\text{finite sets, bijections}]$

and if n is your favorite n -element set, $\text{Aut}(n) = S_n$, the symmetric group

So Fin Set_0 "unifies" all the symmetric groups

Ex Suppose G is a group acting on a set X :

$$\alpha: G \times X \longrightarrow X$$

$$(g, x) \longmapsto gx$$

Often people form the set X/G , the quotient set where an element $[x]$ is an equivalence class of elements $x \in X$ where $x \sim y$ iff $y = gx$ for some $g \in G$.

But a "better" thing is to form the translation groupoid $X//G$, where objects are elements $x \in X$

a morphism from x to y is a pair (g, x) where $g \in G$ and $gx = y$.

$$x \xrightarrow{(g, x)} y$$

the composite of

$$x \xrightarrow{(g, x)} y$$

and

$$y \xrightarrow{(h, y)} z$$

is

$$x \xrightarrow{(hg, x)} z$$

In X/G , we say x & y are equal if $gx = y$

In $X//G$, we say they are isomorphic, or more precisely, we have a chosen isomorphism $(g, x): x \rightarrow y$.

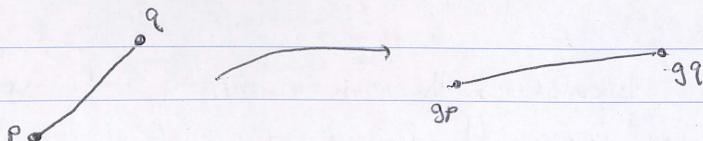
To a first approximation, a "moduli space" is a set X/G , given some obvious topology, while a "moduli stack" is a group $X//G$, where the sets of objects & morphisms have topologies.

Ex Let X be the set of line segments in the Euclidean plane. Let G be the Euclidean group of the plane:

$$\text{all bijections } g: \mathbb{R}^2 \rightarrow \mathbb{R}^2 \text{ s.t. } |gp - gq| = |p - q|$$

i.e. g preserves distances

G acts on X :



More precisely, $X = \mathbb{R}^2 \times \mathbb{R}^2$ & G acts on it via $g(p, q) = (gp, gq)$

We're not counting (p, q) as the same as (q, p) .

We're allowing $p = q$.

(continued)

X/G is the "moduli space of line segments"

$$X/G \cong [0, \infty)$$

There's a line segment (p, q) and also a line segment (q, p) , but $(p, q) \sim (q, p)$. So they have some equivalence class: they give a single point in X/G .

Next consider $X//G$

Now objects are line segments & morphisms are like:

$$(p, q) \xrightarrow{(g, p, q)} (p', q')$$

where $gp = p'$ & $gq = q'$, as in the picture.

Given a groupoid C , we can form:

1) the set \underline{C} of isomorphism classes of objects:

$$[x] \text{ where } [x] = [y] \text{ iff } x \cong y.$$

2) for any $[x] \in \underline{C}$, a group $\text{Aut}(x)$, where x is any representative of $[x]$. (Note if $x \cong y$, then $\text{Aut}(x) \cong \text{Aut}(y)$ as groups.)

Thm Given a groupoid C , we can recover C (up to equivalence) from \underline{C} and all the groups $\text{Aut}(x)$ (one for each isomorphism class in \underline{C}).

Ex $C = \text{Fin Set}$

$$\underline{C} \cong \mathbb{N}$$

and for each $n \in \mathbb{N}$ we get a group S_n which we've seen is (iso. to) $\text{Aut}(x)$ for any $x \in \text{Fin Set}$ with n elements.

Ex $C = X//G$ where X is the set of line segments & G is the Euclidean group of the plane.

$$\underline{C} \cong [0, \infty)$$

In general, $X//G = X/G$ because both are names for the set of equivalence classes $[x]$ where $x \sim y$ iff $y = gx \exists g \in G$.

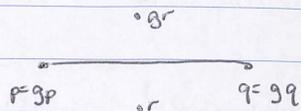
But $X//G$ has more information, namely all the automorphism groups $\text{Aut}(x)$, one for each equivalence class.

So in our example, what's $\text{Aut}(\langle p, q \rangle)$?

