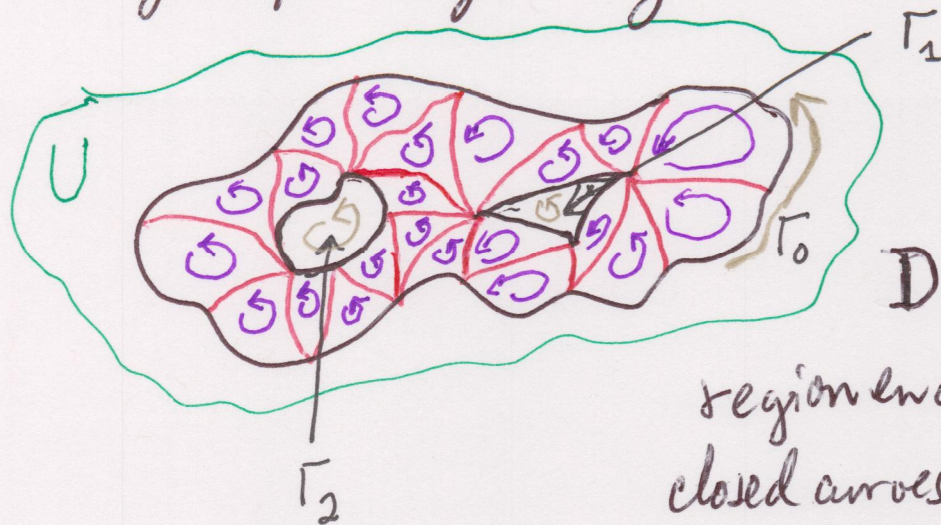


Triangulations

Motivation \vec{F} vector field on open set U containing closed, ^{bounded} region D whose boundary is finitely many nice curves Γ_i .

\vec{F}
Smooth



D is the region enclosed by the closed curves $\Gamma_0, \Gamma_1, \Gamma_2$

Intuitively, one curve Γ_0 is outermost (need to make this rigorous eventually!).

Parametrize Γ_i in the counterclockwise sense.

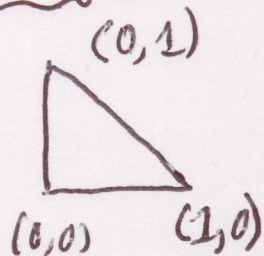
Standard topic in MV calculus

Prove Green's Theorem for D .

$$\vec{F} = (P(x, y), Q(x, y))$$

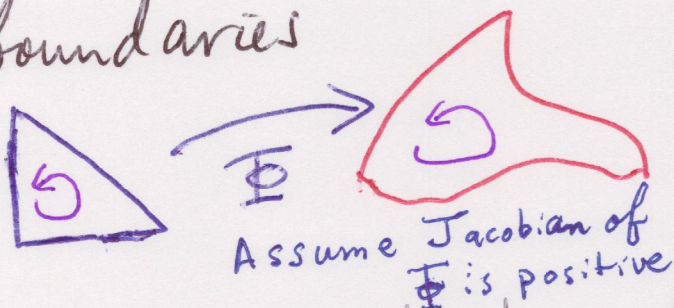
$$\left(\int_{\Gamma_0} - \sum_{i \geq 1} \int_{\Gamma_i} \right) (P dx + Q dy) = \iint_D \left(\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy$$

Idea ① Prove for triangular region



(Straight forward)

② Prove for regions with curved boundaries



Prev. case +
change of
variables
for double integrals

③ Cut D up into triangles as illustrated

(needs rigorous justification). Apply ② to each piece and add.

Ⓐ Sum of double integrals = integral over D .

and

② If we add the line integrals ^{over the Δs} then the integrals over the inside (red) pieces cancel each other as below

and all that's left is the sum of the line integrals over the boundary pieces.

Usually this can be done explicitly for any given example, but we would prefer a general proof.

Triangulation (2D)

$$X = \cup A_i$$

$$A_i \cong \text{triangle } \{x, y \geq 0, x+y \leq 1\}$$

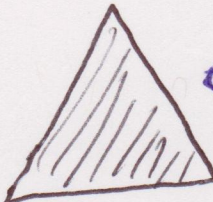
$i \neq j \Rightarrow A_i$ meets A_j in an edge.

Generalize formally to all finite dims.

Building blocks

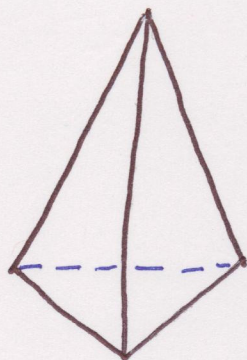
•
point

—
line
segments


solid
triangular
regions

Better than
Square
regions. Every
polygonal region
can be cut up
into Δ
regions

tetrahedra
(3D)



3D polyhedral regions
can be cut up into
tetrahedral regions.

<http://faculty.grsu.edu/goldenij/TriangularPrismDissection.html>
(interactive picture)

n-dimensions n-simplex in \mathbb{R}^n

$$\{(x_1, \dots, x_n) : x_i \geq 0, \sum x_j \leq 1\}$$

More symmetric version

In \mathbb{R}^{n+1} , take all points (x_1, \dots, x_{n+1})
 s.t. $x_i \geq 0$ and $\sum x_j = 1$

1st \rightsquigarrow 2nd Let $x_{n+1} = 1 - \sum_{i=1}^n x_i$.

2nd \rightsquigarrow 1st Drop last coord.

General n -dim simplex in \mathbb{R}^m .

Vertices v_0, \dots, v_n s.t.

$v_1 - v_0, \dots, v_n - v_0$ are linearly indep.

Take all points $\sum t_i v_i$ s.t. all $t_i \geq 0$

and $\sum t_j = 1$.

Prop. In above setting, \exists 1 way of writing
 such a point as $\sum s_i v_i$ with $s_i \geq 0$ and
 $\sum s_i = 1$.

Proof $x = \sum t_i v_i = \sum s_i v_i \Rightarrow$

$$x - v_0 = \sum t_i (v_i - v_0)^* = \sum s_i (v_i - v_0)$$

* Because $(\sum t_i v_i) - v_0 =$

Linear indep now
 $\Rightarrow t_i = s_i, i > 0$

$$\sum t_i v_i - \sum t_i v_0 = \sum_{i>0} t_i (v_i - v_0)$$

$$\boxed{\sum t_i = 1}$$

$$\boxed{t_0 (v_0 - v_0) = 0}$$

Then show
 $s_0 = t_0$

Continuation in

math 246A notes 2010. pdf

in the course directory.

see Section I.2, pp. 15-21.

Def A triangulation of a space X is a homeomorphism $P \rightarrow X$, where P is a simplicial complex. (see p. 17 of above file).