L7-1 Parallel Projection A' Limutually 11 B' M&N transversals C' M&N transversals First observation A*B*C=>A*B*C Proof B his instrip between L, +L3 (see HW). BELZ-side Lin Li-side Lz (dy alstrip) Either B'= Bor BB'= L2. Ineithercase, B' & Eprevious] => A & B' & C'. Thop 1 |AB|= |AC| => |A'B'|= |A'C'|. Proof. Cases A=A', B=B' or C=C' A A to switchingroler Moin of variables MENM, Cast A=A' Take N through B' N 11Mg C' 1×ABB'1=1×ACC'1 (Corresp × S) 1×AB'BI=1×AC'C1 1×AB'BI=1×AC'C1 1×AB'BI=1×AC'C1 M. KM2 Do MillMz artend DENALZ AdC'on oppsidur N ALC on Same Side N ⇒ C+C'on opp sides N. hence C+D+C'

17-2 This implie BB, D, C ventues of I, Rad LACC'= LACD and B'DC' are cowelp J.S. Henrie |AB| = |BD((= |BC|) and Also recall | AB'B | =12 AC e' = [ZABB'] = [A C'B'D] Hence (asa) / AC'C by $\triangle ABB' \cong \triangle B'DC'$, so that remarks on prev. page. 1A'B']=|B'C'| Care B=B' A A' B=B' |AB|= |BC| given 1×ABA'I= 1×CBC' Vertical angles 1×AA'CI= 1×CC'AI Alternate Interna Angles. So A ABA' = ACBC' by ASA, and hence IA'B=A'B'I= [C'B=C'B'] (mo condition Case EA, B, C'S ~ EA, B'C'Z= & Swhether (M, 11 M, orNO. AVA' Choose NIIM2 SO B-JB" B' AEN C-JC" C' Han N=M Otherwise M, NM2 AEM, which is not the case

27-3 Previous reasoning shows [A'B''] = 1B''C'']. But now we have I's AA'B''B + B''B'C''C, So 1AB" = 1A'B', 1B" C" = 1B'C' and finally 1A'B' = 1AB" = B'C" = B'C" = Notebook Papartheorem Li, My as before with AieLinMi, AitAitAinertAme. Bi ELinMz. Then B1 # ... * Bmo. Furtherwore, if 1A1 Ail= |AzA1 | for all i, then (BiBi 1=1B2B1 for alli, A. Dr. Apply breeding A. B. Apply breeding A. B. Successively to A. B. Linglight Rational Proportionality Thun - 2, Lz Lz as before ArcMinLi, BrEMinLi- DE 1A2A31 is rational, then IB2B31 = 1A2A31 1A,A21 IS rational, TR,B21 = 1A,A21

17-4. By algebra we also have 1B, B2 [B2B3] [A,A2] [A2A3]= Proof Lot $|A_1 A_3| = p$, so $\frac{|A_1 A_3|}{|A_1 A_2|} = q$, so $\frac{|A_1 A_3|}{|A_1 A_2|} = q$ Previous results shows that [Ying Y, 1= v all i, so also by letweenness $\begin{aligned} |A_1, A_2| &= qw, |A_2A_3| &= pu \\ |B_1B_2| &= qv, |B_2B_3| &= pv, \end{aligned} \qquad \begin{array}{c} |B_2B_3| \\ \hline B_1B_2| &= qv, \\ \hline B_2B_3| &= pv, \end{array} \end{aligned}$ I mational ArAzi Need more sophisticated MARZI Need more sophisticated ideal, took Greeks about 200 years to find a solution.

L7-5 Condition of Eudoxus OLX, yETR. Then X=y => (a) every positive rational & LX patisfies & Ly. (b) Same with in equality directions reversed. Key rolea If Olxly there is some I & X < Z < y . (Look at decimal expansions, for example? Del Moi'le 11.3-11.4 fe-dotailes read & understand passively! Similar Triangles ABCN ADEF vertices (Similar, with ratio of sim ilitude=r) =? IXARCI=1XDEF1, 1XBCA1=14DFE, 1XCAB1=14FDE1 and <u>DEI DEI IDEI IEEI V(Some</u>). IABI IACI IBCI (r). Adstract stuff: MABC = DDEFC= DABC= DDEF AABC~ AABC, AABC~ ADEF ⇒ ADEF~ AABC AARC.~ ADEF & ADEF~ DGHK = DAARCNAGHK

17-6

A.A. Similarity Theorem Given AABC+ ADEF S.J. IXBACI=1XEDFI, 1XABCI=1XDEFI, Then DABENDDEF. Proof. We also have [ZACDI= ZDFE]. Let v = IDEI IARI Nour Suppose r < 1. Switching the valos of the triangles if neec., can assume r > 10 A Take B'E (AB so that My B C AB' = 1 (AB). Let hine L through B' So that L 11 BC Notice A+B+B' (r>1) CLAIM L 2. Not par allel; otherwise LHAC+LIBC => BCHAC) false since these meet at C. Suppose L+AC meet at C# Then LEB'-side AC + A+B+B'=) C'+ A on opporter

17-7 So C, C'on Same side and A*C*C'. By prov. thm. $\frac{1CC'I}{IACI} = \frac{IBB'I}{IABI} = 4 - 1, so \frac{IAB'I}{IABI} = \frac{IABI + IBB'I}{IABI}$ $= \frac{IBB'I}{IABI} + 1, likewise \frac{IAC'I}{IACI} = \frac{ICC''I}{IACI} + 1, Since$ $= \frac{ICC'I}{IABI} = \frac{IBB'I}{IACI} + 1 - \frac{IAB'I}{IACI} = \frac{IAC'I}{IACI}$ Also & BAC= XB'A"C', [XARC]= |XAB'C'], |ZACB1= 1×AC"B"L. Combining with hypotheses, get 1×B'AC'L=1×EDFI, 1×AB'C'L= [*DEF], and IDE = ~ LABI = LABI, & AAB'C' > DEF by SAS. Will suffici to prove DABC~DDEF. Have shown AAA + IAB' | IAC'] Dewesubstitute IABI ZACI Dewesubstitute D -> E' Then ID*FT IDET IEFI IDFI E -> F' Then ID*FT IDET IEFI IDFI IAOT IATON IEFI IDET IBOU IACI » equivo [DE] IDFI IABI IAC Hence DABC ~ DAB'C' DEF. Construction yields Cov. Given DABC & rZ1. Then can find DAB'C' SO B'E (AB, C'E (AC, DABC ~ DAB'C'.