It's \mathbb{Z}_2 if $p \neq q$, since there is

a reflection preserving $\langle p, q \rangle$

If $p = q$ it's group $O(2)$ of all
orthogonal 2×2 matrices, i.e. all
rotations & reflections fixing $p \in \mathbb{R}^2$.



Moduli Spaces & Moduli Stacks

Given a groupoid \mathcal{C} , let $\underline{\mathcal{C}}$ be the set of isomorphism classes of objects. Often $\underline{\mathcal{C}}$ will have the structure of a space (e.g. a topological space, a manifold, an algebraic variety, a scheme, ...). Then $\underline{\mathcal{C}}$ is called a moduli space.

[Ex] If G is a group acting on a set X , we get a groupoid $X//G$, the translation groupoid, where:
 objects are elements of X
 morphisms $x \xrightarrow{(g,x)} y$ are pairs $x \in X, g \in G$, where $y = gx$.
 Then $X//G \cong X/G$ where X/G has elements $[x]$ with $x \sim y$ when $y = gx$ for some $g \in G$.

Recall

[Thm] The groupoid $X//G$ is equivalent to the groupoid with:

- one object $[x]$ for each $[x] \in X/G$
- one morphism $f: [x] \rightarrow [x]$ for each morphism $f: x \rightarrow x$ where x is any chosen representative of the equivalence class $[x]$.

If $[x] \neq [y]$ there are no morphisms between them.

We often call X/G a moduli space, and $X//G$ the moduli stack.

Last time we looked at an example:

[Ex] "The moduli stack of line segments" in Euclidean geometry.

Here $X = \mathbb{R}^2 \times \mathbb{R}^2 \ni (p, q)$, $G = O(2) \times \mathbb{R}^2$

Here G is the Euclidean group of the plane and we think of (p, q) as a line segment with a chosen 1st & 2nd endpoint, which can be equal.

Then the moduli space is $X/G \cong [0, \infty)$ the space of lengths.

$$[(p, q)] \mapsto |p - q|$$

The moduli stack $X//G$ keeps track of symmetries: $\text{Aut}[(p, q)] \cong \text{Aut}((p, q))$

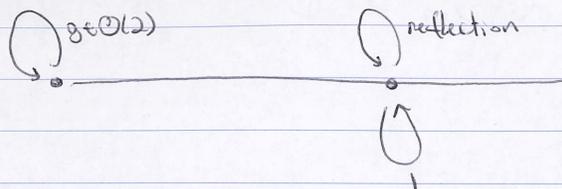
is the subgroup of G consisting of all $g \in G$ with $(gp, gq) = (p, q)$

$$\text{Aut}((p, q)) \cong \mathbb{Z}/2 \quad \text{if } p \neq q \quad \begin{array}{c} p \longrightarrow q \\ \end{array}$$

$$\text{Aut}((p, q)) \cong O(2) \quad \text{if } p = q \quad \begin{array}{c} p \circlearrowleft q \\ \end{array}$$

$$\cong SO(2) \times \mathbb{Z}_2$$

So the moduli stack looks like



Ex "The moduli stack of triangles"

Let G = the Euclidean group as before, but now let X be the set of triangles:

$$X = \mathbb{R}^2 \times \mathbb{R}^2 \times \mathbb{R}^2$$

These are triangles with named vertices that can be equal.

The moduli space X/G is the set of isomorphism classes of triangles.

$$\text{Now } X/G \cong [0, \infty)^3$$

$$[(p, q, r)] \mapsto (|p-q|, |q-r|, |r-p|)$$

Here it seems that if p, q, r are all distinct, (p, q, r) has as automorphisms only the identity.

If we define a triangle to be an unordered triple of points in \mathbb{R}^2 , an equilateral triangle would have S_3 as automorphisms, and isosceles would have $S_2 = \mathbb{Z}/2$.

This gives a more interesting moduli stack.

