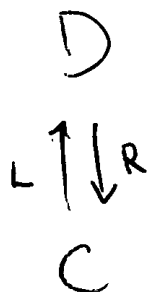
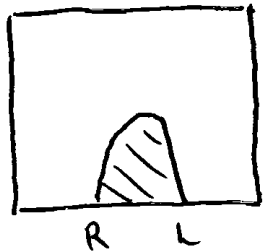


The Bar Construction

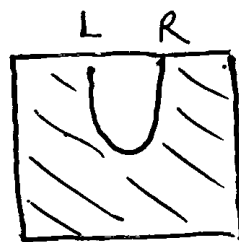
Suppose we have an adjunction:



We get a unit η & counit ε :



unit $\eta: I_C \Rightarrow RL$



counit $\varepsilon: LR \Rightarrow I_D$

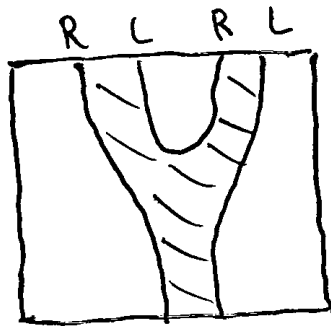
satisfying zig-zag identities. Last time

we saw this gives a monad on C ,

i.e. a monoid in $\text{End}(C)$, namely

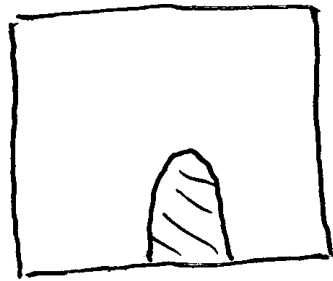
$RL \in \text{End}(C)$.

This is indeed a monoid:



multiplication

$$RLRL \xRightarrow{RL} RL$$



unit

$$I_c \xRightarrow{\epsilon} RL$$

satisfying associativity: $Y = Y$

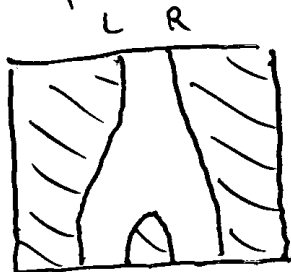
l/r unit laws: $Y = I = Y$

We also get a comonoid on D , i.e.

a comonoid in $\text{End}(D)$, namely $LR \in \text{End}(D)$.

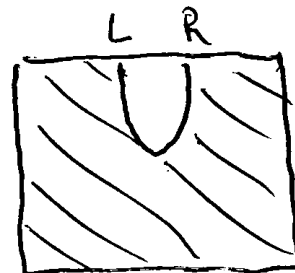
A "comonoid" is just like a monoid,

but upside down:



comultiplication

$$LR \xRightarrow{\epsilon} LRLR$$



counit

$$LR \xRightarrow{\epsilon} 1_D$$

These satisfy coassociativity: $\Delta = \Delta$
i.e. l/r comit laws: $\lambda = \lambda = \lambda$

More tersely, if M is a monoidal category, M^{op} is also monoidal with same \otimes , and:

Def - A comonoid in M is a monoid in M^{op} .

E.g. - An algebra is a monoid in Vect ,
a coalgebra is a comonoid in Vect ,
or monoid in Vect^{op} .

So: our adjunction gives a comonad (R) ,
which is a monoid in $\text{End}(D)^{\text{op}}$.

Since Δ is the free monoidal category on a monoid, this gives a monoidal functor

$$\Delta \longrightarrow \text{End}(D)^{\text{op}}$$

i.e. a monoidal functor

$$\Delta^{op} \xrightarrow{\alpha} \text{End}(D)$$

and thus a simplicial object in $\text{End}(D)$!

Taking a specific object $d \in D$, we get:

$$\begin{aligned} \text{ev}_d: \text{End}(D) &\longrightarrow D \\ F &\longmapsto Fd \end{aligned}$$

so we get a simplicial object in D :

$$\Delta^{op} \xrightarrow{\alpha} \text{End}(D) \xrightarrow{\text{ev}_d} D$$

This simplicial object in D is called

$\underline{d}: \Delta^{op} \rightarrow D$; we call this the

bar construction.