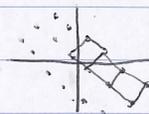
Ex A Riemann surface is a 2-dim. smooth manifold with charts $\varphi_i: U_i \rightarrow \mathbb{C}$ s.t. $\varphi_i \circ \varphi_j^{-1}$ is analytic (=holomorphic)

- Every Riemann surface that's homeomorphic to the plane is isomorphic (as a Riemann surface) to \mathbb{C} .
- Every Riemann surface homeomorphic to the sphere is isomorphic to the Riemann sphere $\mathbb{C}P^1 \cong \mathbb{C} \cup \{\infty\}$

There are lots of nonisomorphic ways to make a torus into a Riemann surface - these are elliptic curves.

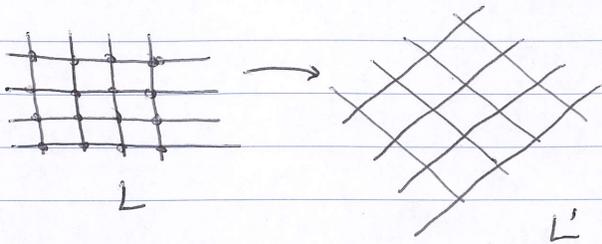
Every elliptic curve is isomorphic to one of this form:

take a lattice $L \subseteq \mathbb{C}$, i.e. a subgroup of $(\mathbb{C}, +, \partial)$ that's isomorphic to \mathbb{Z}^2 , and form \mathbb{C}/L , getting a torus with obvious charts $\varphi_i: U_i \rightarrow \mathbb{C}$, and thus an elliptic curve.



When do 2 lattices L & L' give isomorphic elliptic curves: $\mathbb{C}/L \cong \mathbb{C}/L'$?

Answer: iff $L' = \alpha L$ for some nonzero $\alpha \in \mathbb{C}$



There's a groupoid \mathcal{C} with

- elliptic curves as objects
- isomorphisms of Riemann surfaces as morphisms

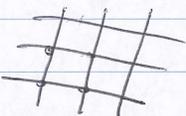
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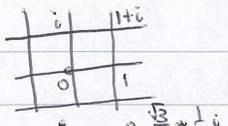
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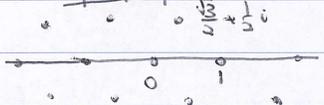
where X is the set of lattices & $G = \mathbb{C}^*$ (nonzero complex numbers with multiplication)

So X/G is called the moduli space of elliptic curves, and $X//G$ is the moduli stack of elliptic curves.

There are 2 elliptic curves with a bigger automorphism group:

typical elliptic curve  has $\mathbb{Z}/2$ as symmetries
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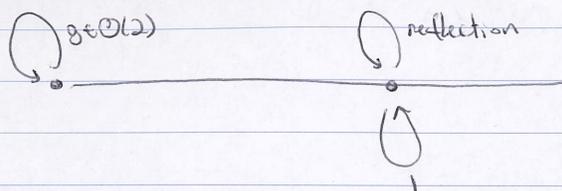
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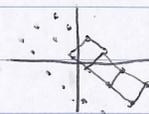
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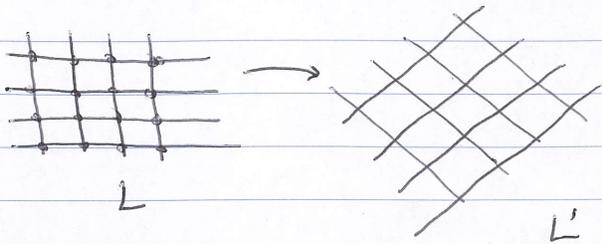
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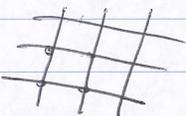
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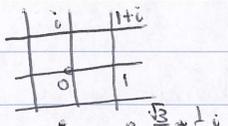
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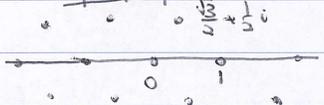
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Klein Geometry

We've seen that:

- a geometry is a group G
- a type of figure in this geometry is a subgroup $H \subseteq G$
- the set of figures of that type is G/H : a homogeneous G -space

How can we do geometry this way?

We need G -invariant relations between figures.

Ex projective plane geometry:

$$G = PGL(3, \mathbb{R})$$

$$X = \{\text{lines through the origin in } \mathbb{R}^3\} = \{\text{pts in } \mathbb{RP}^2\}$$

X is a homogeneous G -space, so

$X \cong G/H$ where $H \subseteq G$ is the stabilizer of your favorite point $p \in X$:

$$H = \{h \in G : hp = p\}$$

An invariant relation between points is a relation, i.e. a subset $R \subseteq X \times X$ s.t. $(p, q) \in R \Rightarrow (gp, gq) \in R$ for all $p, q \in X, g \in G$

But the only invariant relations in this example are

$$p = q \quad \text{and} \quad p \neq q$$

because distance is not preserved by G .

More interestingly, let $Y = \{A \subseteq B : A \text{ is a 1-dim. subspace of } \mathbb{R}^3 \text{ \& } B \text{ is a 2-dim. subspace of } \mathbb{R}^3\}$
 $= \{\text{flags}\}$

where a flag is a point $p \in \mathbb{RP}^2$ lying on a line $L \in \mathbb{RP}^2$

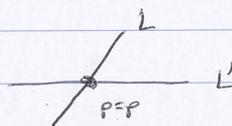


G acts transitively on Y (even the Euclidean group does), and there are various invariant relations between flags, i.e. subsets $R \subseteq Y \times Y$ invariant under G .

For example:

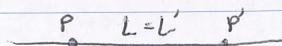
One invariant relation between (p, L)

& (p', L') says " $p = p'$ and $L \neq L'$ "

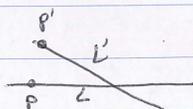


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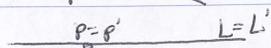
Or: " $L=L'$ and $p \neq p'$ "



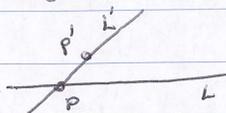
Or: " $p \neq p'$ and $L \neq L'$ "



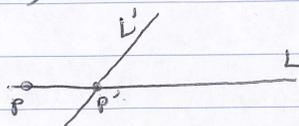
Or: " $p = p'$ and $L \neq L'$ "



Or: " $p \in L'$ but $L \neq L'$ and $p \neq p'$ "



Or: " $p' \in L$ but $L \neq L'$ and $p \neq p'$ "



All 6 of these relations are visible here:



"6 flags"

For any group G , we can make up a category $G\text{Rel}$ where:

• objects are G -sets

• morphisms are invariant relations

where an invariant relation $R: X \rightarrow Y$ from the G -set to the G -set Y is a relation, i.e. a subset $R \subseteq X \times Y$ such that:

$$(x, y) \in R \Rightarrow (gx, gy) \in R \quad \forall x \in X, y \in Y, g \in G$$

How do we compose morphisms?

Given any relations $R: X \rightarrow Y$ and $S: Y \rightarrow Z$ (not nec. invariant)

we can compose them to get $S \circ R: X \rightarrow Z$:

$$S \circ R = \{(x, z) \in X \times Z : \exists y \in Y \text{ s.t. } (x, y) \in R \text{ and } (y, z) \in S\}$$

If R & S are invariant so is $S \circ R$.

There's a category Rel where

• objects are sets

• morphisms are relations

Here

$$\text{hom}(X, Y) = 2^{X \times Y}$$

Recall: for any set S , 2^S is a CABA: a complete atomic boolean algebra, with

$$\begin{array}{lcl} \subseteq & \text{as} & \leq \\ \cap & \text{as} & \wedge (=g\wedge b) \\ \cup & \text{as} & \vee (=l\vee b) \\ \subset & \text{as} & \neg \end{array}$$

So in Rel , $\text{hom}(X, Y)$ is not merely a set, it's a CABA. The same is true for GRel : e.g. if $R: X \rightarrow Y$, $S: X \rightarrow Y$ are invariant, so is $R \cap S$, $R \cup S$, R^c

In fact Rel & GRel are "CABA-enriched categories"

What's an enriched category?