Moral: given an adjunction $\begin{matrix} D \\ \lrcorner \\ C \end{matrix} \hat{=} R$,

any object $d \in D$ gives a simplicial object \underline{d} in D .

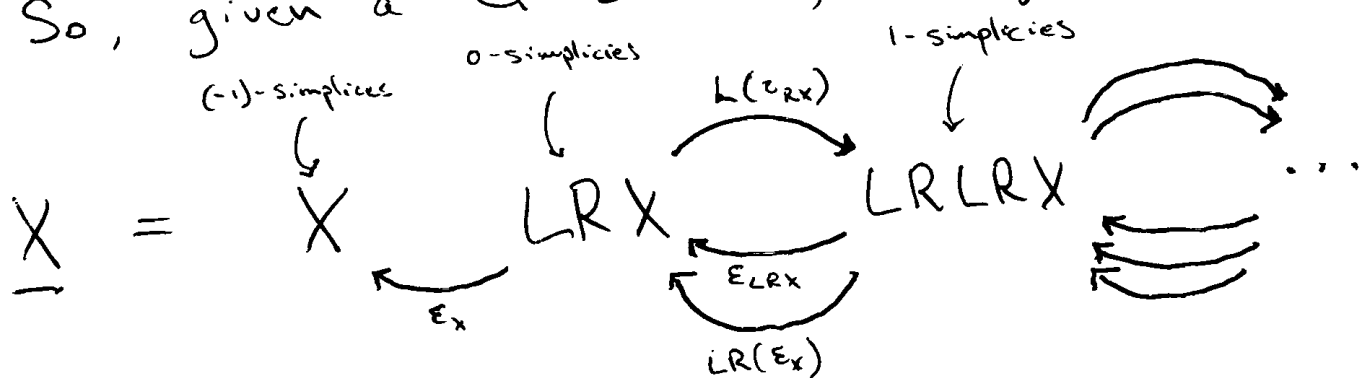
Example: The Cohomology of Groups

Here we take a group G ; get an adjunction

$$\begin{array}{c}
 G\text{-Set} \\
 \begin{array}{c} \uparrow L \\ \downarrow R \end{array} \\
 \text{Set}
 \end{array}$$

where $G\text{-Set}$ is the category of sets w. left G -action.

So, given a $G\text{-Set } X$, we get:



a simplicial G -set!

What's a 1-simplex, or 2-simplex, in this simplicial G -set like?

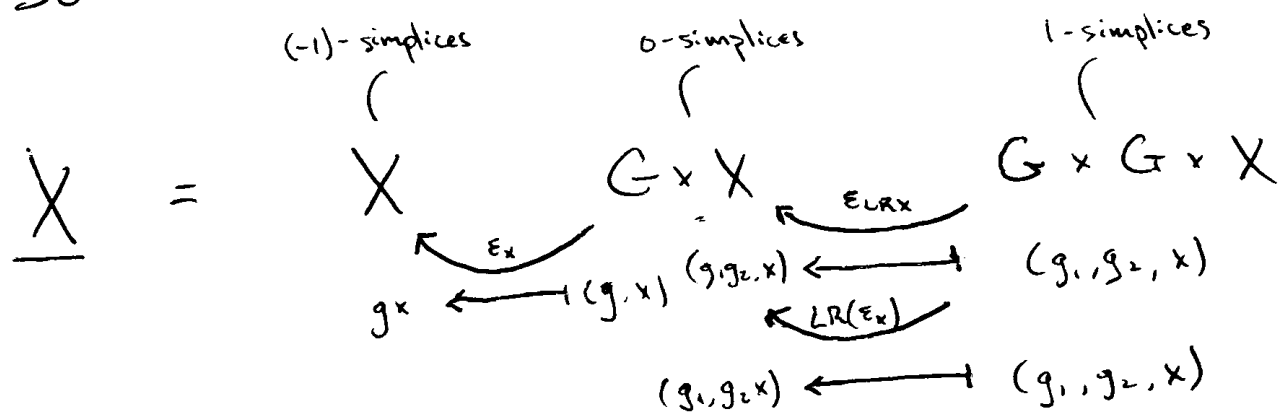
Given our G -set X , what's LRX ?

RX is usually just called " X " - the underlying set of our G -set X . LRX has elements " gx " for each $g \in G$; $x \in X$.

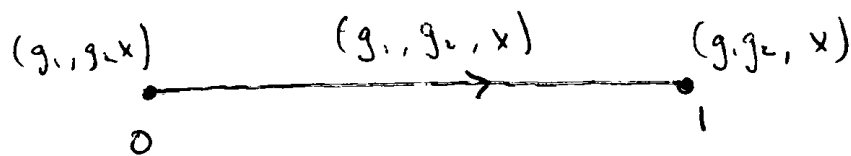
Really " gx " is just (g, x) , so $LRX = G \times X$, which is a G -set with

$$g_1 \cdot (g_2, x) = (g_1 g_2, x).$$

So

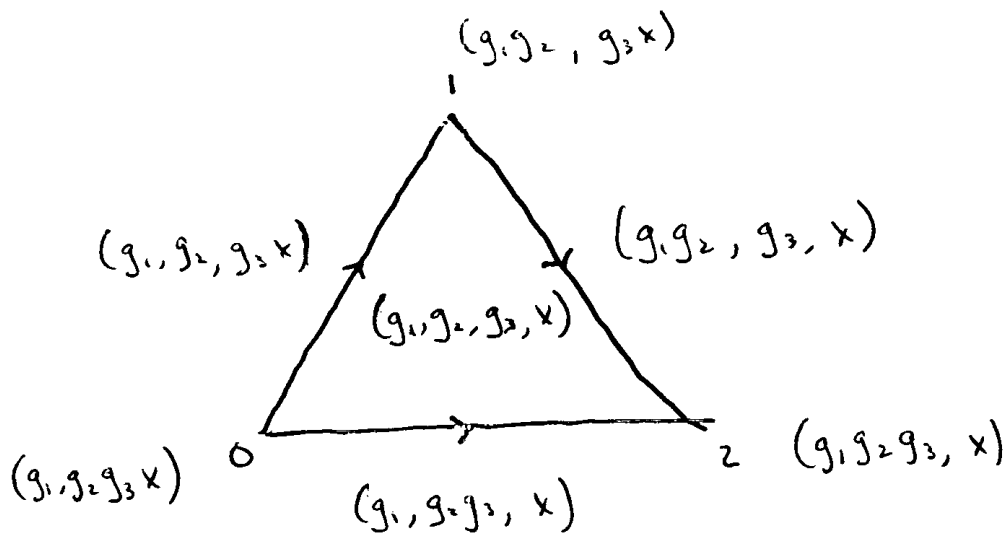


So, a typical 1 -simplex in \underline{X} looks like:

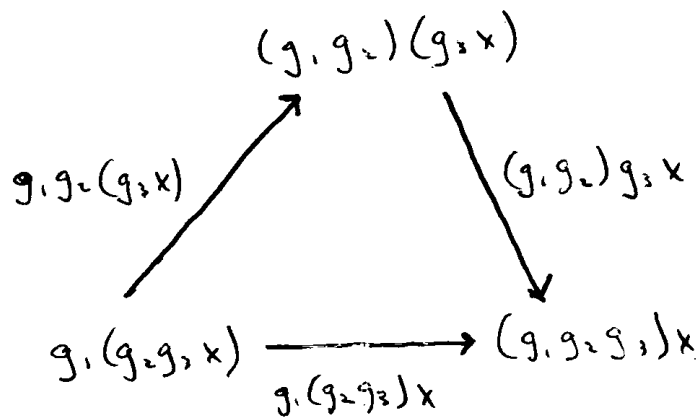


Note: both 0-simplices here have as face the (-1)-simplex $g_1 g_2 x$. So this 1-simplex is a proof that $g_1(g_2 x) = (g_1 g_2)x$ — the 2 formal expressions $(g_1, g_2 x)$ & $(g_1 g_2, x)$ evaluate via ε_x to the same element of X , namely $g_1 g_2 x$.

How about a 2-simplex in X ?



Here we see 2 proofs that $g_1(g_2 g_3 x) = (g_1 g_2 g_3)x$:



Using one step or 2. The triangle is a "metaproof" or "syzygy" - a "homotopy between proofs".

The simplicial G -set \underline{X} is called EG when $X = *$. In general \underline{X} has contractible components, one for each element of X - these are the 1-simplices of \underline{X} !

So EG has one contractible component.

It's like a "puffed-up point" - a contractible space on which G acts freely.