In category theory we want to overthrow the tyranny of sets: instead of working in Set all the time, we try to prove results that hold in many categories. But the very definition of category use sets:

A category is a class of objects, and for each pair of objects x, y a set $\text{hom}(x, y)$, and a composition function

$$\circ: \text{hom}(x, y) \times \text{hom}(y, z) \rightarrow \text{hom}(x, z)$$

etc...

The idea in enriched cat. theory is to generalize, replacing Set by some other category V and say:

A V -enriched category is a class of objects, and for each pair of x, y an object $\text{hom}(x, y) \in V$, and a composition morphism in V :

$$\circ: \text{hom}(x, y) \otimes \text{hom}(y, z) \rightarrow \text{hom}(x, z)$$

etc...

Here we need V to be a "monoidal category", i.e. a category with some sort of "tensor product" \otimes .

It turns out that CABA's form a monoidal category, so it makes sense to talk about a CABA-enriched category, & Rel & GRel are such.

Enriched Categories & Internal Monoids

A monoid is "the same" as a 1-object category: if you have a category \mathcal{C} with one object x , there's a monoid $\text{hom}(x, x)$ with multiplication $\circ: \text{hom}(x, x) \times \text{hom}(x, x) \rightarrow \text{hom}(x, x)$

Conversely given a monoid M you can build a category with one object x and $\text{hom}(x, x) = M$, with composition being multiplication in M .

More generally suppose V is a monoidal category, i.e. a category with a tensor product: $\otimes: V \times V \rightarrow V$ obeying some rules.

Then recall a V -enriched category \mathcal{C} has a class of objects and for any objects $x, y \in \mathcal{C}$, a "hom-object" $\text{hom}(x, y) \in V$ & composition morphisms: $\circ: \text{hom}(x, y) \otimes \text{hom}(y, z) \rightarrow \text{hom}(x, z)$

A 1-object V -enriched category is the same as a monoid internal to V , or monoid in V , i.e. an object $M \in V$ with a multiplication $m: M \otimes M \rightarrow M$ that's associative and unital.

Ex Suppose $V = \text{AbGrp}$ with the usual tensor product of abelian groups. Then a monoid in V is called a ring. (ring has unit, but ring does not)
It's an abelian group M , with a multiplication $m: M \otimes M \rightarrow M$ an abelian group homomorphism, i.e. a function $m: M \times M \rightarrow M$ that's linear in each argument:

$$(a+b) \cdot c = a \cdot c + b \cdot c$$

$$a \cdot (b+c) = a \cdot b + a \cdot c$$

Ex If $V = \text{RMod}$ for some comm. ring R , a monoid in V is called an R -algebra.

Ex If $V = \text{Top}$ with usual product \times of top'd spaces as \otimes , a monoid in V is called a topological monoid.

Back to our favorite example: Klein geometry

Let G be a group, and let $G\text{Rel}$ be the category with G -sets as objects
 G -invariant relations as morphisms.

This a CABA-enriched category. So if we take one object, i.e. one G -set X , we can form a 1-object CABA-enriched category with X as the only object $\text{hom}(X, X)$ is the only homset, or "hom-CABA".

Ex Projective plane geometry

Take $G = \text{PGL}(3, \mathbb{R})$

$Y = \{\text{flags}\}$

$= \{(p, L) : p \in \mathbb{R}^3 \text{ is a 1-dim'l subspace,}$
 $L \in \mathbb{R}^3 \text{ is a 2-dim'l subspace,}$
 $p \subseteq L\}$

$\text{hom}(Y, Y)$ is a monoid in CABA. What is it like?

Instead of describing all the elements, let's just describe the atoms.

In general, given any group G and any G -sets X, Y , what are the atoms in $\text{hom}(X, Y)$ like?

They're invariant relations $R: X \rightarrow Y$

i.e. $R \subseteq X \times Y$ s.t. $(x, y) \in R \Rightarrow (gx, gy) \in R$

But they are the smallest nonempty subsets of this form. So, any atom R must contain a point (x, y) , and thus all points of the form (gx, gy) with $g \in G$. Indeed, any orbit $\{(gx, gy) : g \in G\} \subseteq X \times Y$ is an atom in $\text{hom}(X, Y)$.

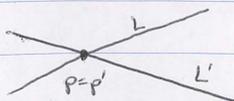
So, if

$G = \text{PGL}(3, \mathbb{R})$

$Y = \{\text{flags}\}$

the atoms in $\text{hom}(Y, Y)$ are the orbits of G acting on $Y \times Y$.

e.g. the orbit of this

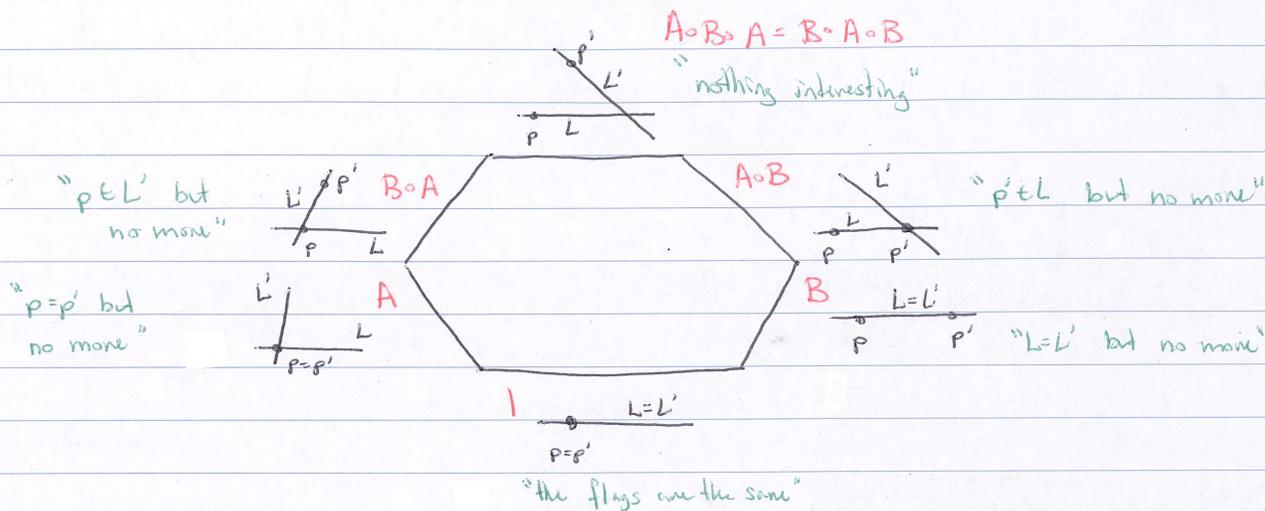
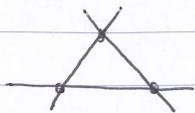


$x = (p, L)$

$y = (p', L')$

is the set of all pairs of flags sharing the point (and no more!).

Last time we saw all 6 atoms in $\text{hom}(Y, Y)$:



The identity $I \in \text{hom}(Y, Y)$ is "two flags are the same"

Note we can compose invariant relation & $I \circ I = I$

Let $A \in \text{hom}(Y, Y)$ be "having the same point but no more"

$$A \circ A = A \vee I$$

If you change the line on a flag twice, the result could be changing the line or getting back the original flag.

Let $B \in \text{hom}(Y, Y)$ be "having the same line but no more" and "changing the point"

$$B \circ B = B \vee I$$

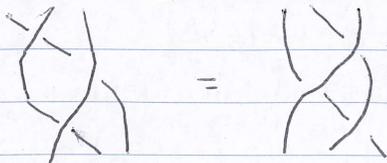
$A \circ B$ is one of our atoms, " $p' \in L$ but no more"

$B \circ A$ is another atom, " $p \in L'$ but no more"

$A \circ B \circ A = B \circ A \circ B$ is "nothing interesting"

In fact this is a presentation for our monoid in CABA, $\text{hom}(Y, Y)$.

If we draw A as $X|$ and B as $|X$ then $A \circ B \circ A = B \circ A \circ B$ is called the "3rd Reidemeister move" or "Vang-Baxter equation":



This is the only relation in B_3 , the 3-strand braid group.